



# MLIR Tutorial:

## Building a Compiler with MLIR

MLIR 4 HPC, 2019

Jacques Pienaar  
Google

Sana Damani  
Georgia Tech

Presenting the work of many people!

## Introduction

- ML != Machine Learning in MLIR
- ... but Machine Learning is one of first application domains
- And where MLIR started
- ... but not what MLIR is limited to :)

# Why a new compiler infrastructure?



## The LLVM Ecosystem: Clang Compiler



Green boxes are SSA IRs:

- Different levels of abstraction - operations and types are different
- Abstraction-specific optimization at both levels

Progressive lowering:

- Simpler lowering, reuse across other front/back ends



# Azul Falcon JVM



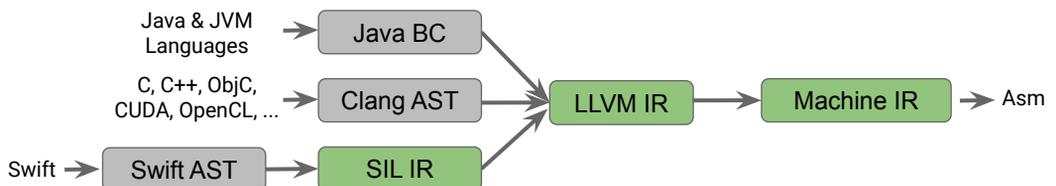
Uses LLVM IR for high level domain specific optimization:

- Encodes information in lots of ways: IR Metadata, well known functions, intrinsics, ...
- Reuses LLVM infrastructure: pass manager, passes like inliner, etc.



[“Falcon: An Optimizing Java JIT”](#) - LLVM Developer Meeting Oct'2017

# Swift Compiler



3-address SSA IR with Swift-specific operations and types:

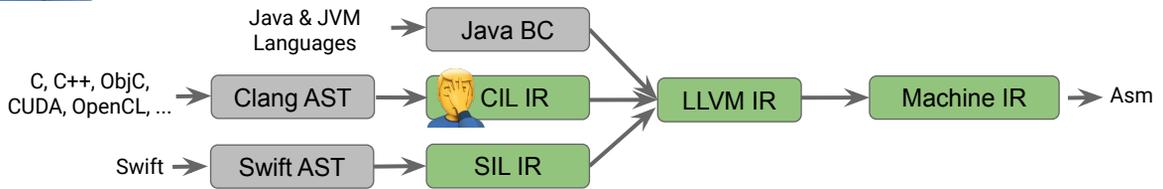
- Domain specific optimizations: generic specialization, devirt, ref count optzns, library-specific optzns, etc
- Dataflow driven type checking passes: e.g. definitive initialization, “static analysis” checks
- Progressive lowering makes each edge simpler!



[“Swift's High-Level IR”](#) - LLVM Developer Meeting Oct'2015



## A sad aside: Clang should have a CIL!



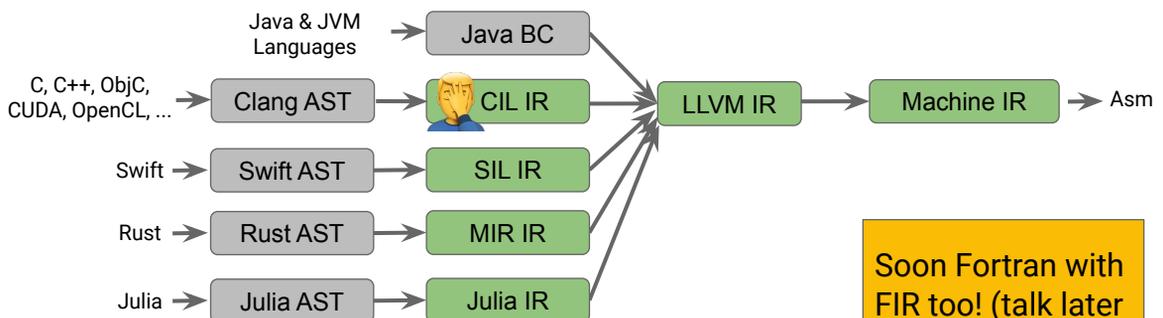
3-address SSA IR with **Clang**-specific operations and types:

- Optimizations for `std::vector`, `std::shared_ptr`, `std::string`, ...
- Better IR for Clang Static Analyzer + Tooling
- Progressive lowering for better reuse

