

The E1 Distributed Operating System Project

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Goals



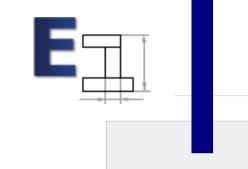
- Provide efficient access to the resources of computer network.
- Implement a convenient programming model, isolating software developers from the intrinsic complexity of asynchronous distributed environment.



Distributed Objects



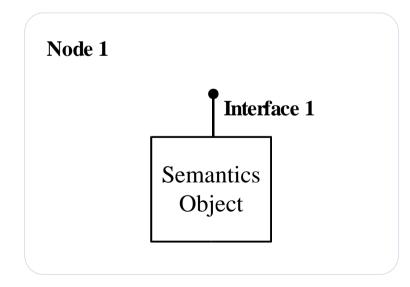
Distributed Object = U_{nodes} (Semantics Object + Replication Object)

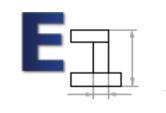




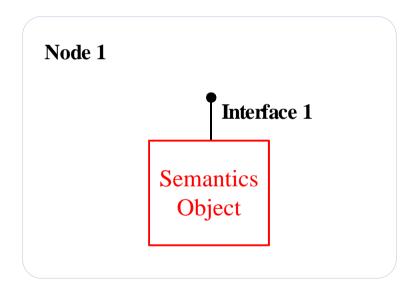
When distributed object is used only in one node:

Distributed Object = Semantics Object





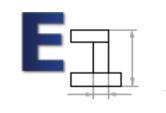
Semantics Object



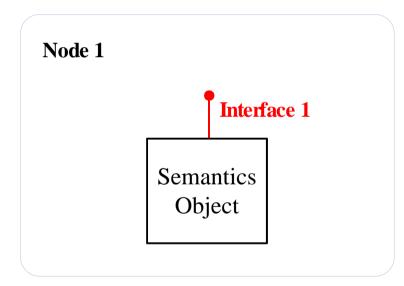
Semantics object is much like a C++ object, it:

- Stores object state
- Implements local object functionality

but ...



Semantics Object



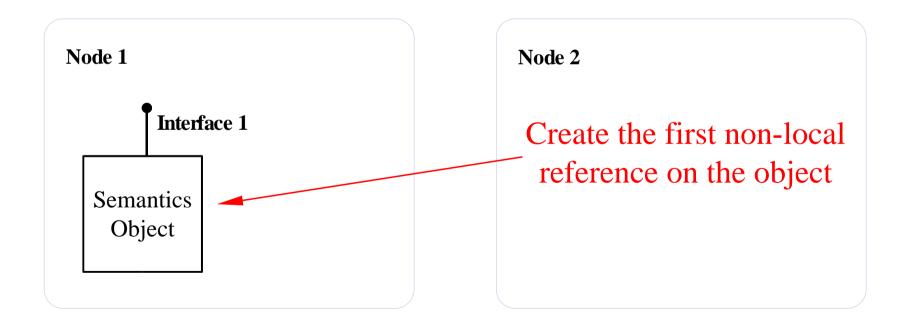
Semantics object is much like a C++ object, it:

- Stores object data
- Implements local object functionality

but it's accessible only via object interfaces

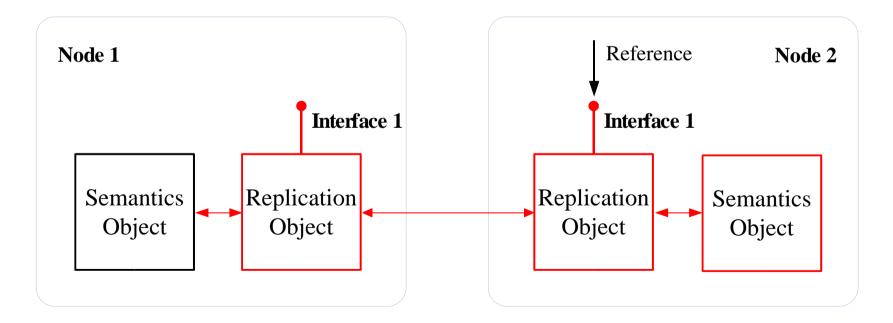


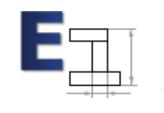
Creation of the first non-local reference initiates creation of replication objects



