## Lecture 22: Cache Hierarchies

- Today's topics:
- Cache access details
- Examples

Accessing the Cache

## Accessing the Cache



## The Tag Array



## Example Access Pattern



## Increasing Line Size



## Associativity



## Associativity



How many offset/index/tag bits if the cache has 64 sets,
each set has 64 bytes, 4 ways

Way-1 Way-2

## Example 1

- 32 KB 4-way set-associative data cache array with 32 byte line sizes
- How many sets?
- How many index bits, offset bits, tag bits?
- How large is the tag array?

$$
\begin{aligned}
& \text { Cache size }=\text { \#sets } \times \text { \#ways } \times \text { blocksize } \\
& \text { Index bits }=\log _{2} \text { (sets) } \\
& \text { Offset bits }=\log _{2} \text { (blocksize) } \\
& \text { Addr width }=\text { tag }+ \text { index }+ \text { offset }
\end{aligned}
$$

## Example 1

- 32 KB 4-way set-associative data cache array with 32 byte line sizes
cache size = \#sets x \#ways x block size
- How many sets? 256
- How many index bits, offset bits, tag bits?

$$
\begin{array}{ccc}
8 & 5 & 19 \\
\log _{2} \text { (sets) } & \log _{2}(\text { blksize }) & \text { addrsize-index-offset }
\end{array}
$$

- How large is the tag array?
tag array size = \#sets x \#ways x tag size

$$
=19 \mathrm{~Kb}=2.375 \mathrm{~KB}
$$

## Example 2

Show how the following addresses map to the cache and yield hits or misses. The cache is direct-mapped, has 16 sets, and a 64-byte block size.
Addresses: 8, 96, 32, 480, 976, 1040, 1096


Offset = address \% 64 (address modulo 64, extract last 6) Index = address/64 \% 16 (shift right by 6, extract last 4)
Tag $=$ address $/ 1024 \quad$ (shift address right by 10)

|  | 32-bit address |  |  |  |
| :--- | :---: | :---: | :---: | :---: |
|  | 22 bits tag | 4 bits index | 6 bits offset |  |
| 8: | 0 | 0 | 8 | M |
| 96: | 0 | 1 | 32 | M |
| 32: | 0 | 0 | 32 | H |
| 480: | 0 | 7 | 32 | M |
| 976: | 0 | 15 | 16 | M |
| 1040: | 1 | 0 | 16 | M |
| 1096: | 1 | 1 | 8 | M |

## Example 3

- A pipeline has CPI 1 if all loads/stores are L1 cache hits $40 \%$ of all instructions are loads/stores $85 \%$ of all loads/stores hit in 1-cycle L1 $50 \%$ of all (10-cycle) L2 accesses are misses
Memory access takes 100 cycles
What is the CPI?


## Example 3

- A pipeline has CPI 1 if all loads/stores are L1 cache hits $40 \%$ of all instructions are loads/stores
$85 \%$ of all loads/stores hit in 1-cycle L1
$50 \%$ of all (10-cycle) L2 accesses are misses
Memory access takes 100 cycles
What is the CPI?

Start with 1000 instructions
1000 cycles (includes all $400 \mathrm{L1}$ accesses)
+400 (ld/st) x 15\% x 10 cycles (the L2 accesses)
$+400 \times 15 \% \times 50 \% \times 100$ cycles (the mem accesses)
= 4,600 cycles
CPI $=4.6$

## Example 4



## Example 4



