Instructions You will have eighty minutes to complete the actual open-book, open-note exam. Electronic devices will be allowed only to consult notes or books from local storage; network use will be prohibited. The actual exam will be shorter than this practice exam and will not include a problem like number 21.

The first three questions refer to the code template

```c
long f(long x, long y) {
    return ....;
}
```

and the code fragments

1. `x*(2+y)`
2. `x+8*y`
3. `((long *)x)[y]`
4. `*(long *)x = y`
5. `x == y`

1. Which of the above code fragments most likely appears in place of .... if the generated machine code includes the following instruction?

   leaq (%rdi, %rsi, 8), %rax

2. Which of the above code fragments most likely appears in place of .... if the generated machine code includes the following instruction?

   movq (%rdi, %rsi, 8), %rax

3. Which of the above code fragments most likely appears in place of .... if the generated machine code includes the following sequence?

   cmpq %rdi, %rsi
   sete %al
   movzbq %rax


The next four questions refer to the following assembly sequence:

\[
\begin{align*}
&\text{leal (, %eax, 4), %ebx} \\
&\text{shrl $0x2, %ebx} \\
&\text{subl %ebx, %eax} \\
&\text{cmpl $0x0, %eax} \\
&\text{jg .L1} \\
&\text{andl $0x7FFFFFFF, %eax} \\
&\text{.L1: } \\
&\text{shrl $29, %eax} \\
&\text{ret}
\end{align*}
\]

4. If %eax starts as 0, what is the ending value of register %eax?

5. If %eax starts as -1, what is the ending value of register %eax?

6. If %eax starts as 0x4FFFFFFF, what is the ending value of register %eax?

7. If %eax starts as unknown, what are all the possible ending values of %eax?
Given that

```c
int sum_element(int c,
    long mat1[][M],
    long mat2[][N]) {
    int k;
    long sum = 0;
    for (k = 0; k < c; k++)
        sum += mat1[k][c] - mat2[k][c];
    return sum;
}
```

compiles as

```assembly
    testl %edi, %edi
    jle zero
    movslq %edi, %rax
    subl $1, %edi
    leaq 0(,%rax,8), %r8
    leaq (%rax,%rdi,4), %rax
    leaq (%rsi,%r8), %rcx
    addq %r8, %rdx
    leaq 32(%rsi,%rax,8), %rsi
    xorl %eax, %eax
    loop:
        addq (%rcx), %rax
        addq $32, %rcx
        subq (%rdx), %rax
        addq $40, %rdx
        cmpq %rsi, %rcx
        jne loop
    ret
    zero:
        xorl %eax, %eax
    ret
```

then what are the values of the constants M and N among the following possibilities?

- M = 10 and N = 6
- M = 12 and N = 2
- M = 3 and N = 5
- M = 4 and N = 5
- M = 5 and N = 12
- M = 8 and N = 17

To enable partial credit, show your reasoning for how the assembly code corresponds to elements of the C code.
The next two questions refer to the iterate function defined as

```c
double iterate(double x, double y, double z, int steps) {
    int i;
    for (i = 0; i < steps; i++) {
        z = x * y;
        if (!(i & 0x1))
            x = y + z;
        else
            y = x + z;
    }
    return z;
}
```

9. Which of the following correctly represents the dependency graph of iterate over six iterations, where each column corresponds to a single iteration?

10. Based on the dependency graph, how many cycles will iterate take (expressed as a multiple of steps)?
The following four questions refer to the following declarations:

```c
typedef struct jumble {
    struct {
        int i;
    } r;
    double d;
    char c;
    union {
        int i;
        char *s;
    } u;
    short s;
} jumble;

jumble ja[7][5];
```

For the questions, assume that `ja` starts at address 0x10000.

11. What is `sizeof(jumble)`?

12. What is the address of `ja[0][2].d`?

13. What is the address of `ja[4][2].u.s`?

14. What is the address of `ja[4][2].s`?
The next two questions refer to the following declarations and function:

```c
typedef struct {
    int a[2];
} pair;

pair p[32][32];

int sum(int how_many) {
    int i, j, k, a = 0;
    for (i = 0; i < 32; i++)
        for (j = 0; j < 32; j++) {
            for (k = 0; k < how_many; k++) {
                a += s[i][j].a[k];
            }
        }
    return a;
}
```

For each question below, assume a 1kB direct-mapped cache that uses 16-byte blocks, the cache is initially empty, and local variables are in registers. Also assume that the array p is at the address 0x10000 in memory.