*Anyway, back to the talk...*



## Rust and Julia have things similar to SIL



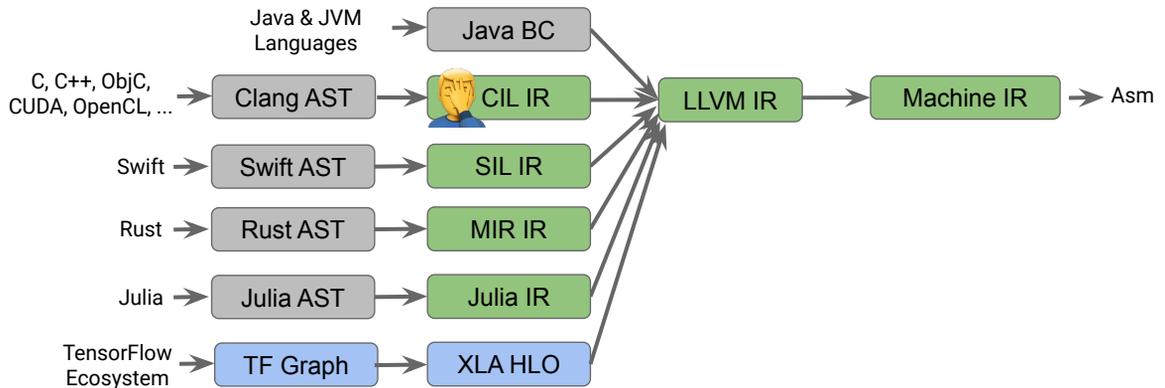
Soon Fortran with FIR too! (talk later today)

- Dataflow driven type checking - e.g. borrow checker
- Domain specific optimizations, progressive lowering



“[Introducing MIR](#)”: Rust Language Blog, “[Julia SSA-form IR](#)”: Julia docs

# TensorFlow XLA Compiler



- Domain specific optimizations, progressive lowering



“XLA Overview”: <https://tensorflow.org/xla/overview> (video overview)

## Domain Specific SSA-based IRs

Great!

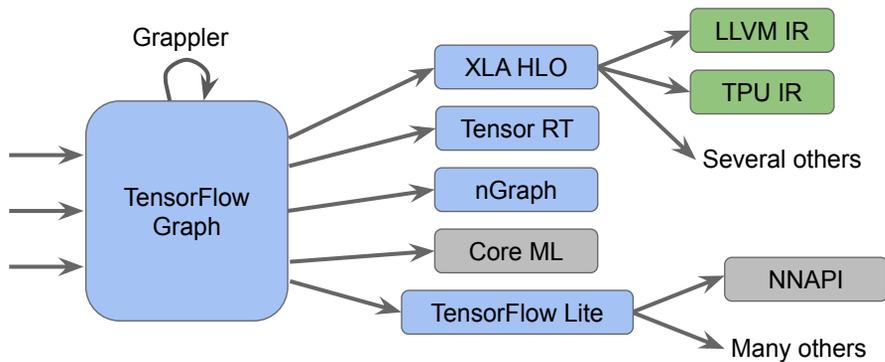
- High-level domain-specific optimizations
- Progressive lowering encourages reuse between levels
- Great location tracking enables flow-sensitive “type checking”

Not great!

- Huge expense to build this infrastructure
- Reimplementation of all the same stuff:
  - pass managers, location tracking, use-def chains, inlining, constant folding, CSE, testing tools, ...
- Innovations in one community don't benefit the others



# The TensorFlow compiler ecosystem



Many "Graph" IRs, each with challenges:

- Similar-but-different proprietary technologies: not going away anytime soon
- Fragile, poor UI when failures happen: e.g. poor/no location info, or even crashes
- Duplication of infrastructure at all levels



## Goal: Global improvements to TensorFlow infrastructure

SSA-based designs to generalize and improve ML

- Better side effect modeling and control flow
- Improve generality of the lowering passes
- Dramatically increase code reuse
- Fix location tracking and other pervasive issues for better user experience

**But HPC has similar needs so why stop there?**

No reasonable existing answers!

- ... and we refuse to copy and paste SSA-based optimizers 6 more times!

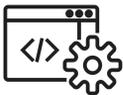


# What is MLIR?

A collection of **modular and reusable** software components that enables the **progressive lowering of operations**, to efficiently **target hardware in a common way**



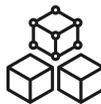
# How is MLIR different?



## State of Art Compiler Technology

MLIR is NOT just a common graph serialization format nor is there anything like it

New shared industry abstractions spanning languages ("OMP" dialect?)



## Modular & Extensible

From graph representation through optimization to code generation

Mix and match representations to fit problem space



## Not opinionated

Choose the level of representation that is right for your device

We want to enable whole new class of compiler research



# A toolkit for representing and transforming “code”

Represent and transform IR  $\rightleftharpoons$   $\updownarrow$

Represent [Multiple Levels](#) of

- tree-based IRs (ASTs),
- graph-based IRs (TF Graph, HLO),
- machine instructions (LLVM IR)

IR at the same time

While enabling

Common compiler infrastructure

- location tracking
- richer type system
- common set of conversion passes

And much more



## What about HPC?

Could talk about:

- reusing abstractions for parallelism (new parallelism constructs?),
- polyhedral code generations
- stencil abstractions

Instead:

- here to listen what are the problems domain specific abstractions during compilation could lead to much simpler/better world
- Improvements in one community benefiting others



