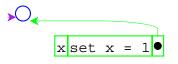
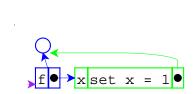
What is the result of this program?

Is it 0 or 1?

2

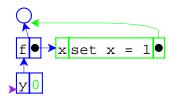
4.

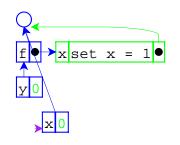


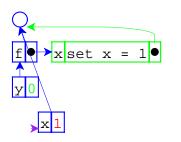


1

3







f • xset x = 1 • y 0 x 1

}

8.

So the answer is 0.

```
void f(int x) {
    x = 1;
}
int main() {
    int y = 0;
    f(y);
    return y;
}
```

```
The result above is 0, too.
```

```
void f(int& x) {
   x = 1;
}
int main() {
   int y = 0;
   f(y);
   return y;
}
```

But the result above is 1.

```
,
```

11

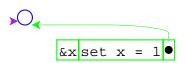
```
void f(int& x) {
    x = 1;
}
int main() {
    int y = 0;
    f(y);
    return y;
}
```

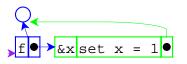
This example shows call-by-reference.

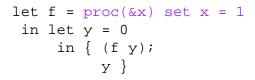
The previous example showed **call-by-value**.

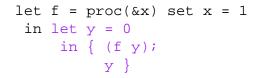
Adding call-by-reference parameters to our language.

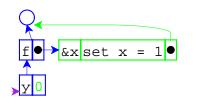
10

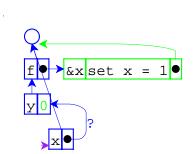




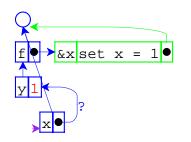


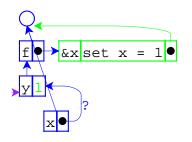






The pointer from one environment frame to another is questionable, because frames are supposed to point to values.





What changes in the interpreter?

}

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20

17

19

Same as before:

- Expressed values: Number + Proc
- **Denoted values:** Ref(Expressed Value)

Same as before:

- Expressed values: Number + Proc
- **Denoted values:** Ref(Expressed Value)

The difference is that application doesn't always create a new location for a new variable binding.

=> Separate location creation from environment extension



The old way



The new way

let x = 10 y = 12 in +(x,y)

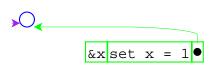
```
•
```

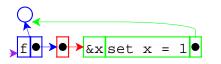
21

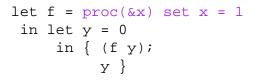
22

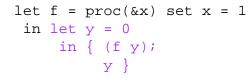
24

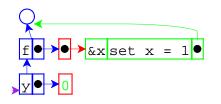
Do the previous evaluation the new way...

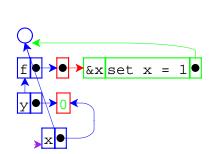




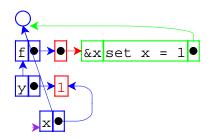


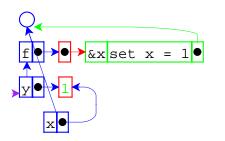






This time, the new environment frame points to a location box, which is consistent with other frames.

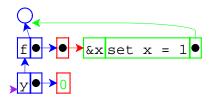


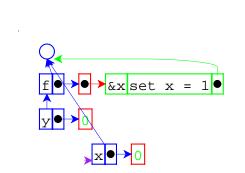


}

30

32





29

31

If call-by-reference argument is not a variable...

... always create a location.

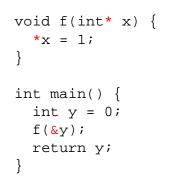
Interpreter changes (starting with pre-letrec version):

- Add call-by-reference arguments (indicated by &). • New var type, with cbv-var and cbr-var
- Create explicit locations for variables. ° location, location-val, location-set!
- Change variable lookup to deference locations.
- Change set to work on locations.
- Change eval-rands and apply-proc. <sup>O</sup> make-var-location helper proc

```
void f(int* x) {
    *x = 1;
}
int main() {
    int y = 0;
    f(&y);
    return y;
}
```

33

35

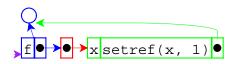


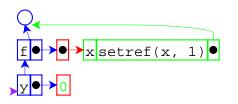
This is back to **call-by-value**, but with a reference as a value.

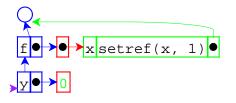
To study this form of call, we can add explicit references to our language, too.

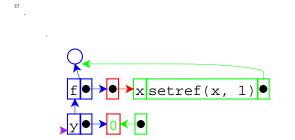
```
xsetref(x, 1)●
```

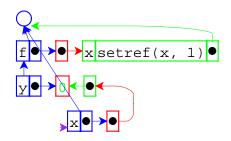
34

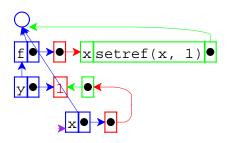


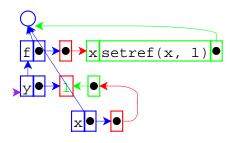












Revised language:

41

}

43

• **Expressed vals:** Number + Proc + Ref(Expressed Val)

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• **Denoted vals:** Ref(Expressed Val)

Interpreter changes:

- Add reference values.
- Add ref form and setref primitive.