CS4230 Parallel Programming

Lecture 9: Locality and Data Parallel Algorithms

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Administrative

• SVD assignment status
  - Problems with speedup on water
  - I suspect OMP implementation
  - Will resolve today in some way with Axel
More locality: Image correlation

... Initialize th[i][j] = 0 ...

/* compute array convolution */

for(m = 0; m < IMAGE_NROWS - TEMPLATE_NROWS + 1; m++){
    for(n = 0; n < IMAGE_NCOLS - TEMPLATE_NCOLS + 1; n++){
        for(i=0; i < TEMPLATE_NROWS; i++){
            for(j=0; j < TEMPLATE_NCOLS; j++){
                if(mask[i][j] != 0) {
                    th[m][n] += image[i+m][j+n];
                }
            }
        }
    }
}

/* scale array with bright count and template bias */

... th[i][j] = th[i][j] * bc - bias;
The Red/Blue computation simulates two interactive flows. An n x n board is initialized so cells have one of three colors: red, white, and blue, where white is empty, red moves right, and blue moves down. Colors wrap around on the opposite side when reaching the edge.

In the first half step of an iteration, any red color can move right one cell if the cell to the right is unoccupied (white). On the second half step, any blue color can move down one cell if the cell below it is unoccupied. The case where red vacates a cell (first half) and blue moves into it (second half) is okay.

Viewing the board as overlaid with t x t tiles (where t divides n evenly), the computation terminates if any tile’s colored squares are more than c% one color. Use Peril-L to write a solution to the Red/Blue computation.
Steps for “Localized” Solution

Step 1. Partition global grid for n/t x n/t processors
Step 2. Initialize half iteration (red) data structure
Step 3. Iterate within each t x t tile until convergence (guaranteed?)
  Step 4. Compute new positions of red elts & copy blue elts
Step 5. Communicate red boundary values
Step 6. Compute new positions of blue elts
Step 7. Communicate blue boundary values
Step 8. Check locally if DONE
Steps 1, 3 & 8. Partition Global Grid and Iterate

```java
int grid[n,n], lgrid[t,t];
boolean gdone = FALSE;
int thr = (n/t)*(n/t);

forall (index in(0..thr-1))
    int myx = (n/t) * index/(n/t);
    int myy =(n/t) * (index % (n/t));
lgrid[] = localize(grid[myx,myy]);
while (!gdone) {
    // Steps 3-7: compute new locations and determine if done
    if (Bsum > threshold || Wsum > threshold || Rsum > threshold)
        exclusive { gdone = TRUE; }
    barrier;
}
```
Step 2: Initialization of lr the first time

```plaintext
int lr[t+2,t+2], lb[t+2,t+2];
// copy from lgrid and grid to lr
lr[1:t,1:t] = lgrid[0:t-1,0:t-1];
lr[0,1:t] = grid[myx,myy:myy+t-1];
lr[t,1:t+1] = ...; lr[1:t,0] = ...; lr[1:t,t+1] = 
lr[0,0] = lr[0,t+1] = lr[t+1,0] = lr[t+1,t+1] = W; // don't care about these
```

```
Boundary is ghost cells
```
Steps 4 & 5: Compute new positions of red elts and Communicate Boundary Values

```c
int lb[t+1,t+1];
lb[0:t+1,0:t+1] = W;
for (i=0; i<t+1; i++) {
    for (j=0; j<t+1; j++) {
        if (lr[i,j] == B) lb[i,j] = B;
        if (lr[i,j] == R && lr[i+1,j] = W) lb[i+1,j] = R;
        else lb[i,j] = R;
    }
}

barrier;
lgrid[0:t-1,0:t-1] = lb[1:t,1:t]; //update local portion of global ds
barrier;

//copy leftmost ghosts from grid
if (myx-1 >= 0) lb[0,1:t] = grid[myx-1, 0:t-1];
else lb[0,1:t] = grid[n-1,0:t-1];
```
Steps 6 & 7: Compute new positions of blue elts and communicate Boundary Values

```c
int lb[t+1,t+1];

lr[0:t+1,0:t+1] = W;
for (i=0; i<t+1; i++) {
    for (j=0; j<t+1; j++) {
        if (lb[i,j] == R) then lr[i,j] = R;
        else lr[i,j] = B;
    }
}
barrier;
lgrid[0:t-1,0:t-1] = lr[1:t,1:t]; //update local portion of global ds
barrier;

//copy top ghosts from grid
if (myy-1 >= 0) lr[1:t,0] = grid[0:t-1,myy-1];
else lr[1:t,0] = grid[0:t-1,0];
```
Step 8: Compute locally if DONE

for (i=1; i<=t+1; i++) {
    for (j=1; j<=t+1; j++) {
        if (lr[i,j] == R) then Rsum++;
        if (lr[i,j] == B) then Bsum++;
        if (lr[i,j] == W) then Wsum++;
    }
}