# Introduction: a Toy Language

(e.g., enough talking, let's get to code)



## Overview

Tour of MLIR by way of implementing basic toy language

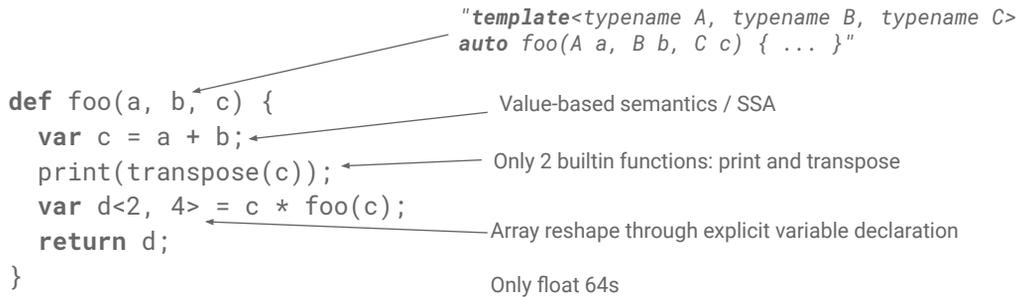
- Define a Toy language
- Represent Toy using MLIR
  - Introducing dialect, operations, ODS, verifications
- Attaching semantics to custom operations
- High-level language specific optimizations
  - Pattern rewrite framework
- Writing passes for structure rather than ops
  - Op interfaces for the win
- Lowering to lower-level dialects
  - The road to LLVM IR



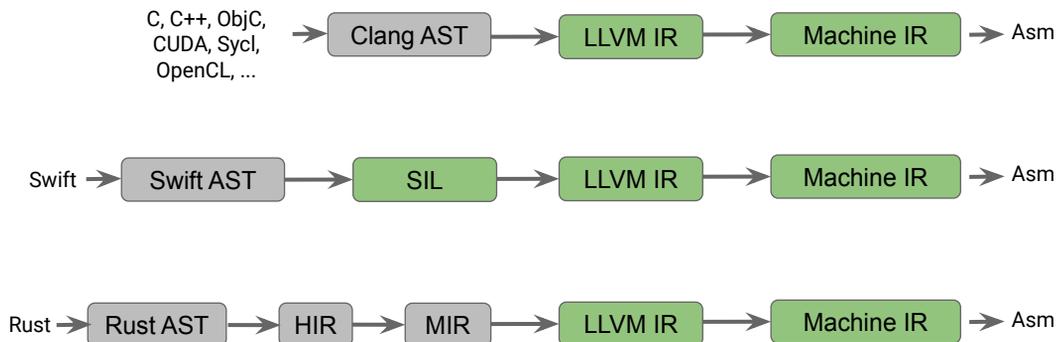
[The full tutorial on the MLIRs GitHub](#)

# Let's Build Our Toy Language

- Mix of scalar and array computations, as well as I/O
- Array shape Inference
- Generic functions
- Very limited set of operators (it's just a Toy language!):

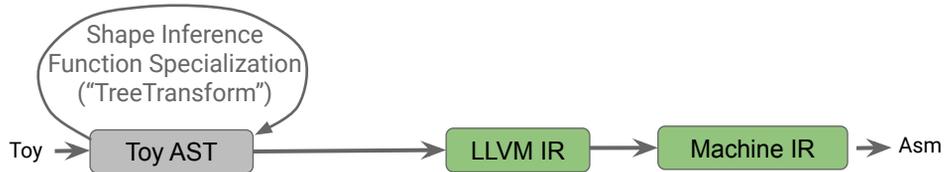


# Existing Successful Model



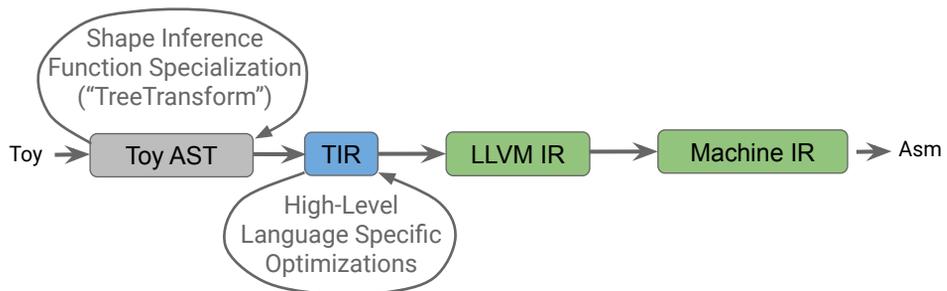
# The Toy Compiler: the "Simpler" Approach of Clang

Need to analyze and transform the AST  
-> heavy infrastructure!  
And is the AST really the most friendly  
representation we can get?



# The Toy Compiler: With Language Specific Optimizations

Need to analyze and transform the AST  
-> heavy infrastructure!  
And is the AST really the most friendly  
representation we can get?