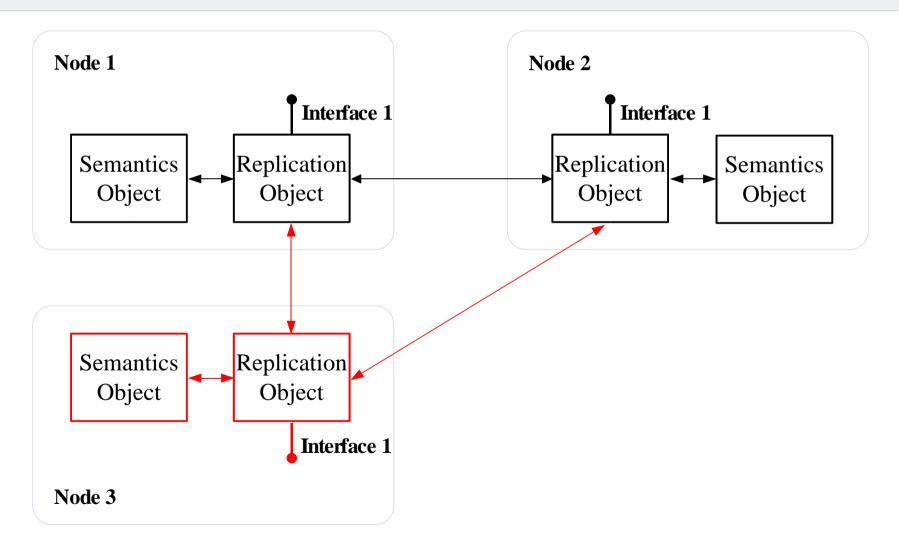


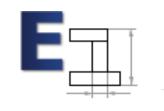
A pair of replication and semantics objects is created in each node



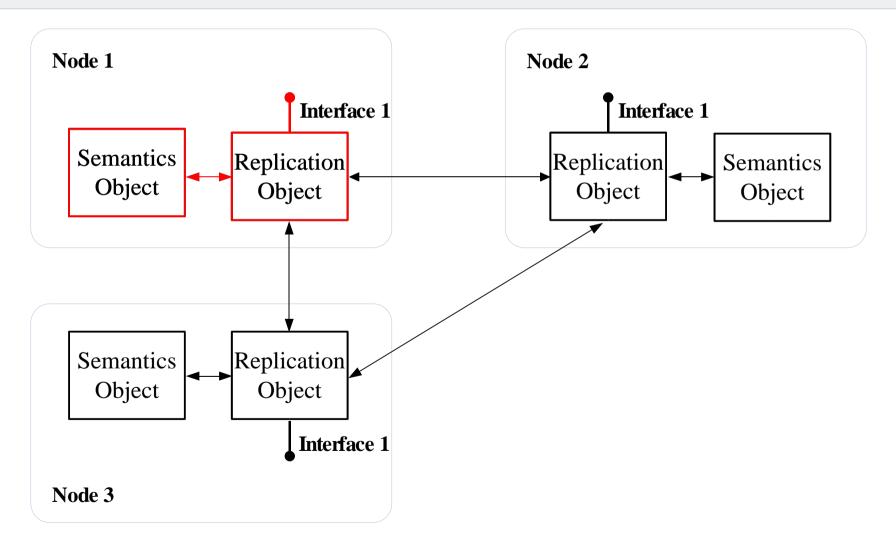


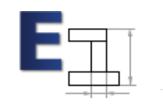
More Nodes...



