
L17: Generalizing CUDA: Concurrent Dynamic Execution, and Unified Address Space

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Administrative

- Remaining Lectures
 - Wednesday, April 17: Parallel Architectures and Getting to Exascale
- Final projects
 - Poster session, April 24 (dry run April 22)
 - Final report, May 1

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Final Project Presentation

- Dry run on April 22, Presentation on April 24
 - Easels, tape and poster board provided
 - Tape a set of Powerpoint slides to a standard 2'x3' poster, or bring your own poster.
- Final Report on Projects due May 1
 - Submit code
 - And written document, roughly 10 pages, based on earlier submission.
 - In addition to original proposal, include
 - Project Plan and How Decomposed (from DR)
 - Description of CUDA implementation
 - Performance Measurement
 - Related Work (from DR)

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Sources for Today's Lecture

References:

CUDA Dynamic Parallelism Programming Guide:

[http://docs.nvidia.com/cuda/pdf/
CUDA_Dynamic_Parallelism_Programming_Guide.pdf](http://docs.nvidia.com/cuda/pdf/CUDA_Dynamic_Parallelism_Programming_Guide.pdf)

 Examples: [http://developer.download.nvidia.com/GTC/PDF/
GTC2012/PresentationPDF/S0338-GTC2012-CUDA-
Programming-Model.pdf](http://developer.download.nvidia.com/GTC/PDF/GTC2012/PresentationPDF/S0338-GTC2012-CUDA-Programming-Model.pdf)

CUDA 5 Programming Guide

concurrentKernels example in /usr/local/cuda/samples/
6_AdvancedcdpAdvancedQuicksort example in /usr/local/cuda/samples/
6_Advanced

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Progress in Generalizing CUDA

- Early versions of CUDA had a lot of restrictions on parallelism that limited applicability and performance
 - First, synchronous kernel calls (host stalled)
 - Then, asynchronous, but only a single kernel executing on the GPU at a time (what we've been doing all semester)
 - More recently (CUDA 3 and compute capability 2), up to 4 kernels could be concurrently launched in the host using streams
 - Now in CUDA 5, kernels can launch other kernels (even recursive ones!)
- Early versions of CUDA dealt with memory in restrictive ways
 - Dynamic global and shared memory allocation only on host, otherwise static
 - Explicit copying between host and GPU required
 - Unified virtual address space added to CUDA 4, "pointers returned by `cudaHostAlloc()` can be used directly from within kernels running on these devices"

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Recall some forms of asynchronous overlap and concurrent execution (Lecture 7)

- Asynchronous copy and kernel execution
 - Using page-locked memory on host
- Concurrent Execution
 - Using streams (up to 4)

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Asynchronous Copy To/From Host (compute capability 1.1 and above)

- Concept:
 - Memory bandwidth can be a limiting factor on GPUs
 - Sometimes computation cost dominated by copy cost
 - But for some computations, data can be "tiled" and computation of tiles can proceed in parallel (some of your projects may want to do this, particularly for large data sets)
 - Can we be computing on one tile while copying another?
- Strategy:
 - Use page-locked memory on host, and asynchronous copies
 - Primitive `cudaMemcpyAsync`
 - Effect is GPU performs DMA from Host Memory
 - Synchronize with `cudaThreadSynchronize()`

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Page-Locked Host Memory

- How the Async copy works:
 - DMA performed by GPU memory controller
 - CUDA driver takes virtual addresses and translates them to physical addresses
 - Then copies physical addresses onto GPU
 - Now what happens if the host OS decides to swap out the page???
- Special malloc holds page in place on host
 - Prevents host OS from moving the page
 - `CudaMallocHost()`
- But performance could degrade if this is done on lots of pages!
 - Bypassing virtual memory mechanisms

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Example of Asynchronous Data Transfer

```

cudaStreamCreate(&stream1);
cudaStreamCreate(&stream2);
cudaMemcpyAsync(dst1, src1, size, dir, stream1);
kernel<<<grid, block, 0, stream1>>>(…);
cudaMemcpyAsync(dst2, src2, size, dir, stream2);
kernel<<<grid, block, 0, stream2>>>(…);

```

src1 and src2 must have been allocated using cudaMallocHost
 stream1 and stream2 identify streams associated with asynchronous
 call (note 4th "parameter" to kernel invocation, by default there is one
 stream)

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Code from asyncAPI SDK project

```

// allocate host memory
CUDA_SAFE_CALL( cudaMallocHost((void**)&a, nbytes) );
memset(a, 0, nbytes);

// allocate device memory
CUDA_SAFE_CALL( cudaMalloc((void**)&d_a, nbytes) );
CUDA_SAFE_CALL( cudaMemset(d_a, 255, nbytes) );

... // declare grid and thread dimensions and create start and stop events

// asynchronously issue work to the GPU (all to stream 0)
cudaEventRecord(start, 0);
cudaMemcpyAsync(d_a, a, nbytes, cudaMemcpyHostToDevice, 0);
increment_kernel<<<blocks, threads, 0, 0>>>(d_a, value);
cudaMemcpyAsync(a, d_a, nbytes, cudaMemcpyDeviceToHost, 0);
cudaEventRecord(stop, 0);

// have CPU do some work while waiting for GPU to finish

// release resources
CUDA_SAFE_CALL( cudaFreeHost(a) );
CUDA_SAFE_CALL( cudaFree(d_a) );

```

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Concurrent Execution (Compute Capability 2.0)

- Stream concept: create, destroy, tag asynchronous operations with stream
 - Special synchronization mechanisms for streams: queries, waits and synchronize functions
- Concurrent Kernel Execution
 - Execute multiple kernels (up to 4) simultaneously
 - Example: concurrentKernels in samples/6_Advanced

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CUDA Dynamic Parallelism

- Key Idea: Ability to dynamically create CUDA threads on the GPU from kernel functions
- Things to think about
 - What is the syntax?
 - How does this affect the execution model?
 - How does a thread know when launched threads have completed?
 - What does a __syncthreads() mean?
 - What happens to memory state, consistency?
- Let's look at some examples
 - Hello World
 - From GTC presentation
 - From Quicksort

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