15. For each of first five memory accesses via p in sum(1), what is the accessed element, what is the accessed address, and is the access a cache hit or miss?

<table>
<thead>
<tr>
<th>Access expression</th>
<th>Access address</th>
<th>Hit or miss</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

16. What is the overall expected cache-miss rate of sum(1)?
17. For each of first five memory accesses via \( p \) in \( \text{sum}(2) \), what is the accessed element, what is the accessed address, and is the access a cache hit or miss?

<table>
<thead>
<tr>
<th>Access expression</th>
<th>Access address</th>
<th>Hit or miss</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

18. What is the overall expected cache-miss rate of \( \text{sum}(2) \)?
19. What are all of the possible outputs of the following program?
Assume that the program runs on Linux, where the child process created by fork does \textit{not} include all the threads of the parent process. The child process includes only the thread that called \texttt{fork}.

In case you don't list all of the possible outputs correctly, to improve opportunities for partial credit, show how you arrived at your answer by sketching one or more process/thread graphs.

```
#include "csapp.h"

static sem_t done;

static void *go(void *p) {
    pid_t pid = fork();
    if (pid == 0) {
        Write(1, "c", 1);
        V(&done);
        exit(0);
    } else {
        Write(1, "p", 1);
        Waitpid(pid, NULL, 0);
        V(&done);
        return NULL;
    }
}

int main() {
    pid_t pid;
    pthread_t th1, th2;

    Sem_init(&done, 0, 0);

    Write(1, "+", 1);
    Pthread_create(&th1, NULL, go, NULL);
    Pthread_create(&th2, NULL, go, NULL);

    P(&done);
    P(&done);
    Write(1, "-", 1);

    return 0;
}
```
20. For this question, a word is defined to be 22 bytes, each cell in a diagram represents a word, and an underlined number \( N \) is a shorthand for \( N \) times 16.

An allocator produces word-aligned payload pointers, uses a word-sized header and footer for each allocated block, uses a 2-word prolog block and a 1-word terminator block, coalesces unallocated blocks, uses an explicit free list as a singly-linked and NULL-terminated list starting from \( fp \) and using the first word of a block payload for a “next” pointer, adds free blocks to the start of the free list, does not split a block to create an unallocated block with an empty payload, uses a best-fit allocation strategy, and is confined to 16 words of memory that is initially filled with 0s. Show a header in a diagram as a value for the block size over a 0 or 1 to indicate the block’s allocation status; draw a footer as just the block size. Show the free list as either equal to NULL or as a chain of pointers.

The left-hand column below contains a sequence of malloc and free calls that are handled by the allocator. Fill in the left-hand column by showing relevant header and footer values and the free list just after each step on the left. The first row of the left column is blank so that you can show the initial state of memory in the first row of the right column.

\[
\begin{array}{c}
\text{fp}= \\
p1 = \text{malloc}(2) \\
p2 = \text{malloc}(2) \\
p3 = \text{malloc}(2) \\
p4 = \text{malloc}(2) \\
\text{free}(p1) \\
\text{free}(p3) \\
p5 = \text{malloc}(2)
\end{array}
\]
The last page contains a server implementation that implements a vote-counting service, except that many parts of the program have been deleted and replaced with blanks. Your task is to fill in each blank with a number that corresponds to one of the following program fragments. Each fragment can be used multiple times, and not all fragments need to be used, but you can only fill in blanks with these fragments. If you use a fragment that itself has blanks, then those blanks must be filled in, too; blanks in unused fragments do not need to be filled in.

The server listens for both TCP connections and UDP connections at a given port. Each message sent to the server by UDP is received as a vote, which is recorded by incrementing a counter for the message string as stored in a dictionary. For example, a client could send the message apple to increase the apple count. For each connection made to the server via TCP, the server sends back vote results so far. Vote results are reported as an int for the number of voted-on strings, then, for each voted-on string: each string’s length as an int, its content as bytes, and its vote count as an int.

The server must listen for UDP and TCP connections currently. Since each UDP message can be handled immediately, there’s no need to have multiple messages processed concurrently to each other. Each TCP connection may take a while to receive all the votes, however, so the server is meant to support multiple TCP connections currently. Finally, each TCP connection should be sent a snapshot of vote tallies at the time of the connection, and new votes via UDP should be allowed meanwhile.

