

Functional Programs

So far, the language that we've implemented is purely *functional*

- A function produces the same result every time for the same arguments
- Also, lazy and eager results are the same
 - ... except that eager evaluation might loop forever or raise an exception where the lazy version produces a result

Non-Functional Procedures

```
(define (f x)
  (+ x (read)))
```

```
(define counter 0)
(define (f x)
  (begin
    (set! counter (+ x counter))
    counter))
```

```
(define f
  (local [(define b (box 0))]
    (lambda (x)
      (begin
        (set-box! b (+ x (unbox b)))
        (unbox b))))))
```

BCFAE = FAE + Boxes

```
<BCFAE> ::= <num>
          | { + <BCFAE> <BCFAE> }
          | { - <BCFAE> <BCFAE> }
          | <id>
          | { fun { <id> } <BCFAE> }
          | { <BCFAE> <BCFAE> }
          | { if0 <BCFAE> <BCFAE> <BCFAE> }
          | { newbox <BCFAE> }
          | { setbox <BCFAE> <BCFAE> }
          | { openbox <BCFAE> }
          | { seqn <BCFAE> <BCFAE> }
```

NEW

NEW

NEW

NEW

```
{with {b {newbox 0}}
  {seqn
    {setbox b 10}
    {openbox b}}}
```

⇒ 10

Implementing Boxes with Boxes

```
(define-type BCFAE-Value
  [numV (n number?)]
  [closureV (param symbol?)
            (body BCFAE?)
            (sc SubCache?)]
  [boxV (container (box-of BCFAE?))])
```

Implementing Boxes with Boxes

```
; interp : BCFAE SubCache -> BCFAE-Value
(define (interp a-bcfae sc)
  (type-case RCFAE a-bcfae
    ...
    [newbox (val-expr)
            (boxV (box (interp val-expr sc)))]
    [setbox (box-expr val-expr)
            (set-box! (boxV-container
                      (interp box-expr sc))
                      (interp val-expr sc))]
    [openbox (box-expr)
             (unbox (boxV-container
                    (interp box-expr sc)))]))
```

But this doesn't explain anything about boxes!

Boxes and Memory

`{with {b {newbox 7}}}` \Rightarrow `...`
`...`

Memory:

Memory:

			7	

Boxes and Memory

... {setbox b 10}
...

⇒

... {openbox b}
...

Memory:

			7	

Memory:

			10	

The Store

We represent memory with a **store**:

```
(define-type Store
  [mtSto]
  [aSto (address integer?)
        (value BCFAE-Value?)
        (rest Store?)])
```

Memory:

			10	

```
(aSto 13 (numV 10)
      (mtSto))
```


Implementing Boxes without State

```
; interp : BCFAE SubCache Store -> Value*Store

(define-type BCFAE-Value
  [numV (n number?)]
  [closureV (param symbol?)
             (body BCFAE?)
             (sc SubCache?)]
  [boxV (address integer?)])

(define-type Value*Store
  [v*s (value BCFAE-Value?)
       (store Store?)])
```

Implementing Boxes without State

```
; interp : BCFAE SubCache Store -> Value*Store
(define (interp expr sc st)
  ...
  [newbox (expr)
    (type-case Value*Store (interp expr sc st)
      [v*s (val st)
        (local [(define a (malloc st))]
          (v*s (boxV a)
              (aSto a val st))))))]
  ... )

; malloc : Store -> integer
```

Implementing Boxes without State

```
; malloc : Store -> integer
(define (malloc st)
  (+ 1 (max-address st)))

; max-address : Store -> integer
(define (max-address st)
  (type-case Store st
    [(mtSto) 0]
    [(aSto n v st)
     (max n (max-address st))]))
```

Implementing Boxes without State

```
; interp : BCFAE SubCache Store -> Value*Store
(define (interp expr sc st)
  ...
  [openbox (bx-expr)
    (type-case Value*Store (interp bx-expr sc st)
      [v*s (bx-val st)
        (v*s (store-lookup (boxV-address bx-val)
                           st)
              st)]])]
  ...)
```

Implementing Boxes without State

```
; interp : BCFAE SubCache Store -> Value*Store
(define (interp expr sc st)
  ...
  [setbox (bx-expr val-expr)
    (type-case Value*Store (interp bx-expr sc st)
      [v*s (bx-val st2)
        (type-case Value*Store (interp val-expr sc st2)
          [v*s (val st3)
            (v*s val
              (aSto (boxV-address bx-val)
                    val
                    st3))]]))]
  ...)
```

seqn, **add**, **sub**, and **app** will need the same sort of sequencing

Implementing Boxes without State

```
; interp-two : (BCFAE BCFAE SubCache Store
;             (Value Value Store -> Value*Store)
;             -> Value*Store)
(define (interp-two expr1 expr2 sc st handle)
  (type-case Value*Store (interp expr1 sc st)
    [v*s (val1 st2)
      (type-case Value*Store (interp expr2 sc st2)
        [v*s (val2 st3)
          (handle val1 val2 st3)]))]))
```

