Week 10: Lecture B Hybrid Fuzzing II

Wednesday, March 20, 2024



How are projects going?

Problems?

Successes?







Questions?





Hybrid Fuzzing Recap



What is hybrid fuzzing?

Combining crude fuzzing with smarter fuzzing

- E.g., random + concolic execution (Driller, QSYM, Savior)
- E.g., random + taint tracking (VUzzer, RedQueen, Angora)

Goal is to balance strengths of both techniques

- Use generic fuzzing for most test cases
 - Use speed to brute-force easy branches
- Deploy more elegant approach selectively
 - Focus its **precision** on harder branches





How most hybrid fuzzers work

Leverage AFL-style **parallel fuzzing** mode with conventional fuzzer as parent





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What could go wrong?

Ineffective seed scheduling

- There are fundamental differences in **speed**
 - AFL can solve basic branch conditionals fast
 - Fancier approaches generally are much slower
- Heavyweight approaches are best applied to a **subset** of paths
 - Invoking on all paths will lead to path explosion
 - E.g., by the time it's solved, fuzzer is already way past

Questions?





Adventures in Hybrid Fuzzing: Driller



Fuzzing

I

0.	def f (x) {
1.	if x > 10 {
2.	if x < 100:
3.	print "You win!"
4.	else:
5.	<pre>print "You lose!"</pre>
6.	<pre>}else:</pre>
7.	print "You lose!"

1	\Rightarrow	"You lose!"

- 593 **⇒** "You lose!"
- 183 **⇒** "You lose!"
- 4 ⇒ "You lose!"
- 498 **⇒** "You lose!"
- 48 \Rightarrow "You win!"

Source: https://www.ndss-symposium.org/wp-content/uploads/2017/09/07_3-ndss2016-slides.pdf

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Where fuzzing falls short





 \Rightarrow



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Symbolic Execution to the rescue!



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Driller

- Idea: invoke concolic execution via demand launch
 - Heuristic 1: a pre-determined # of mutations based on test case length
 - Heuristic 2: after a pre-determined time interval without new coverage
- Concolic executor based on angr
 - Binary-level instrumentation and analysis framework
 - Heavily maintained and used in many research projects
 - Translates, analyzes binary in intermediate form (VEXIR)





Driller in action



AFL-found test cases

















Driller in action

Driller in action









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When to turn solving elsewhere?

When the path is already fully solved

- Track all branches and which have been solved
- A fundamental piece of info that is tracked



When to turn solving elsewhere?

When the path is already fully solved

- Track all branches and which have been solved
- A fundamental piece of info that is tracked
- When symbolic executor cannot solve
 - Biggest culprit: hashes







Questions?





Adventures in Hybrid Fuzzing: QSYM



Problem: relying on an IR is costly

Executor	chksum	md5sum	sha1sum	md5sum(mosml)
Native	0.008	0.014	0.014	0.001
KLEE	26.243	32.212	73.675	0.285
angr	-	-	-	462.418

Table 1: The emulation overhead of KLEE and angr compared to native execution, which are underlying symbolic executors of S2E and Driller, respectively. We used chksum, md5sum, and sha1sum in coreutils to test KLEE, and md5sum (mosml) [12] to test angr because angr does not support the fadvise syscall, which is used in the coreutils applications.



Source: https://taesoo.kim/pubs/2018/yun:qsym.pdf

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QSYM: operate on *native* **instructions**

Omit lifting to intermediate representation

Use Intel PIN dynamic binary instrumentation

Trade-offs:

- A much higher implementation complexity
- Significant decrease in symbolic instructions
 - 4X fewer than Driller







Problem: incomplete environment modeling





Problem: incomplete environment modeling



QSYM: leave the environment as-is

Omit translating the environment

- Use concrete execution to model it
 - Model only relevant system calls
 - E.g., standard input, reads, etc.
- What about kernel state forking?
 - Avoid—just **re-execute** from the start

Trade-offs:

- Re-execution adds more overhead
 - Cannot "go back in time" like Driller





Problem: overconstrained paths





Solver will try to solve these first

> Really just need to solve this last one

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Problem: overconstrained paths





Source: https://www.ndss-symposium.org/wp-content/uploads/2017/09/07_3-ndss2016-slides.pdf



QSYM: optimistically solve last constraint

```
0. def f (x) {
1. if x > 10 {
2.
    if (x > 1000){
        if x<sup>2</sup> == 152399025:
3.
           print "You win!"
4.
5.
        else:
6.
           print "You lose!
7.
      }else:
8.
         print "You lose!"
    }else:
9.
0.
      print "You lose!"
```



Trade-offs:

- Does not always work
- Can just let the fuzzer quickly rule these out

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Questions?





Adventures in Hybrid Fuzzing: RedQueen



Problem: symbolic and concolic execution is slow



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RedQueen's solution: input-to-state tracking

Idea: hook comparison instructions and identify their input bytes

Replace with compared-to value (lifted directly from the operand)





RedQueen's solution: input-to-state tracking

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Supporting other comparisons

Idea: hook comparison instructions and identify their input bytes

Replace with compared-to value (lifted directly from the operand)





What about checksums?

Finding these at the binary-level is difficult

- **Assumption:** can identify input bytes that affect the checksum hash
- **Colorize the input:** inject random bytes and see if they influence the outcome

```
if( u64(input) == hash(input[8..len]) )
```



What about checksums?

- Then, patch-out the checksum with an always-true operation
 - Assumption: checksum is only passed if the input is well-formed





What about checksums?

- Then, patch-out the checksum with an always-true operation
 - Assumption: checksum is only passed if the input is well-formed
 - Thus, skipping over checksum won't matter if well-formed
 - New input found afterwards? Great—restore the checksum





Questions?