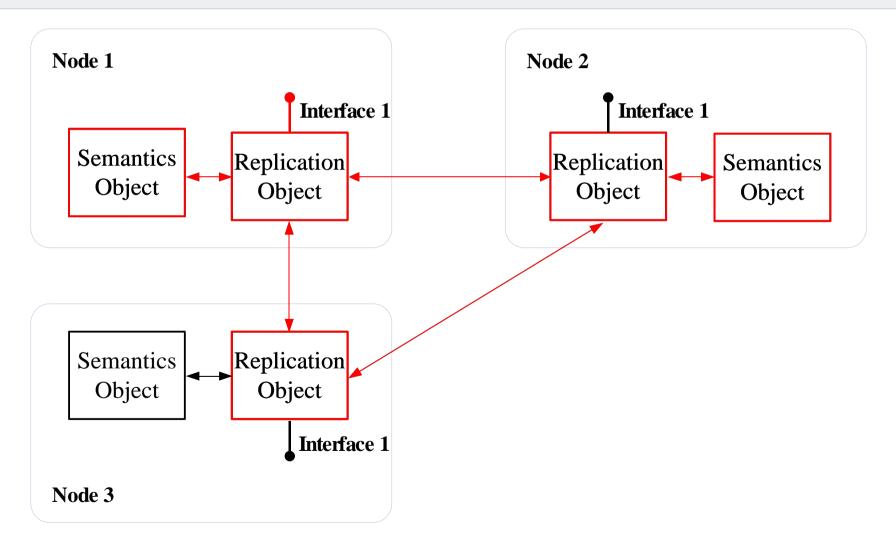


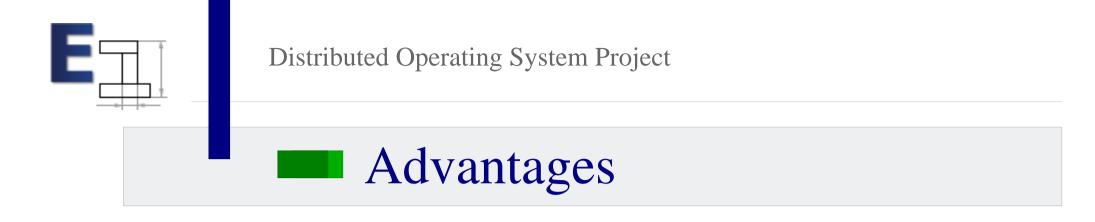
Object Invocations



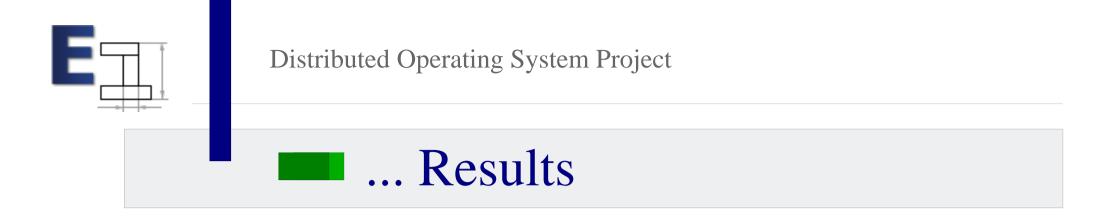


Object Invocations





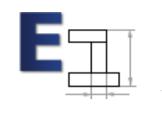
- Effective separation of object semantics and replication strategy
- No imposed restrictions on the replication strategies.

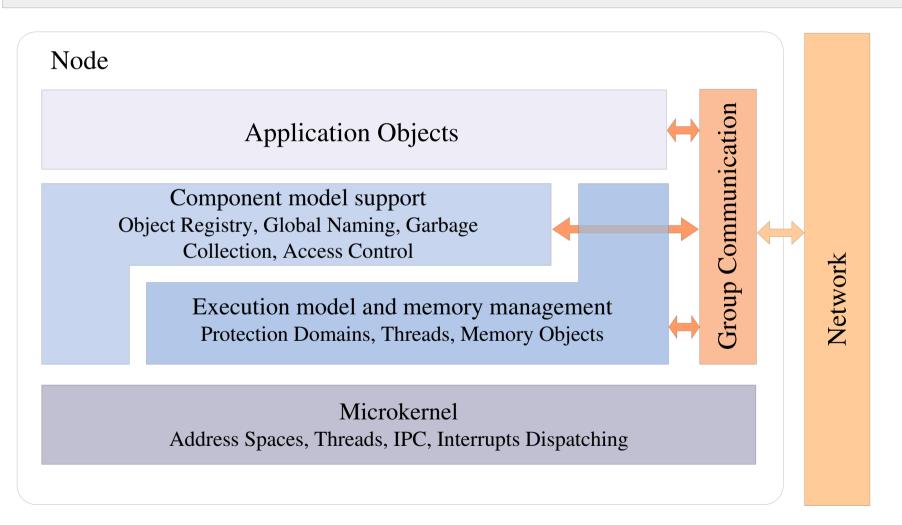


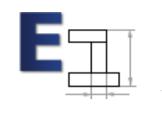
- Object functionality can be implemented separately from its replication strategy
- It's possible to select most efficient replication strategy for each object.