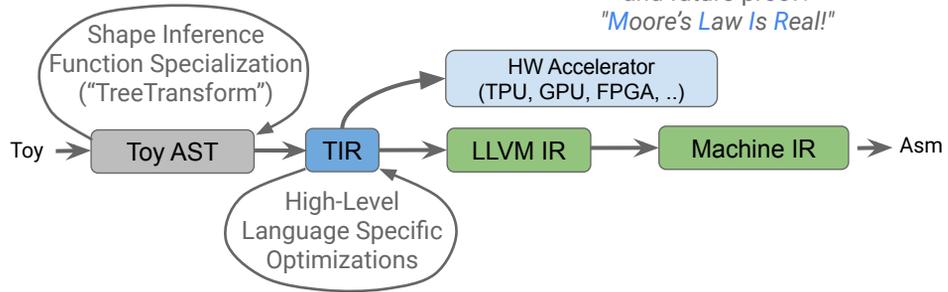
For more optimizations: a custom IR.  
Reimplement again all the LLVM infrastructure?



# Compilers in a Heterogenous World

Need to analyze and transform the AST  
-> heavy infrastructure!  
And is the AST really the most friendly  
representation we can get?

New HW: are we extensible  
and future-proof?  
*"Moore's Law Is Real!"*

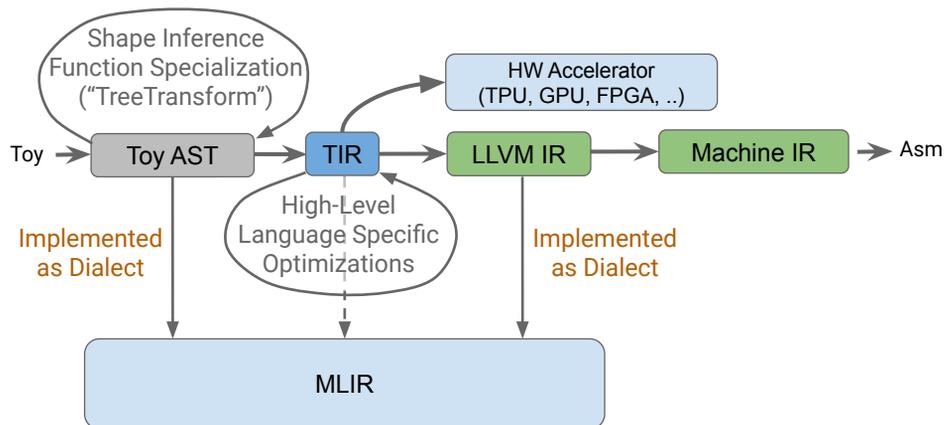


For more optimizations: a custom IR.  
Reimplement again all the LLVM infrastructure?



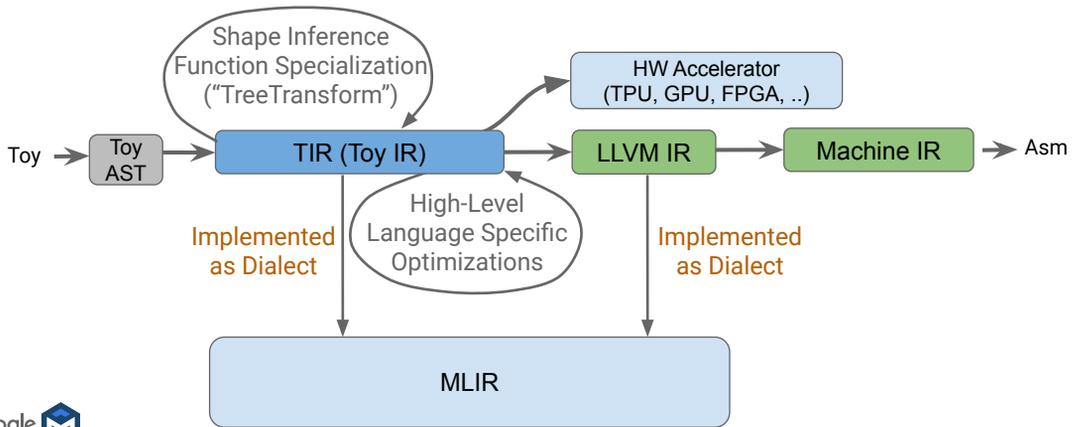
# It Is All About The Dialects!

MLIR allows every level to be  
represented as a Dialect



# Adjust Ambition to Our Budget (let's fit the talk)

Limit ourselves to a single dialect for Toy IR: still flexible enough to perform shape inference and some high-level optimizations.

