Sample Mid-Term Exam 1

CS 3520, Fall 2005

September 21, 2005

1) Given the following grammar:

```
\langle \operatorname{weed} 
angle =  'leaf | (list 'branch \langle \operatorname{weed} 
angle | (list 'stem \langle \operatorname{weed} 
angle)
```

Which of the following expressions are examples of $\langle weed \rangle$?

- a) (list 'leaf)
- b) (list 'stem)
- c) (list 'branch (list 'branch 'leaf 'leaf) 'leaf)
- d) (list 'stem 'leaf)
- **2**) Explain why the following is a $\langle \text{weed} \rangle$:

```
(list 'branch (list 'stem 'leaf) (list 'branch 'leaf 'leaf))
```

3)	Provide a define-type declaration for Weed that is a suitable representation for	$\langle \mathrm{weed} \rangle \mathrm{s}.$	
4)	Implement the function $\mathtt{weed\text{-}forks}$, takes a $\langle \mathtt{weed} \rangle$ and returns the number contains. Your implementation must follow the shape of the data definition.	of branches that it	i

5) Given the following expression in the book language with with, fun, and rec:

- a) Draw arrows on the above expression from each bound variable to its binding occurrence.
- b) List the free variables:

and bound variables:

6) Given the following expression:

```
{with {g {fun {x} {fun {y} {+ y x}}}}
    {with {x 13}
        {with {f {g 6}}
        {f x}}}
```

Describe a trace of the evaluation in terms of arguments to an interp function for every call. (There will be 16 calls.) The interp function takes two arguments — an expression and a substitution cache — so show both for each call. For number, variable, and fun expressions, show the result value, which is immediate. Use the back of the exam for additional space, and use the following abbreviations to save time:

```
\begin{array}{lcl} E_0 & = & \text{the whole expression} \\ E_1 & = & \{\text{fun } \{\text{x}\} \; \{\text{fun } \{\text{y}\} \; \{+ \; \text{y} \; \text{x}\}\}\} \\ E_2 & = & \{\text{with } \{\text{x } \; \text{13}\} \; \{\text{with } \{\text{f } \{\text{g } \; \text{6}\}\} \; \{\text{f } \; \text{x}\}\}\} \\ E_3 & = & \{\text{with } \{\text{f } \{\text{g } \; \text{6}\}\} \; \{\text{f } \; \text{x}\}\} \end{array}
```

Answers

- 1) (c) and (d)
- 2) Since 'leaf is a (weed) by line 1 of the definition, then by line 3, (list 'stem 'leaf) is a (weed), and by line 2, (list 'branch 'leaf 'leaf) is a (weed). Finally, then, by line 2 again, (list 'branch (list 'stem 'leaf) (list 'branch 'leaf 'leaf)) is a (weed).

```
3) (define-type Weed
     [leaf]
      [stem (rest Weed?)]
      [branch (left Weed?)
               (right Weed?)])
4) ; weed-forks : Weed -> num
   (define (weed-forks w)
     (type-case Weed w
       [leaf () 0]
       [stem (rest) (weed-forks rest)]
       [branch (l r) (+ 1
                          (weed-forks 1)
                          (weed-forks r))]))
   (test (weed-forks (leaf)) 0)
   (test (weed-forks (stem (leaf))) 0)
   (test (weed-forks (stem (branch (leaf) (leaf)))) 1)
   (test (weed-forks (branch (branch (leaf) (leaf)) (leaf))) 2)
          v----,
{rec {g {fun {z} {f z}}}}
^-----,
,---+-,
5)
            {rec {f {fun {z} {g z}}}}
               ,---+---, {with {y {with {f {fun {z} {f {+ z x}}}}}
                    -^ ,---^
                         (f y}}}
                 {+ y q}}}
   Free: f, x, y, q Bound: z, g, f, y
6)
                 expr
                              E_0
                              (mtSub)
                 subs
                 expr
                             |E_1|
                 subs
                             (mtSub)
                             (closureV 'x \boxed{\{\texttt{fun } \{\texttt{y}\}\ \{\texttt{+ y x}\}\}} (mtSub)) = C_1
                 result
                 expr
```

 $\overline{\text{(aSub 'g } C_1 \text{ (mtSub))}} = S_1$

subs

$$\begin{array}{lll} \operatorname{expr} &=& 13 \\ \operatorname{subs} &=& S_1 \\ \operatorname{result} &=& (\operatorname{numV 13}) \end{array} \\ \operatorname{expr} &=& E_3 \\ \operatorname{subs} &=& (\operatorname{aSub 'x (numV 13) } S_1) = S_2 \\ \operatorname{expr} &=& \left\{ \mathbf{g 6} \right\} \\ \operatorname{subs} &=& S_2 \\ \operatorname{expr} &=& \mathbf{g} \\ \operatorname{subs} &=& S_2 \\ \operatorname{result} &=& C_1 \\ \end{array} \\ \operatorname{expr} &=& \left\{ \mathbf{6} \\ \operatorname{subs} &=& S_2 \\ \operatorname{result} &=& (\operatorname{numV 6}) \end{array} \\ \operatorname{expr} &=& \left\{ \mathbf{fun \{y\} \{+ \ y \ x\}\}} \right\} \\ \operatorname{subs} &=& (\operatorname{aSub 'x (numV 6) (mtSub))} = S_3 \\ \operatorname{result} &=& (\operatorname{closureV 'y \{+ \ y \ x\}\} } S_3) = C_2 \\ \end{array} \\ \operatorname{expr} &=& \left\{ \mathbf{f x} \right\} \\ \operatorname{subs} &=& (\operatorname{aSub 'f } C_2 \ S_2) = S_4 \\ \end{array} \\ \operatorname{expr} &=& \left\{ \mathbf{f x} \right\} \\ \operatorname{subs} &=& S_4 \\ \operatorname{result} &=& C_2 \\ \end{array} \\ \operatorname{expr} &=& \left\{ \mathbf{f x} \right\} \\ \operatorname{subs} &=& S_4 \\ \operatorname{result} &=& (\operatorname{numV 13}) \\ \operatorname{expr} &=& \left\{ \mathbf{f y x} \right\} \\ \operatorname{env} &=& \left\{ (\operatorname{aSub 'y (numV 13) } S_3) = S_5 \\ \end{array} \\ \operatorname{expr} &=& \left\{ \mathbf{f x y x} \right\} \\ \operatorname{env} &=& \left\{ (\operatorname{aSub 'y (numV 13) } S_3) = S_5 \\ \end{array} \\ \end{array} \\ \operatorname{expr} &=& \left\{ \mathbf{f x y x} \right\} \\ \operatorname{env} &=& \left\{ (\operatorname{aSub 'y (numV 13) } S_3) = S_5 \\ \end{array} \\ \operatorname{expr} &=& \left\{ (\operatorname{aSub 'y (numV 13) } S_3) = S_5 \\ \end{array} \\ \operatorname{expr} &=& \left\{ (\operatorname{aSub 'y (numV 13) } S_3) = S_5 \\ \end{array} \\ \end{array}$$

(numV 6)

result

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