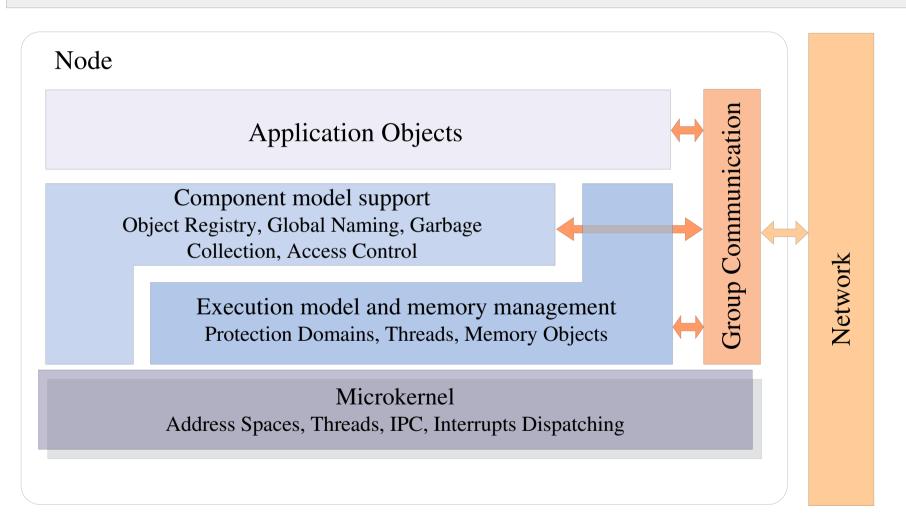


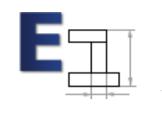
System Architecture

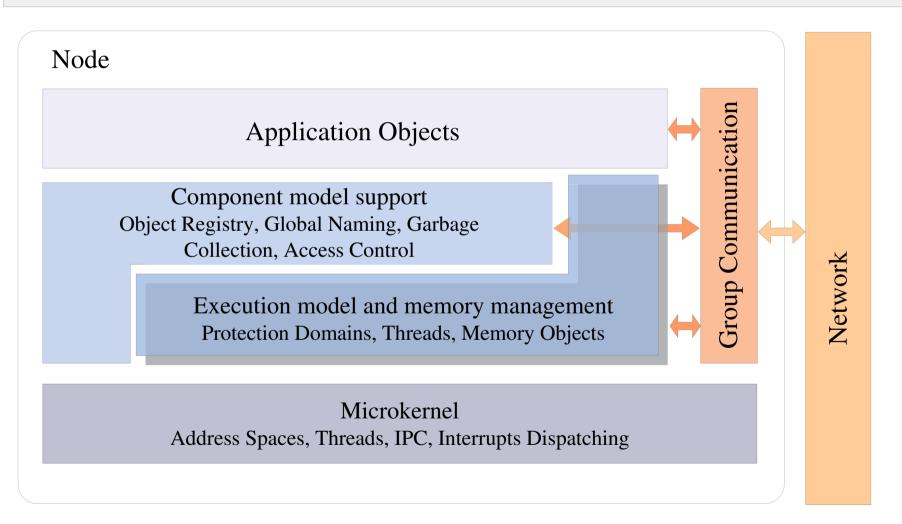


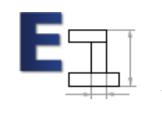


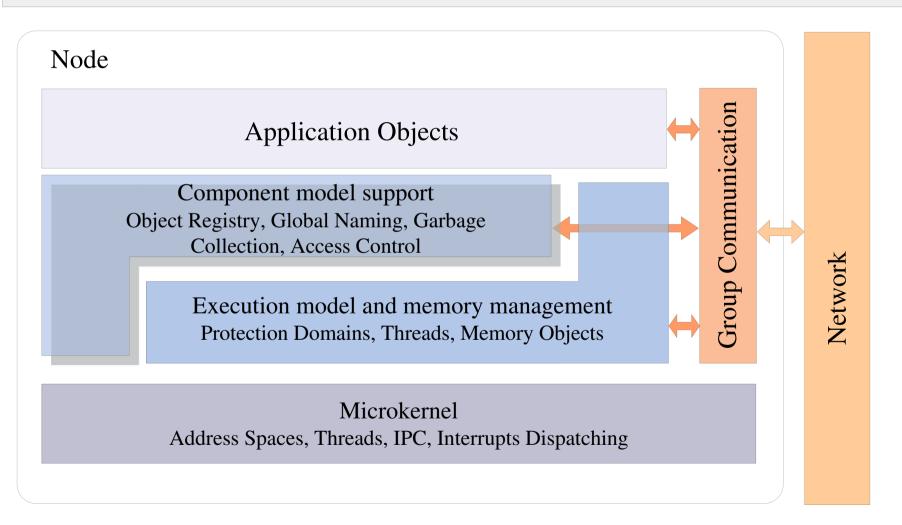






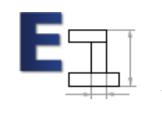


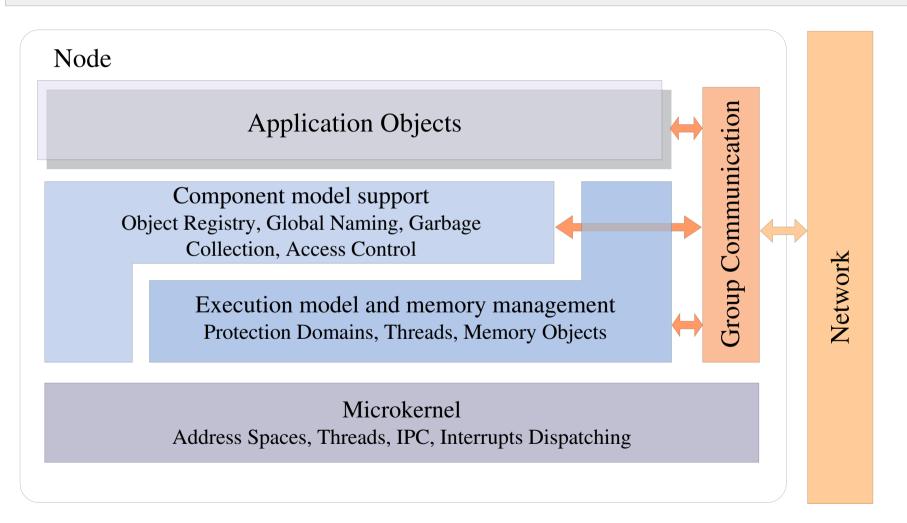


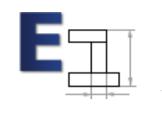