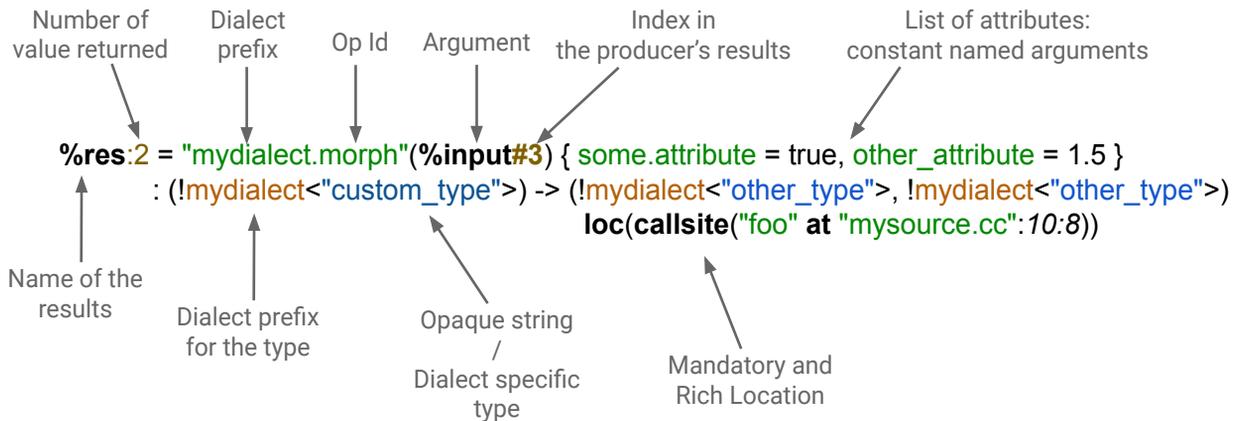


# MLIR Primer



# Operations, Not Instructions

- No predefined set of instructions
- Operations are like “opaque functions” to MLIR



# Example

```
func @some_func(%arg0: !random_dialect<"custom_type">) ->
    !another_dialect<"other_type"> {
    %result = "custom.operation"(%arg0) :
        (!random_dialect<"custom_type">) -> !another_dialect<"other_type">
    return %result : !another_dialect<"other_type">
}
```

Yes: this is a fully valid textual IR module: try round-tripping with *mlir-opt*!



## The “Catch”

```
func @main() {  
  %0 = "toy.print"() : () -> tensor<10xi1>  
}
```

Yes: this is also fully valid textual IR module!

It is not valid though! Broken on many aspects:

- the *toy.print* builtin is not a terminator,
- it should take an operand
- it shouldn't return any value

JSON of compiler IR !?!



## Dialects: Abstractions, Rules and Semantics for the IR

A MLIR dialect is a logical grouping including:

- A prefix (“namespace” reservation)
- A list of custom types, each its C++ class.
- A list of operations, each its name and C++ class implementation:
  - Verifier for operation invariants (e.g. *toy.print* must have a single operand)
  - Semantics (has-no-side-effects, constant-folding, CSE-allowed, ...)
- Possibly custom parser and assembly printer
- Passes: analysis, transformations, and dialect conversions.



## Look Ma, Something Familiar There...

Dialects are powerful enough that you can wrap LLVM IR within an MLIR Dialect

```
%13 = llvm.alloca %arg0 x !llvm<"double"> : (!llvm<"i32">) -> !llvm<"double*">
%14 = llvm.getelementptr %13[%arg0, %arg0] :
    (!llvm<"double*">, !llvm<"i32">, !llvm<"i32">) -> !llvm<"double*">
%15 = llvm.load %14 : !llvm<"double*">
llvm.store %15, %13 : !llvm<"double*">
%16 = llvm.bitcast %13 : !llvm<"double*"> to !llvm<"i64*">
%17 = llvm.call @foo(%arg0) : (!llvm<"i32">) -> !llvm<"{ i32, double, i32 }">
%18 = llvm.extractvalue %17[0] : !llvm<"{ i32, double, i32 }">
%19 = llvm.insertvalue %18, %17[2] : !llvm<"{ i32, double, i32 }">
%20 = llvm.constant(@foo : (!llvm<"i32">) -> !llvm<"{ i32, double, i32 }">) :
    !llvm<"{ i32, double, i32 } (i32)*">
%21 = llvm.call %20(%arg0) : (!llvm<"i32">) -> !llvm<"{ i32, double, i32 }">
```



## Operations: Regions are Powerful

```
%res:2 = "mydialect.morph"(%input#3) ({ Region A }, { Region B })
{ some.attribute = true, other_attribute = 1.5 } :
    (!mydialect<"custom_type">) -> (!mydialect<"other_type">, !mydialect<"other_type">)
    loc(callsite("foo" at "mysource.cc":10:8))
```

- Regions are list of basic blocks nested alongside an operation.
- Opaque to passes by default, not part of the CFG.
- Similar to a function call but can reference SSA value defined outside.
- SSA value defined inside region don't escape



# Affine Dialect: Simplified Polyhedral Form

- Multidimensional loop nests with affine loops/conditionals/memory references
- Goal: Easier to perform loop transforms (skewing, interchange etc.)
  - See presentation later today!
- Originally baked into the core
  - But not all codegen can use this form, so why not make optional?
- Expanded MLIR core so that it could become "just" a dialect
  - Regions in operations enabled moving it



## Region Example: Affine Dialect

```
func @test() {  
  affine.for %k = 0 to 10 {  
    affine.for %l = 0 to 10 {  
      affine.if (d0) : (8*d0 - 4 >= 0, -8*d0 + 7 >= 0) (%k) {  
        // Dead code, because no multiple of 8 lies between 4 and 7.  
        "foo"(%k) : (index) -> ()  
      }  
    }  
  }  
  return  
}
```

With custom parsing/printing: affine.for operations with an attached region feels like a regular for!

