# File Systems

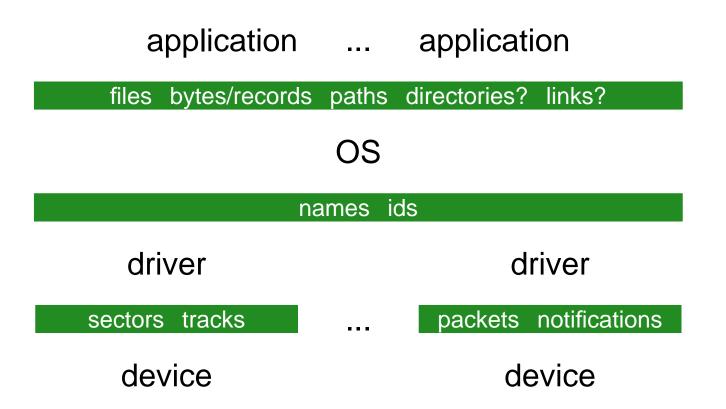
*File system* = most common abstraction for *persistence* 

Also provides

- Large-data storage with random access
- Data organization
- Mobility (e.g., CD ROM, NFS)
- Sharing & protection
- Communication

"File system" sometimes refers to the abstraction and sometimes refers to a particular disk format. We mean the former.

### File system Layers



### Files

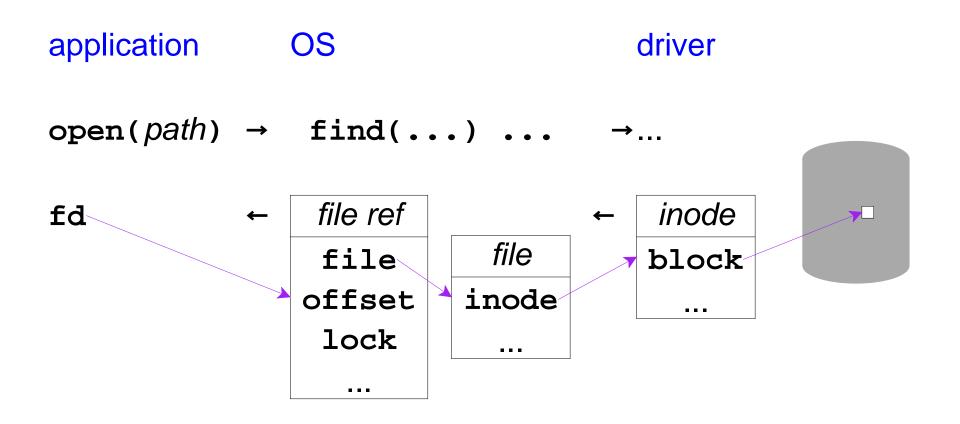
Typically, a *file* is

- A sequence of bytes
- Metadata, including modification time, permissions, and type

Typically, a file is accessed through a *path* 

- Access results in a *file descriptor* or *file handle*
- Descriptor or handle sticks with a file, while the path can change

# Opening a File



# **Common File System Operations**

Data

- Create() Open() Read()
- Delete() Close() Write()
  - Seek()

Naming

- Rename() HardLink()
  - SoftLink()

Metadata

• GetAttribute() • SetAttribute()

#### Create()

Unix:

int open(const char \*path, int oflag, mode\_t mode);
with O\_CREAT

also opens

int mkdir(const char \*path, mode\_t mode);

Windows:

HANDLE CreateFile(LPCTSTR lpFileName, ....); with CREATE\_ALWAYS

also opens

HANDLE CreateDirectory(LPCTSTR lpPathName, ....);

See create.c

### Delete()

Unix:

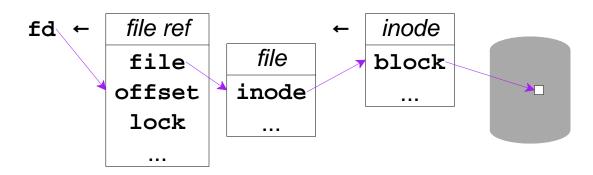
int unlink(const char \*path);

int rmdir(const char \*path);

Windows:

BOOL DeleteFile(LPCTSTR lpPathName);

BOOL RemoveDirectory(LPCTSTR lpPathName);



Removes the path mapping, but doesn't actually delete until all references are closed (see create.c)