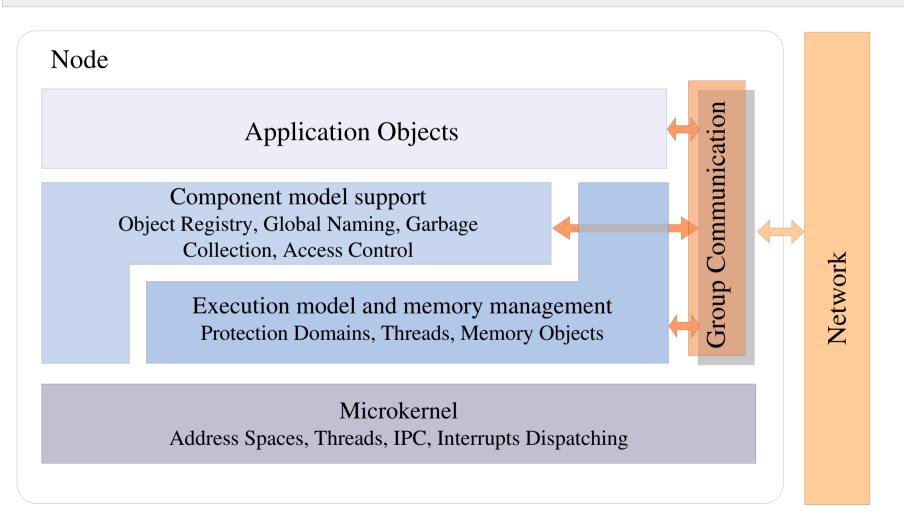


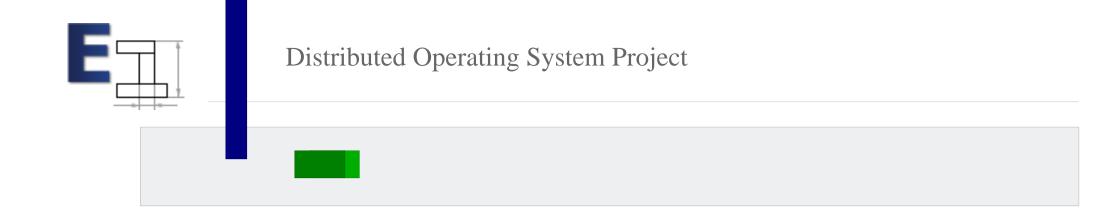
- Lifecycle control
 - Object Registry
 - Distributed Garbage Collection System
- Global Naming Service
- Access Control Server











Replication



Allow developers to define distributed behaviour of applications without implementing distributed algorithms.