The server uses a dictionary library with the following property: the dictionary makes its own copy of keys, but when a value is installed into the dictionary with dictionary_set, the dictionary takes ownership of the value and is responsible for deallocating it when the value is replaced (or when the dictionary is destroyed, but that never happens in this server). A client does not take ownership of a value that is extracted from the dictionary with dictionary_get; that value can be deallocated if it is meanwhile replaced.

1. Close(s);
2. ls = Open_listenfd(portno);
3. struct sockaddr_storage addr;
   unsigned int len = sizeof(addr);
   s = Accept(ls, (struct sockaddr *)&addr, &len);
4. s = Socket(addrs->ai_family, addrs->ai_socktype, addrs->ai_protocol);
5. amt = Sendto(s, buffer, MAXBUF-1, 0,
            (struct sockaddr *)&addr, &len);
6. char buffer[MAXBUF];
   size_t amt;
   amt = Recv(s, buffer, MAXBUF-1, 0);
   buffer[amt] = 0;
7. free(new_p);
8. free(p);
9. static int make_socket(char *portno) {
   int s;
   struct addrinfo hints, *addrs;
   memset(&hints, 0, sizeof(struct addrinfo));
   hints.ai_family = AF_INET;
   hints.ai_socktype = SOCK_STREAM;
   Getaddrinfo("localhost", portno, &hints, &addrs);
   Connet(s, addrs->ai_addr, addrs->ai_addrlen);
   Freeaddrinfo(addrs);
   return s;
}
10. static int make_socket(char *portno) {
    int s;
    struct addrinfo hints, *addrs;
    memset(&hints, 0, sizeof(struct addrinfo));
    hints.ai_family = AF_INET;
    hints.ai_socktype = SOCK_DGRAM;
    hints.ai_flags = AI_PASSIVE;
    Getaddrinfo(NULL, portno, &hints, &addrs);
    Bind(s, addrs->ai_addr, addrs->ai_addrlen);
    Freeaddrinfo(addrs);
    return s;
}
11. s
12. ls
13. Sem_init(&db_lock, 0, 1);
14. static void sigchld_handler(int sig) {
    while (waitpid(-1, NULL, WNOHANG) < 0)
    ;
}
15. Signal(SIGCHLD, sigchld_handler);
16. V(&db_lock);
17. P(&db_lock);
18. Waitpid(-1, NULL, 0);
19. exit(0);
20. receive_votes
21. serve_counts
22. static void write_entry(dictionary_t *db, int i, int s) {
    const char *key = dictionary_key(db, i);
    int len = strlen(key);
    int votes = *(int *)dictionary_value(db, i);
    Rio_writen(s, &len, sizeof(len));
    Rio_writen(s, (void *)key, len);
    Rio_writen(s, &votes, sizeof(votes));
}
23. /* A million threads isn't cool... */
21. Fill in the blanks with numbers for fragments on the previous page:

```c
#include "csapp.h"
#include "dictionary.h"

static dictionary_t *db;
static sem_t db_lock;

```

```
c
d
e

```

static void receive_votes(char *portno)
{
    int s = make_socket(portno);
    while (1) {
        int *p, *new_p;
        
        new_p = malloc(sizeof(int));
        *new_p = 1;
        p = dictionary_get(db, buffer);
        if (p) *new_p += *p;
        dictionary_set(db, buffer, new_p);
    }
}

```

```
f
g
h
```

static void *serve_counts(void *portno)
{
    int ls, s;
    
    while (1) {
        
        if (Fork() == 0) {
            int i, count = dictionary_count(db);
            Rio_write(s, &count, sizeof(count));
            for (i = 0; i < count; i++)
                write_entry(db, i, s);
        }
    }
}

```

```
i
j
k
l
m
```

int main(int argc, char **argv)
{
    pthread_t th;
    if (argc != 2) app_error("need <port>");
    db = make_dictionary(COMPARE_CASE_SENS, free);
    Sem_init(&db_lock, 0, 1);
    
    Pthread_create(&th, NULL, ___________o, argv[1]);
    Pthread_detach(th);
    __________p(argv[1]);
    return 0;
}
```
Answers

1. \( x + 8 \cdot y \)

2. \((\text{long } *)x)[y]\)

3. \(x == y\)

4. 0

5. 2

6. 2

7. 0 and 2

8. \(M = 4\) and \(N = 5\)

