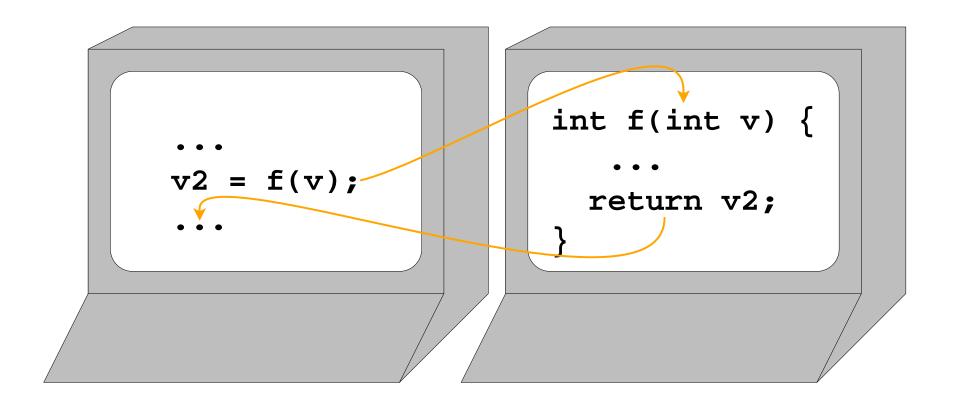
Today's topics:

- Remote procedure calls (RPC)
- Transactions

Procedure Call

Remote Procedure Call



RPC vs. DSM

DSM:

- Illusion of shared memory
- No illusion of shared control

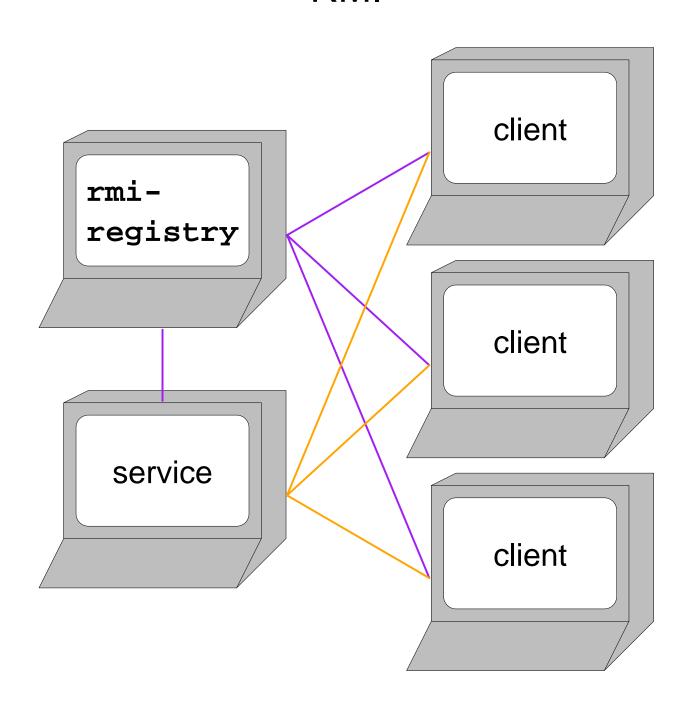
RPC:

- No illusion of shared memory
- Illusion of shared control

Some RPC Systems

- Java RMI
- CORBA
- rpcgen

RMI

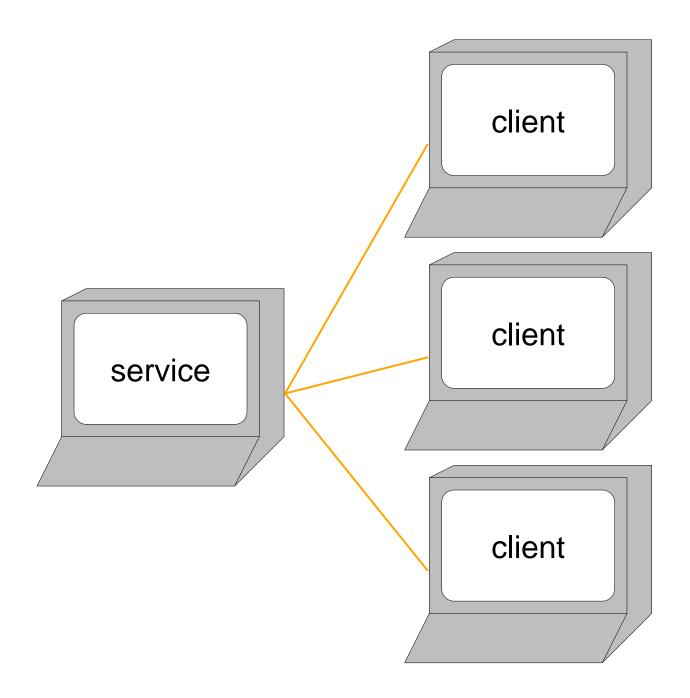


RMI

Each connected client gets its own thread in the service

see rmi

rpcgen



rpcgen

Service is single-threaded

see rpcgen, field

Multiple Clients in the Field Game

Could add locks:

- Add per-client thread in the service
- Use the usual locks on the service side

Instead, we'll use transactions

- Keeps service single-threaded
- More directly implements what we mean

The goal of a lock is to allow a single thread to modify state at a time, preventing interleaving

A *transaction* brackets a set of operations to say:

- Make this appear instantaneous with respect to other threads
- If that's not possible, do nothing and let me try again

```
start_trans();
                   start_trans();
a[2] = a[3] + 5;
                    a[1] = a[2] + 7;
end_trans();
                     ...aborts...
                     start_trans();
                    a[1] = a[2] + 7;
                    end_trans();
```

Simplifed variant:

see field_trans

More general variant (includes early abort):

see field_trans2

Benefits:

- May allow more concurrency, since transaction system figures out granularity
- Easier to reason about

Drawbacks:

Extra runtime overhead

Design issues:

- Log writes and check on read, or perform writes and log for undo
- Is a transaction atomic with respect to non-transactions?
- Nested transactions
- Exception handling