• This is in contrast to DSM and RMI -based operating systems, which try to make distribution completely transparent by sacrificing efficiency.

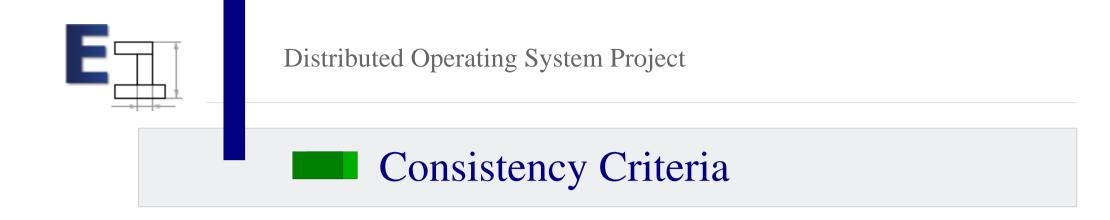


In E1 distributed behaviour of an object is defined by:

- 1. Selecting replication strategy for the object (possibly, at run time)
- 2. Adjusting replication strategy parameters:
 - consistency criteria;
 - required level of redundancy;
 - object topology (replica placement) ...



- timing properties;
- failure detection strategies;
- failure-handling policies;
- handling of network fragmentations;
- network protocol selection;
- etc.



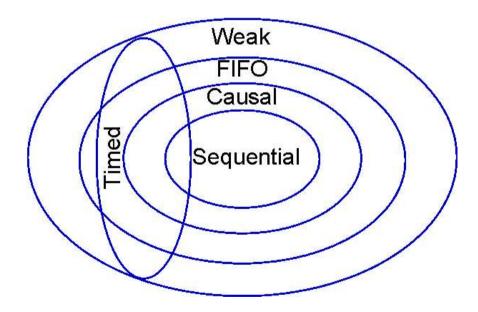
Strict consistency

• Sequential

Relaxed consistency

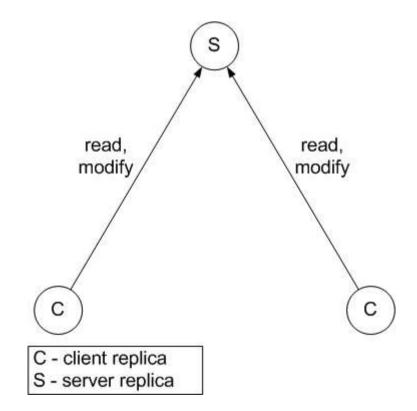
- Causal
- FIFO
- Weak

Timed consistency





Client/Server Replication

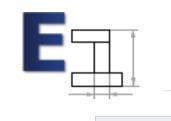


Pros:

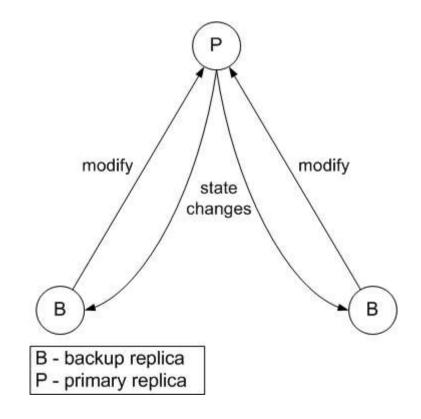
+ universal

Cons:

- inefficient
- unreliable



Passive Replication



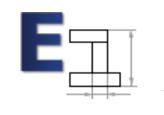
Pros:

+ reliability

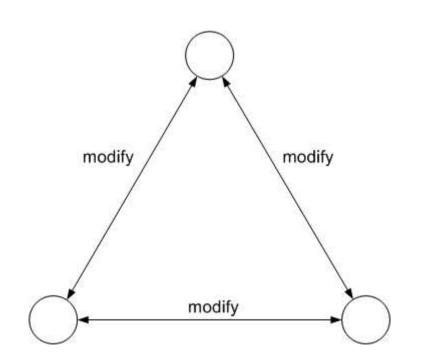
+ reads are local

Cons:

- updates are not local



Active Replication



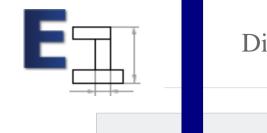
Pros:

+ reliability

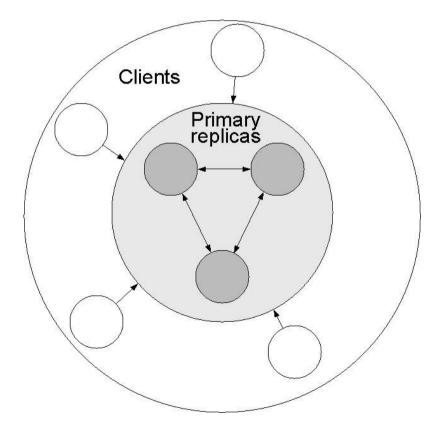
+ both reads and updates are local when allowed by consistency criteria

Cons:

- requires deterministic behaviour
- recursive invocations are difficult



Example of a Real Strategy



A small set of active primary replicas with many stateless client replicas connected to them.

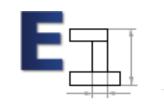


The Replication Strategies Framework

Replication Strategies Library

Virtually Synchronous Group Communication

Network Protocol Stack



Network Protocol stack

Replication Strategies Lib

Virtually Synchronous GC

Network Protocol Stack

- Network protocol layer provides at least unreliable unicast primitive.
- However, more advanced primitives, e.g. unreliable multicast or reliable unicast can also be available.



Virtually Synchronous GC

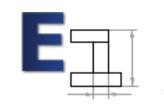
Replication Strategies Lib

Virtually Synchronous GC

Network Protocol Stack

<u>Implements two types of</u> <u>services:</u>

- Group membership service.
- Reliable unicast & multicast message delivery services with various ordering guarantees.



Group Membership Service

Replication Strategies Lib

Virtually Synchronous GC

Network Protocol Stack

Detectscrashedandrecoveredobjectreplicasusingunreliablefailuredetectoranddeliversconsistentviews of the groupof replicas to all its members.



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Message Delivery Service

Replication Strategies Lib

Virtually Synchronous GC

Network Protocol Stack

Message ordering properties:

- FIFO multicast
- Causal multicast
- Totally ordered multicast



Distributed Operating System Project

Replication Strategies Library

Replication Strategies Lib

Virtually Synchronous GC

Network Protocol Stack



• Completeness

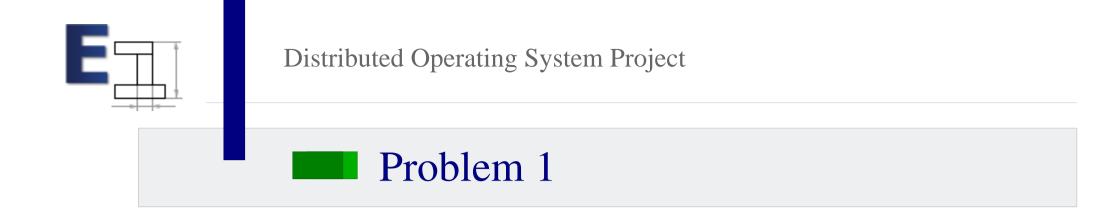
For virtually any object a replication strategy providing "good" performance can be found in the library.

• Customizability

The developer can further fine-tune application performance by adjusting the replication strategy parameters.

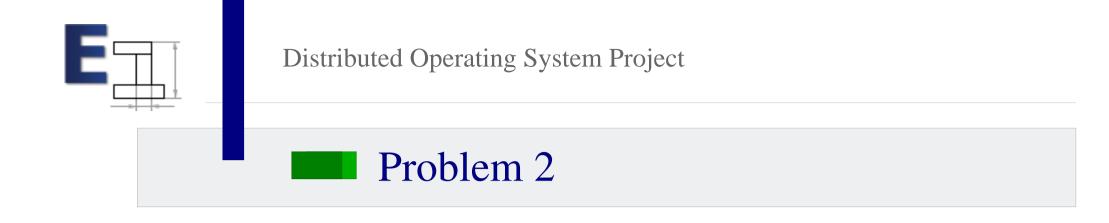
• Extensibility

New replication strategies can be easily developed by reusing existing components and design patterns.