Extra semantics constraints in this dialect: the if condition is an affine relationship on the enclosing loop indices.



# A Toy Dialect

- Dialect & custom types defined in C++
- Dialect can define hooks for
  - type printing and printing
  - constant folding
  - ...
- Custom ops can be defined
  - Programmatically (in C++)
  - Using **Op Definition Spec** ->
  - Custom printing, parsing, folding, canonicalization, verification
  - Documentation

```
def TF_LeakyReluOp : TF_UnaryOp<"LeakyRelu",
    [NoSideEffect, SameValueType]>,
    Results<(outs TF_Tensor:$output)> {
  let arguments = (ins
    TF_FloatTensor:$value,
    DefaultValuedAttr<F32Attr, "0.2">:$alpha
  );

  let summary = "Leaky ReLU operator";
  let description = [{
    The Leaky ReLU operation takes a tensor and returns
    a new tensor element-wise as follows:
    LeakyRelu(x) = x      if x >= 0
                  = alpha*x  else
  }];

  let constantFolding = ...;
  let canonicalizer = ...;
}
```



# A (Robust) Toy Dialect

After registration, operations are now fully checked

```
$ cat test/Examples/Toy/Ch3/invalid.mlir
```

```
func @main() {
  "toy.print"() : () -> ()
}
```

```
$ build/bin/toyc-ch3 test/Examples/Toy/Ch3/invalid.mlir -emit=mlir
```

```
invalid.mlir:8:8: error: 'toy.print' op requires zero results
```

```
%0 = "toy.print"() : () -> tensor<2x3xf64>
```

^



# Toy High-Level Transformations

# Interfaces

# Motivation

- Decouple transformations from dialect and operation definitions
- Apply transformations across dialects
- Design passes to operate on attributes/structure rather than specific ops
- Prevent code duplication
- Easily extend to new dialects/ops

# Interfaces

1. Create an interface
2. Write a pass using the interface
3. Implement interface methods in participating dialects/ops

# Types of Interfaces

- Dialect Interfaces: information across operations of a dialect
  - e.g. Inlining
- Operation Interfaces: information specific to operations
  - e.g. Shape Inference



# Creating an Inliner Dialect Interface

Create a new Inliner Interface

```
class InlinerInterface
    : public DialectInterfaceCollection<DialectInlinerInterface>
{
public:
    virtual bool isLegalToInline(...) const;
    virtual void handleTerminator(...) const;
}
```



# Writing an Opaque Inliner Pass

Create a new Inliner Pass using interface collections

- Use interface collections to obtain a handle to the dialect-specific interface hook to opaquely query interface methods
- Collect all function calls and inline if legal. Also handle block terminators.

```
bool InlinerInterface::isLegalToInline(
    Operation *op, Region *dest, BlockAndValueMapping &valueMapping) const {
    auto *handler = getInterfaceFor(op);
    return handler ? handler->isLegalToInline(op, dest, valueMapping) : false;
}
```



# Inlining in Toy

Inherit DialectInlinerInterface within Toy and specialize methods

```
struct ToyInlinerInterface : public DialectInlinerInterface {
    using DialectInlinerInterface::DialectInlinerInterface;
    bool isLegalToInline(Operation *, Region *,
        BlockAndValueMapping &) const final {
        return true;
    }
    void handleTerminator(Operation *op,
        ArrayRef<Value *> valuesToRepl) const final {
        // Only "toy.return" needs to be handled here.
        auto returnOp = cast<ReturnOp>(op);

        // Replace the values directly with the return operands.
        assert(returnOp.getNumOperands() == valuesToRepl.size());
        for (const auto &it : llvm::enumerate(returnOp.getOperands()))
            valuesToRepl[it.index()->replaceAllUsesWith(it.value());
    }
};
```



## Inlining in Toy

Add the new interface to Toy Dialect

```
ToyDialect::ToyDialect(mlir::MLIRContext *ctx) : mlir::Dialect("toy", ctx) {  
    ...  
    addInterfaces<ToyInlinerInterface>();  
}
```

Add Inliner Pass to Toy's pass manager

```
mlir::LogicalResult optimize(mlir::ModuleOp module) {  
    mlir::PassManager pm(module.getContext());  
    ...  
    pm.addPass(mlir::createInlinerPass());  
    ...  
}
```



