What is the result of this program?

```
let f = proc(x) set x = 1
    in let y = 0
        in { (f y);
```

Is it 0 or $1 ?$
let $f=\operatorname{proc}(x)$ set $x=1$
in let $y=0$
in \{ (f y);
y \}


```
let \(f=\operatorname{proc}(x)\) set \(x=1\)
    in let \(y=0\)
        in \(\{(f y) ;\)
        Y \}
```

```
let f = proc(x) set x = 1
    in let }y=
    in { (f y);
        y }
```

```
```

let f = proc(x) set x = 1

```
```

let f = proc(x) set x = 1
in let }y=
in let }y=
in { (f y);
in { (f y);
Y }

```
```

            Y }
    ```
```

```
let f = proc(x) set x = 1
    in let }\textrm{y}=
        in { (f y);
        Y }
```

in $\{(f y) ;$

```
in let \(y=0\)
y \}
```



```
let \(f=\operatorname{proc}(x)\) set \(x=1\)
    in let \(y=0\)
        in \{ (f y);
            Y \}
```


y 0
$\times 2$


```
let f = proc(x) set x = 1
    in let }y=
        in { (f y);
        Y
```

So the answer is 0 .

```
void f(int x) {
    x = 1;
}
int main() {
    int }y=0
    f(y);
    return y;
}
```

The result above is 0 , too.

```
void f(int& x) {
    x = 1;
}
int main() {
    int y = 0;
    f(y);
    return y;
}
```

This example shows call-by-reference.
The previous example showed call-by-value.

```
void f(int& x) {
    x = 1;
}
int main() {
    int y = 0;
    f(y);
    return y;
}
But the result above is 1 .
```

```
let f = proc(&x) set x = 1
    in let y = 0
        in { (f y);
        Y }
```

Adding call-by-reference parameters to our language.


```
let f = proc(&x) set x = 1
    in let }\textrm{y}=
        in { (f y);
        y }
```



```
let f = proc(&x) set x = 1
```

    in let \(y=0\)
        in \(\{\) (fy);
        y \}
    ```
let f = proc(&x) set x = 1
```

let f = proc(\&x) set x = 1
in let }y=
in let }y=
in { (f y);
in { (f y);
y }

```
        y }
```



```
let f = proc(&x) set x = 1
    in let }y=
        in { (f y);
        y }
```

The pointer from one environment frame to another is questionable, because frames are supposed to point to values.


```
let f = proc(&x) set x = 1
    in let }\textrm{y}=
    in { (f y);
        y }
```

```
let f = proc(&x) set x = 1
    in let }y=
        in { (f y);
        y
```


## Same as before:

- Expressed values: Number + Proc
- Denoted values: Ref(Expressed Value)


## Same as before:

- Expressed values: Number + Proc
- Denoted values: Ref(Expressed Value)

The difference is that application doesn't always create a new location for a new variable binding.
=> Separate location creation from environment extension

```
let x = 10
    y = 12
in +(x,y)
\[
\begin{array}{r}
\text { let } \mathrm{x}=10 \\
\mathrm{y}=12 \\
\text { in }+(\mathrm{x}, \mathrm{y})
\end{array}
\]
```

The new way

```
let x = 10
    y = 12
in +(x,y)
\[
\begin{array}{r}
y=12 \\
\text { in }+(x, y)
\end{array}
\]
```

The old way



```
let f = proc(&x) set x = 1
    in let }y=
        in { (f y);
        y }
```

Do the previous evaluation the new way...
let $\mathrm{f}=\operatorname{proc}(\& \mathrm{x})$ set $\mathrm{x}=1$
in let $y=0$
in $\{(f y) ;$
Y $\}$


```
let f = proc(&x) set x = 1
    in let }y=
        in { (f y);
        y }
```

```
```

let f = proc(\&x) set x = 1

```
```

let f = proc(\&x) set x = 1
in let }y=
in let }y=
in { (f y);
in { (f y);
y }

```
```

        y }
    ```
```



```
let f = proc(&x) set x = 1
    in let }y=
        in { (f y);
            y }
```

This time, the new environment frame points to a location box, which is consistent with other frames.

let $\mathrm{f}=\operatorname{proc}(\& x)$ set $\mathrm{x}=1$
in let $y=0$
in $\{(f y) ;$
y $\}$

```
let f = proc(&x) set x = 1
    in let }\textrm{y}=
        in { (f y);
        y
```



```
let f = proc(&x) set x = 1
    in let y = 0
        in { (f 0);
            y }
```

If call-by-reference argument is not a variable...

```
let f = proc(&x) set x = 1
    in let y = 0
        in { (f 0);
        y }
```

... always create a location.

Interpreter changes (starting with pre-letrec version):

- Add call-by-reference arguments (indicated by \&).
- New var type, with cbv-var and cbr-var
- Create explicit locations for variables.
- location, location-val, location-set!
- Change variable lookup to deference locations.
- Change set to work on locations.
- Change eval-rands and apply-proc.
- make-var-location helper proc

```
void f(int* x) {
    *x = 1;
}
int main() {
    int y = 0;
    f(&y);
    return y;
}
```

This is back to call-by-value, but with a reference as a value.
To study this form of call, we can add explicit references to our language, too.

```
void f(int* x) {
    *x = 1;
}
int main() {
    int y = 0;
    f(&y);
    return y;
}
```


x setref(x, 1)

```
let f = proc(x) setref(x, 1)
    in let y = 0
        in { (f ref(y));
            y }
```

```
let f = proc(x) setref(x, 1)
    in let }y=
        in { (f ref(y));
        y }
```

```
let f = proc(x) setref(x, 1)
    in let }\textrm{y}=
        in { (f ref(y));
        y }
```




```
let f = proc(x) setref(x, 1)
    in let }y=
        in { (f ref(y));
        y }
```

let $f=\operatorname{proc}(x)$ setref( $x, 1)$
in let $y=0$
in $\{$ (f ref(y));
Y \}


```
let f = proc(x) setref(x, 1)
    in let }\textrm{y}=
        in { (f ref(y));
        Y }
```

```
let f = proc(x) setref(x, 1)
    in let }y=
        in { (f ref(y));
        Y }
```

Revised language:

- Expressed vals: Number + Proc + Ref(Expressed Val)
- Denoted vals: Ref(Expressed Val)

Interpreter changes:

- Add reference values.
- Add ref form and setref primitive.