Implementing Boxes without State

```
; interp : BCFAE SubCache Store -> Value*Store
(define (interp expr sc st)
  ...
  [add (r l) (interp-two r l sc st
                        (lambda (v1 v2 st)
                          (v*s (num+ v1 v2) st)))]
  ...
  [seqn (a b) (interp-two a b sc st
                          (lambda (v1 v2 st)
                            (v*s v2 st)))]
  ...
  [setbox (bx-expr val-expr)
          (interp-two bx-expr val-expr sc st
                    (lambda (bx-val val st3)
                      (v*s val
                          (aSto (boxV-address bx-val)
                                val
                                st3)))))]
  ...)
```

Variables

Boxes don't explain one of our earlier Scheme examples:

```
(define counter 0)
(define (f x)
  (begin
    (set! counter (+ x counter))
    counter))
```

In a program like this, an identifier no longer stands for a ***value***; instead, an identifier stands for a ***variable***

Implementing Variables

Option 1:

```
(define counter 0)
(define (f x)
  (begin
    (set! counter (+ x counter))
    counter))
(f 10)
```

```
⇒ (define counter (box 0))
   (define (f x)
     (begin
       (set-box! counter (+ (unbox x)
                             (unbox counter)))
       (unbox counter)))
   (f (box 10)))
```

Option 2:

- Essentially the same, but hide the boxes in the interpreter

BMCFAE = BCFAE + variables

```
<BMCFAE> ::= <num>
| { + <BMCFAE> <BMCFAE> }
| { - <BMCFAE> <BMCFAE> }
| <id>
| { fun { <id> } <BMCFAE> }
| { <BMCFAE> <BMCFAE> }
| { if0 <BMCFAE> <BMCFAE> <BMCFAE> }
| { newbox <BMCFAE> }
| { setbox <BMCFAE> <BMCFAE> }
| { openbox <BMCFAE> }
| { seqn <BMCFAE> <BMCFAE> }
| { set <id> <BMCFAE> }
```



Implementing Variables

```
(define-type SubCache  
  [mtSub]  
  [aSub (name symbol?)  
        (address integer?)  
        (sc SubCache?)])
```

Implementing Variables

```
; interp : BCFAE SubCache Store -> Value*Store
(define (interp expr sc st)
  ...
  [id (name) (v*s (store-lookup (lookup name sc) st)
                    st)]
  ...)
```

Implementing Variables

```
; interp : BCFAE SubCache Store -> Value*Store
(define (interp expr sc st)
  ...
  [app (fun-expr arg-expr)
    (interp-two fun-expr arg-expr sc st
      (lambda (fun-val arg-val st)
        (local [(define a (malloc st))]
          (interp (closureV-body fun-val)
            (aSub name
              a
              (closureV-sc fun-val))
            (aSto a
              arg-val
              st))))))]
  ...)
```

Implementing Variables

```
; interp : BCFAE SubCache Store -> Value*Store
(define (interp expr sc st)
  ...
  [set (id val-expr)
    (local [(define a (lookup id sc))]
      (type-case Store*Value (interp val-expr sc st)
        [v*s (val st)
          (v*s val
            (aSto a
              val
              st))])))]
  ...)
```

Variables and Function Calls

```
(define (swap x y)
  (local [(define z y)]
    (set! y x)
    (set! x z)))
```

```
(local [(define a 10)
        (define b 20)]
  (begin
    (swap a b)
    a))
```

Result is **10**; assignment in **swap** cannot affect **a**

Call-by-Reference

What if we wanted `swap` to change `a`?

```
(define (swap x y)
  (local [(define z y)]
    (set! y x)
    (set! x z)))

(local [(define a 10)
       (define b 20)]
  (begin
    (swap a b)
    a))

⇒

(define (swap x y)
  (local [(define z (box (unbox y)))]
    (set-box! y (unbox x))
    (set-box! x (unbox z))))

(local [(define a (box 10))
       (define b (box 20))]
  (begin
    ; (swap (box (unbox a))
    ;      (box (unbox b)))
    (swap a b)
    (unbox a)))
```

This is called ***call-by-reference***, as opposed to ***call-by-value***

Terminology alert: this “call-by-value” is orthogonal to the use in “call-by-value” vs. “call-by-name”

Implementing Call-by-Reference

```
; interp : BCFAE SubCache Store -> Value*Store
(define (interp expr sc st)
  ...
  [app (fun-expr arg-expr)
    (if (id? arg-expr)
      ; call-by-ref handling for id arg:
      (type-case Value*Store (interp fun-expr sc st)
        [v*s (fun-val st)
          (local [(define a
                    (lookup (id-name arg-expr) sc))]
                    (interp (closureV-body fun-val)
                          (aSub name
                              a
                              (closureV-sc fun-val))
                          st)))]
        ; as before:
        ...)]
    ...)
```