## Operation Interfaces: Shape Inference

- We'll use Shape Inference as an example application of operation interfaces
- We define the following rules for shape inference in Toy
  - $A = B + C$  //  $A.shape = B.shape = C.shape$
  - $A = B * C$  //  $A.shape = B.rows, C.cols$
  - $A = \text{transpose}(B)$  //  $A.shape = B.cols, B.rows$



## Creating a Shape Inference Interface

Create a ShapeInference OpInterface:

```
def ShapeInferenceOpInterface : OpInterface<"ShapeInference"> {  
  let methods = [  
    InterfaceMethod<"Infer output shape for the current operation.",  
                  "void", "inferShapes", (ins), []>  
  ];  
}
```



## Writing an Opaque Shape Inference Pass

Thanks to operation interfaces, we can write an opaque ShapeInference Pass:

```
while (!opWorklist.empty()) {  
  ...  
  op = ...  
  // Use inferShape if `op` implements the Shape Inference interface  
  if (auto shapeOp = dyn_cast<ShapeInference>(op)) {  
    shapeOp.inferShapes();  
  }  
  ...  
}
```



## Shape Inference in Toy

Specialize interface methods in Toy's op definitions:

```
def AddOp : Toy_Op<"add", [NoSideEffect]> {  
  ...  
  void inferShapes() {  
    getResult()->setType(getOperand(0)->getType());  
    return;  
  }  
}
```

And then add ShapeInference pass to Toy's pass manager.



## Pattern-Match and Rewrite



# Language Specific Optimizations

```
def no_op(b) {  
    return transpose(transpose(b));  
}
```

Clang can't optimize away these loops:

```
#define N 100  
#define M 100  
  
void sink(void *);  
void double_transpose(int A[N][M]) {  
    int B[M][N];  
    for(int i = 0; i < N; ++i) {  
        for(int j = 0; j < M; ++j) {  
            B[j][i] = A[i][j];  
        }  
    }  
    for(int i = 0; i < N; ++i) {  
        for(int j = 0; j < M; ++j) {  
            A[i][j] = B[j][i];  
        }  
    }  
    sink(A);  
}
```



# Generic DAG Rewriter

- Graph-to-graph rewrites
- Decouple pattern definition and transformation
- Greedy worklist combiner



# Pattern Match and Rewrite

`transpose(transpose(x)) => x`

Two ways:

- C++ style using RewritePattern
- Table-driven using DRR



## C++ Style using RewritePattern

`transpose(transpose(x)) => x`

Override matchAndRewrite(op):

```
input = op.getOperand();
if (input->getDefiningOp() == TransposeOp)
    x = op->getOperand();
rewriter.replaceOp(op, {x});
```

Register Pattern with Canonicalization Framework

```
void TransposeOp::getCanonicalizationPatterns(...) {
    results.insert<SimplifyRedundantTranspose>(context);
}
```



## Declarative, rule-based pattern-match and rewrite

```
transpose(transpose(x)) => x

// Transpose(Transpose(x)) = x
def TransposeTransposeOptPattern : Pat<
  (TransposeOp(TransposeOp $arg)),
  (replaceWithValue $arg)>;

class Pattern<
  dag sourcePattern,
  list<dag> resultPatterns,
  list<dag> additionalConstraints = [],
  dag benefitsAdded = (addBenefit 0)
>;
```



See another example in the repo with Reshape(Reshape(x)):  
<https://github.com/tensorflow/mlir/blob/master/examples/toy/Ch4/mlir/ToyCombine.td#L43-L44>

## Declarative, rule-based pattern-match and rewrite

Conditional pattern match:

```
Reshape(x) = x, if input and output shapes are identical
```

Adding Constraints:

```
def TypesAreIdentical : Constraint<CPred<"$0->getType() == $1->getType()">>;
```

Transformation:

```
def ReshapeOptPattern : Pat<(ReshapeOp:$res $arg), (replaceWithValue $arg), \
  [(TypesAreIdentical $res, $arg)]>;
```



<https://github.com/tensorflow/mlir/blob/master/examples/toy/Ch4/mlir/ToyCombine.td#L67-L71>

## Declarative, rule-based pattern-match and rewrite

Complex Transformation:

`Reshape(Constant(x)) = x'`, where `x'` is `Reshape(x)`

Native Code Call:

```
def ReshapeConstant :  
NativeCodeCall<"$0.reshape(($1->getType()).cast<ShapedType>())">;
```

Transformation:

```
def ConstantReshapeOptPattern : Pat<(ReshapeOp:$res (ConstantOp $arg)), \  
    (ConstantOp (ReshapeConstant $arg, $res))>;
```



## Dialect Lowering

All the way to LLVM!