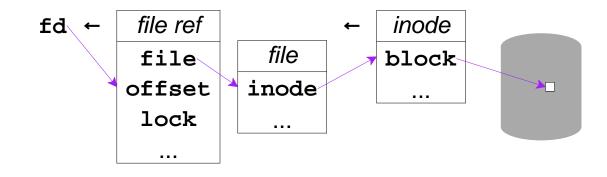
### Open()

Unix:

int open(const char \*path, int oflag);

Windows:

HANDLE CreateFile(LPCTSTR lpFileName, ....);



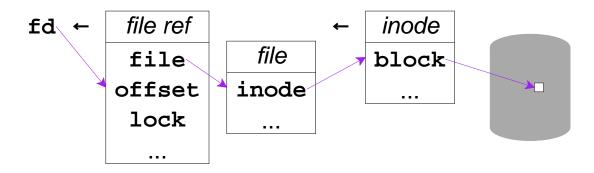
### Close()

Unix:

int close(int filedes);

Windows:

BOOL CloseHandle(HANDLE hFile);



Last copy of decriptor/handle  $\Rightarrow$  free decriptor/handle Last decriptor/handle  $\Rightarrow$  close file

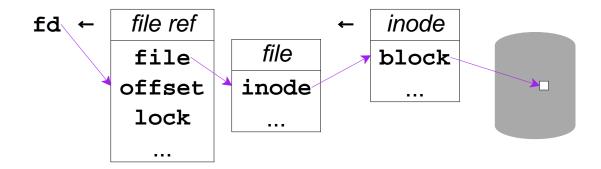
#### Read()

Unix:

ssize\_t read(int fildes, void \*buf, size\_t nbyte);

Windows:

BOOL ReadFile(HANDLE hFile, LPVOID lpBuf, ....);



Updates descriptor/handle offset

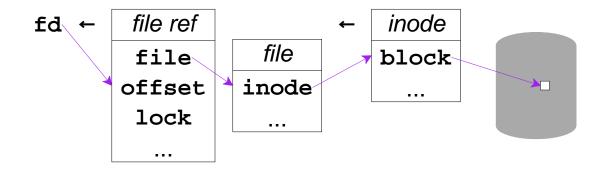
#### Write()

Unix:

ssize\_t write(int fildes, void \*buf, size\_t nbyte);

Windows:

BOOL WriteFile(HANDLE hFile, LPVOID lpBuf, ....);



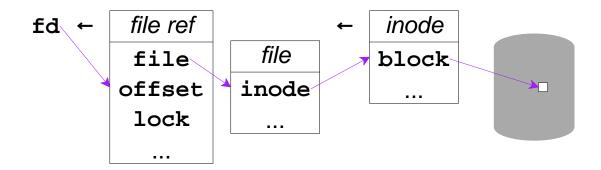
Updates descriptor/handle offset

Unix:

off\_t lseek(int fildes, off\_t offset, int whence);

Windows:

DWORD SetFilePointer(HANDLE hFile, LONG lOff, ...);



Updates descriptor/handle offset

#### Rename()

Unix:

int rename(const char \*old, const char \*new);

Windows:

BOOL MoveFile(LPCTSTR lpOld, LPCTSTR lpNew);



No effect on open descriptors/handles

Atomic update when on the same device

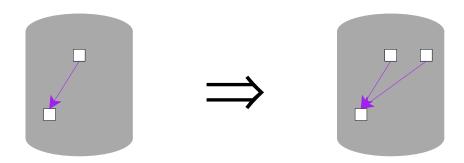
#### HardLink()

Unix:

int link(const char \*old, const char \*new);

Windows:

BOOL CreateHardLink(LPCTSTR lpNew, LPCTSTR lpOld, ...);



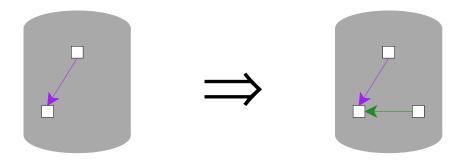
No effect on open descriptors/handles

See share2.c

#### SoftLink()

Unix:

int symlink(const char \*path, const char \*new);



No effect on open descriptors/handles

#### GetAttribute()

Unix:

int fstat(int filedes, struct stat \*buf);

Windows:

BOOL GetFileInformationByHandle(HANDLE hFile, ...);

File type, size, maybe permissions

#### SetAttribute()

Unix:

- int fchmod(int fildes, mode\_t mode);
- int futimes(int fildes, struct timeval times[2]);

Windows:

BOOL SetFileInformationByHandle(HANDLE hFile, ...);

File type, size, maybe permissions

## Unix Paths

 A *path* is a sequence of byte-strings elements, where / is disallowed in an element

usr local bin pdf2ps

- A path is normally written as a single byte string using / as a separator
  - Path starts with / ⇒ absolute
     /usr/local/bin/pdf2ps
  - Path does not start with / ⇒ relative
     bin/pdf2ps
- Each process has a *working directory* that prefixes relative paths

### **Unix Paths**

#### • A device is *mounted* at one or more path prefixes

\$ /usr/bin/mount /dev/sda2 on / type ext3 (rw) proc on /proc type proc (rw) sysfs on /sys type sysfs (rw) devpts on /dev/pts type devpts (rw,gid=5,mode=620) /dev/sda5 on /usr/vice type ext3 (rw) /dev/sda1 on /boot type ext3 (rw) zfs:/server/home2/mflatt on /home/mflatt type nfs ...

# **Unix Paths**

- OS communicates to driver in terms of IDs, known as inodes and immediate names
  - $^{\circ}$  A file is a kind of inode
  - A directory is a kind of inode
  - A hard link is when a directory points to a file's inode
  - A soft link is an inode that contains another path, automatically followed (usually) by the OS
- Case sensivity is managed by the driver
  - ext3 (Linux) is case-sensitive
  - HFS+ (Mac OS) is case-insensitive by default

# Windows Paths

- A *path* combines a drive with a UTF-16 code unit sequence
- A path is normally written as a single string using a letter name for a drive and \ as a separator, in which case <, >, :, ", /, \, and | are disallowed in an element C:\Program Files\PLT\DrScheme.exe
- A drive can also be \\machine\volume
- Except that special files names like **aux** refer to devices, independent of the drive, path, or extension
- At some layers of the Windows API, various automatic transformations are applied, such as converting / to \ and dropping trailing spaces

C:/Program Files/PLT\DrScheme.exe

### Windows Paths

- Path starts with drive and \ ⇒ absolute
   C:\Program Files\PLT\DrScheme.exe
- Path does not start with drive or  $\ \Rightarrow$  relative PLT\DrScheme.exe
- Path starts with drive but not  $\ \Rightarrow$  drive-relative C:PLT\DrScheme.exe
- Path starts with \ ⇒ drive-absolute
   \Program Files\PLT\DrScheme.exe
- Each process has a working drive and each drive per process has a working directory

### Windows Paths

- OS communicates to driver in terms of paths

   Use the \\?\ prefix to specify driver path directly
   \\?\c:\wE|Rd\<path>
- Case sensivity is managed by the OS

### Paths

- Generally cannot get a cannonical path for a file
  - The path can change
  - May have multiple mount points
  - May have multiple links
- File descriptor/handle provides cannonical references
  - e.g, get inode
  - Only works for open files

What if cooperating processes want to modify a file, and only one process should modify the file at a time?

• Advisory locks — provided by the OS to let cooperating programs declare exclusive access

• Unix, typically

 Mandatory locks — provided by the OS to let programs (cooperative or not) gain exclusive access

 $\circ$  Windows

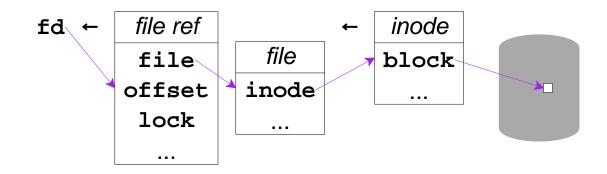
### Lock() and Unlock()

Unix:

int flock(int fd, int operation);

Windows:

HANDLE CreateFile(LPCTSTR lpFileName, ..., DWORD dwShareMode, ...);



See locks.c, locks2.c, locks2.c

# Permissions

- Access-control list (ACL) determines for each file which userid can perform which of a handful of operations
  - Typical operations: read, write, execute, append, delete, list
- Unix-style simplified mapping:
  - owner vs. group vs. everyone
  - read, write, execute