```c
    test1 %edi, %edi          # check whether c is 0
       jle zero               # if c is zero, just return 0
    movelq %edi, %rax         # copy c to rax
    subl $1, %edi             # rdi = c - 1
    leaq 0(%rax,8), %r8       # r8 = c*sizeof(long)
    leaq (%rax,%rdi,4), %rax  # rax = 5*(c-1)
    leaq (%rsi,%r8), %rcx     # rcx = matrix1 + c*sizeof(long)
    addq %r8, %rdx            # rdx = matrix2 + c*sizeof(long)
    leaq 32(%rsi,%rax,8), %rsi # rsi = matrix1 + 5*(c-1)*8+32; used as end condition
    xorl %eax, %eax           # sum = 0

loop:
    addq (%rcx), %rax         # sum += rcx[0]
    addq $32, %rcx            # rcx += 4 * sizeof(long) => M = 4
    subq (%rdx), %rax         # sum -= rdx[0]
    addq $40, %rdx            # rdx += 5 * sizeof(long) => N = 5
    cmpq %rsi, %rcx           # rcx == matrix1+c*sizeof(long)?
    jne loop                 # if we haven’t reached last row....
    ret
zero:
    xorl %eax, %eax           # sum = 0
    ret
```

9. The third one

10. 8, since there’s a path that goes through both FP* and FP+ at each iteration.

11. 40

```c
typedef struct jumble {
    struct {
        int i;       // size 4, alignment 4 => offset 0
    } r;          // also size 4, offset 0
    double d;     // size 8, alignment 8 => offset 8
    char c;       // size 1, alignment 1 => offset 16
} union {
    int i;       // size 4, alignment 2 => offset 20... but 24 for s
```
char *s; // size 8, alignment 8 => offset 24
} u; // so, size 8, offset 24
short s; // size 2, alignment 2 => offset 32
} jumble; // size >= 34, alignment 8 => size 40

jumble ja[7][5]; // each row is 40*5 bytes

12. 0x10000 + 80 + 8 = 0x10058
13. 0x10000 + 40*5*4 + 80 + 24 = 0x10388
14. 0x10000 + 40*5*4 + 80 + 32 = 0x10390
15.

<table>
<thead>
<tr>
<th>Access expression</th>
<th>Access address</th>
<th>Hit or miss</th>
</tr>
</thead>
<tbody>
<tr>
<td>s[0][0].a[0]</td>
<td>0x10000</td>
<td>miss</td>
</tr>
<tr>
<td>s[0][1].a[0]</td>
<td>0x10008</td>
<td>hit</td>
</tr>
<tr>
<td>s[0][2].a[0]</td>
<td>0x10010</td>
<td>miss</td>
</tr>
<tr>
<td>s[0][3].a[0]</td>
<td>0x10018</td>
<td>hit</td>
</tr>
<tr>
<td>s[0][4].a[0]</td>
<td>0x10020</td>
<td>miss</td>
</tr>
</tbody>
</table>

16. 50%
17.

<table>
<thead>
<tr>
<th>Access expression</th>
<th>Access address</th>
<th>Hit or miss</th>
</tr>
</thead>
<tbody>
<tr>
<td>s[0][0].a[0]</td>
<td>0x10000</td>
<td>miss</td>
</tr>
<tr>
<td>s[0][0].a[1]</td>
<td>0x10004</td>
<td>hit</td>
</tr>
<tr>
<td>s[0][1].a[0]</td>
<td>0x10008</td>
<td>hit</td>
</tr>
<tr>
<td>s[0][1].a[1]</td>
<td>0x1000c</td>
<td>hit</td>
</tr>
<tr>
<td>s[0][2].a[0]</td>
<td>0x10010</td>
<td>miss</td>
</tr>
</tbody>
</table>

18. 25%
19. Six possible outputs: +ppcc-, +pcpc-, +ccpp-, +pccp-, +cpcp-, and +ccpp-.
The outputs reflect that + must be first, - must be last, and the two cs and two ps can be in any order. A possible progress/thread graph is
It’s also possible for the Vs and Ps to pair up in the other way, but that makes no difference to the output. The V(done_copy) nodes reflect that the child processes perform a semaphore post on done, but it has no effect since the semaphore is not shared with the parent process.

20.

In the last step, allocating the middle free block would also be valid.

21.

\begin{align*}
b &= 4 \\
c &= 14 \\
d &= 10 \\
e &= 22 \\
f &= 6 \\
g &= 17 \\
h &= 16 \\
i &= 2 \\
j &= 3 \\
k &= 17 \\
l &= 19 \\
m &= 16 \\
n &= 15 \\
o &= 21 \\
p &= 20
\end{align*}