A classical software engineering problem: given a number of related algorithms, construct a framework for unified development, evaluation, utilization and modification of these and similar algorithms.

For example, a similar problem has been successfully solved in the domain of group communication systems (Horus, Transis).



Under what conditions can two replicated objects (with different replication strategies) interact without breaking the consistency of each other?

• Completely ignored in previous research.



- Any non-trivial replication strategy involves operations requiring serialization/deserialization of object state:
 - Creation of a new replica;
 - Migration of existing replica to a new node;
 - State synchronization (passive replication).
- Objects are required to provide a serialization interface.

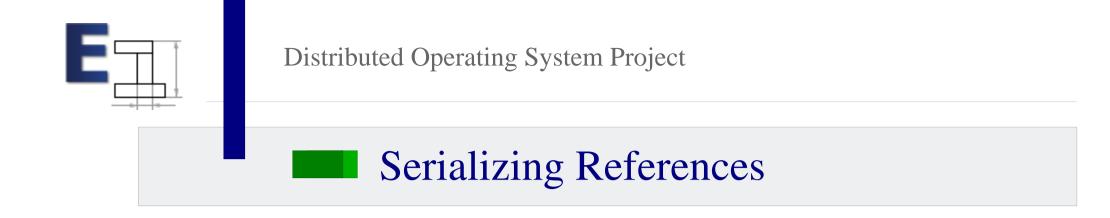


- Serialization can be cumbersome.
- Languages like Java and C# support automatic serialization based on RTTI.

• Problem: Implement automatic serialization for objects written in C/C++



- Dynamically allocated data;
- Static data (global variables);
- References to other objects;



References to other objects can be easily serialized by the operating system



- Dynamically allocated data;
- Static data (global variables);
- References to other objects;



We allow each object to have a separate private heap. Serialization of object dynamic data is then reduced to serialization of the heap.

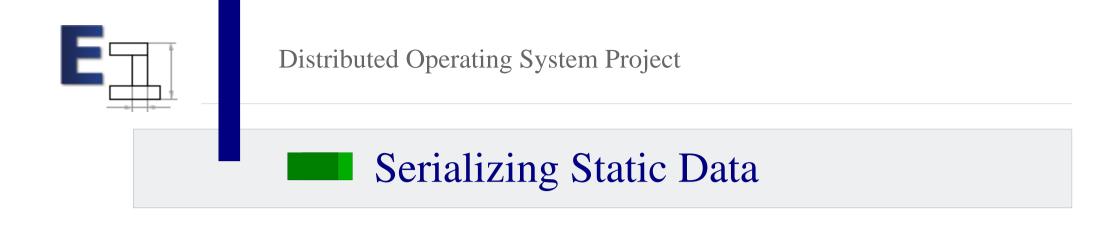
- Possible only in single address space.
- Memory overhead.
 - Should be acceptable for medium-grained objects.
 - For small objects manual serialization is not difficult.



- Dynamically allocated data;
- Static data (global variables);
- References to other objects;

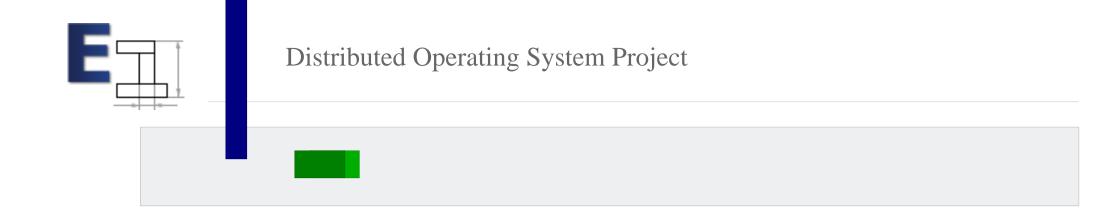


- We allow each object to have a separate copy of writable data segment for each module it depends on.
- It is allocated from the heap on object creation and is serialized/deserialized together with the heap.
- Problem: we have to switch data segment when crossing module and object boundaries.



We adopt Mungi approach with one modification.

- In Mungi all function pointers including C++ virtual method pointers are extended with *global pointer* field.
- In E1 this would require storing copies of all virtual tables in object heap.
- Instead, we store global pointer together with virtual table pointer in the object header.



Thank you