# Lowering

- Goal: Translating source dialect into one or more target dialects
- Full or Partial
- Procedure:
  - Provide target dialects
  - Operation Conversion
  - Type Conversion



# DialectConversion framework

Goal: Transform illegal operations to legal ones

Components of the Framework:

- Conversion Target: Which dialects/ops are legal after lowering?
- Rewrite Patterns: Convert illegal ops to legal ops
- Type Converter: Convert types



## Conversion Targets

- Legal Dialects (target dialects)  
`target.addLegalDialect<mlir::AffineOpsDialect, mlir::StandardOpsDialect>();`
- Illegal Dialects (fail if not converted)  
`target.addIllegalDialect<ToyDialect>();`
- Legal and Illegal Ops  
`target.addLegalOp<PrintOp>(); // preserve toy.print`  
`target.addIllegalOp<BranchOp>(); // must convert`
- Dynamically Legal Ops/Dialects (legality constraints such as operand type)  
`target.addDynamicallyLegalOp<ReturnOp>();`



## Operation Conversion using ConversionPattern Rewriter

Convert illegal ops into legal ops using a pattern match and rewrite

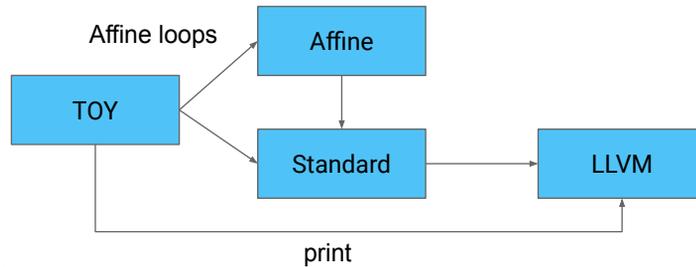
Transitive conversion: [bar.add -> baz.add, baz.add -> foo.add]

ConversionPattern rewriter vs PatternMatch rewriter:

- Additional operands parameter to specify newly rewritten values
- No N->1 or N->M conversions
- Roll-back on failure



# Conceptually: Graph Of Lowering



A->B->C lowering

Lowering Graph:

- Nodes: Dialects/Ops
- Edges: Conversion
- Open Problem: Finding optimal\* route



## Affine to LLVM

- Now let's generate some executable code
- Same conversion as before but with Type conversion
- Full Lowering

Populate the Lowering Graph:

```
mllir::OwningRewritePatternList patterns;  
mllir::populateAffineToStdConversionPatterns(patterns, &getContext());  
mllir::populateLoopToStdConversionPatterns(patterns, &getContext());  
mllir::populateStdToLLVMConversionPatterns(typeConverter, patterns);
```



# MLIR LLVM dialect to LLVM IR

Mapping from LLVM Dialect ops to LLVM IR:

```
auto llvmModule = mlir::translateModuleToLLVMIR(module);
```

LLVM Dialect:

```
%223 = llvm.mlir.constant(2 : index) : !llvm.i64  
%224 = llvm.mul %214, %223 : !llvm.i64
```

LLVM IR:

```
%104 = mul i64 %96, 2
```

# Conclusion

# MLIR : Reusable Compiler Abstraction Toolbox

IR design involves multiple tradeoffs

- Iterative process, constant learning experience

MLIR allows mixing levels of abstraction with non-obvious compounding benefits

- Dialect-to-dialect lowering is easy
- Ops from different dialects can mix in same IR
  - Lowering from “A” to “D” may skip “B” and “C”
- Avoid lowering too early and losing information
  - Help define hard analyses away

} No forced IR impedance mismatch  
} Fresh look at problems



## Not shown today

- Heterogeneous compilation
- MLIR also includes GPU dialect to target
  - CUDA,
  - RocM, and
  - SPIR-V/Vulkan
- New converters to
  - TFLite
  - XLA



## Recap

MLIR is a great infrastructure for higher-level compilation

- Gradual and partial lowerings to mixed dialects
  - All the way to LLVMIR and execution
- Reduce impedance mismatch at each level

MLIR provides all the infrastructure to build dialects and transformations

- **At each level it is the same infrastructure**

Demonstrated this on a Toy language

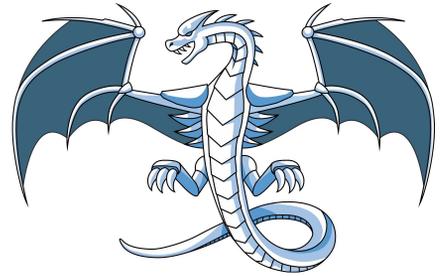
- Tutorial available on github

## Getting Involved

# MLIR is Open Source!

Visit us at [github.com/tensorflow/mlir](https://github.com/tensorflow/mlir):

- Code, documentation, examples
  - Core moving to LLVM repo soon
- Developer mailing list at: [mlir@tensorflow.org](mailto:mlir@tensorflow.org)
- Open design meetings every Thursday
- Contributions welcome!



# MLIR

Thank you to the team!

Questions?

We are hiring!  
[mlir-hiring@google.com](mailto:mlir-hiring@google.com)