

Benchmarks Can Make Sense

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Motivation

Benchmarks are designed to compare systems

- Return a single number (e.g. throughput, average delay, completion time, CPU utilization, transactions per second)

No way to identify performance bottlenecks

- A simple configuration error can invalidate all results
- Netfork file system (NFS) benchmark below:
- No disk write buffering (50%)
- No RAID buffering
- Synchronous NFS mount (60%)
- Dafault NFS request buffer (30%)

- Bandwidth delay product (20%)

(50%)

Implementation

Full-system replay:



Xen virtual machine monitor

- Full-featured virtualization platform
- Support for production workloads

We extend it with a low-overhead recording

Goal

Benchmarks should make more sense

- Verify benchmark setup
- Aid perofmance analysis
 - Performance bottlenecks, verbose performance model

Idea

Full-system deterministic replay

- Lightweight non-intrusive way to record and replay execution

Replay is a mechanism allowing us to combine benchmarks and performance analysis

- Realistic execution
- Sophisticated computation heavy analysis

Analysis runs off-line on a cloned copy of execution

- Global comprehension
- No irreproducibility and observability problems - No restrictions on complexity of analysis (no probe effect)

Performance model:

Do performance measures stay sound during replay?

Replay

- 1) Set branch counters to overflow (cause exception)
- 2) Iterate in a single-step CPU mode to a target IP
- 3) Inject external event

Analysis interface:

Additional research needed to understand which information can actually help performance analysis

- Performance metrics

- Cooperative logging for network of machines
- Versioning storage for deterministic disk communication

Develop tools for execution comparison and automatic detection of non-determinism

- Hardware model during replay - Multiple exceptions (flushes CPU pipeline)
- Keeps caches almost warm

Analyze ILP during original and replay run

- Embarrasingly parallel

- Functional properties of the system



file system on a client machine. Processing of every file-system write involves two machines and multiple operating system components until it reaches the physical disk (request path is shown with a yellow line).

Multiple replay sessions allow us to run analysis multiple times. Constructing new analyses, we can compute various properties of the original run on demand.

Challenges

- Efficient full-system replay
- Faithful performance model
 - Replay may interfere with the execution of a system

Debugging and analysis interface

- Non-intrusive probes
- DTrace-like language interface to collect information about execution

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Implemented a basic deterministic logging and replay infrastructure

- Can replay beginning of the Linux boot (650K instructions)
- Replay mechanisms are designed to treat the state of a guest system as a set of memory pages
- Right choice to support heterogenity of replay in the future
- Support most non-deterministic events
 - Lack support for logging device driver communication