## Outline

$\#$ Data Definitions and Templates
Syntax and Semantics
Defensive Programming

## Data Definitions

Question 1:
Are both of the following data definitions ok?

```
; A w-grade is either
; - num
; - posn
; - empty
    with ; A posn is
    ; (make-posn num num)
; A z-grade is either
; - num
; - (make-posn num num)
; - empty
```


## Data Definitions

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; A z-grade is either
; - num
; - (make-posn num num)
; - empty
```

Yes.

## Data Definitions

## Question 2:

Do w-grade and $\mathbf{z}$-grade identify the same set of values?

```
; A w-grade is either
; - num
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                with ; A posn is
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```

; A z-grade is either
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; - (make-posn num num)
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## Data Definitions

## Question 2:

Do w-grade and $\mathbf{z}$-grade identify the same set of values?

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; - posn
; - empty
                                with ; A posn is
                                (make-posn num num)
```

; A z-grade is either
; - num
; - (make-posn num num)
; - empty

Yes, every w-grade is a w-grade, and every $\mathbf{z}$-grade is a w-grade

## Data Definitions

## Question 3:

Are w-grade and w-grade the same data definition?

```
; A w-grade is either
; - num
; - posn
; - empty
                with ; A posn is
                                (make-posn num num)
```

; A z-grade is either
; - num
; - (make-posn num num)
; - empty

## Data Definitions

Question 3:
Are w-grade and w-grade the same data definition?

```
; A w-grade is either
; - num
; - posn
; - empty
                with ; A posn is
                                (make-posn num num)
```

; A z-grade is either
; - num
; - (make-posn num num)
; - empty

No, in the sense that they generate different templates

## Data Definitions and Templates

The template depends on the static, textual content of a data definition, only

```
; A w-grade is either
; - num
; - posn
; - empty
; A posl. is
; (make-posn num num)
```

; A z-grade is either
; - num
; - (make-posn num num)
; - empty

```
(define (func-for-w-grade w)
    (cond
        [(number? w) ...]
    [(posn? w) ... (func-for-posn w) ...]
    [(empty? w) ...]))
(define (funcinor-posn p)
    ... (posn-x p) ... (posn-y p) ...)
(define (func-for-z-grade z)
    (cond
        [(number? z) ...]
    [(posn? z) ... (posn-x z) ... (posn-y z) ...]
    [(empty? z) ...]))
```


## Data Definitions and Templates

Why we treat the data definition statically to generate a template:

- Provides well-defined, simple rules for generating a template
- "Dynamic" coverage is difficult in general
- Recall 3520 anecdote: thinking in terms of dynamic coverage $\Rightarrow$ broken programs
- Similar to the way that data choices affect modularity
- Details of modularity are beyond the scope of this class, but we want to build the right instincts


## Outline

Data Definitions and Templates
M Syntax and Semantics
Defensive Programming

## Execution in DrScheme

## Suppose that DrScheme's definition window contains

$$
\begin{aligned}
& \text { (define (f x) } \\
& \quad(/ \mathbf{x} 2)) \\
& (\mathbf{f} 10)
\end{aligned}
$$

What's the result of clicking Execute?

## Execution in DrScheme

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## Execution in DrScheme

## Suppose that DrScheme's definition window contains

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## Execution in DrScheme

Suppose that DrScheme's definition window contains

$$
\begin{aligned}
& \text { (define (f x) } \\
& \left(/ \mathbf{x}^{(10))}\right. \\
& (\mathrm{f} 10)
\end{aligned}
$$

What's the result of clicking Execute?

$$
\text { /: divide by } 0
$$

## Execution in DrScheme

## Suppose that DrScheme's definition window contains

$$
\begin{gathered}
\text { (define (f x) } \\
(/ \mathbf{x} 0))
\end{gathered}
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What's the result of clicking Execute?

## Execution in DrScheme

Suppose that DrScheme's definition window contains

$$
\begin{gathered}
\text { (define (f x) } \\
(/ \mathbf{x} 0))
\end{gathered}
$$

What's the result of clicking Execute ?
Nothing (although f would produce an error if it were used)

## Execution in DrScheme

## Suppose that DrScheme's definition window contains

$$
\begin{gathered}
(\text { define }(\mathbf{f} \mathbf{x}) \\
(/ \mathbf{x}(0)))
\end{gathered}
$$

What's the result of clicking Execute ?

## Execution in DrScheme

Suppose that DrScheme's definition window contains

$$
\begin{gathered}
(\text { define }(\mathbf{f} \mathbf{x}) \\
(/ \mathbf{x}(0)))
\end{gathered}
$$

What's the result of clicking Execute ?

```
expected a name after an open parenthesis,
    found a number - even without using f
```


## Execution in DrScheme

## Suppose that DrScheme's definition window contains

$$
\begin{aligned}
& \text { (define (f x) } \\
& (\text { cond } x))
\end{aligned}
$$

What's the result of clicking Execute?

## Execution in DrScheme

Suppose that DrScheme's definition window contains

$$
\begin{gathered}
(\text { define }(f x) \\
(\text { cond } x))
\end{gathered}
$$

What's the result of clicking Execute ?
> cond: expected a question--answer clause-even without using $\mathbf{f}$

## Execution in DrScheme

Suppose that DrScheme's definition window contains

$$
\begin{aligned}
& \text { (define ( } \mathrm{f} \text { ) } \\
& \text { (cond }
\end{aligned}
$$

[false x]))

What's the result of clicking Execute ?

## Execution in DrScheme

Suppose that DrScheme's definition window contains

$$
\begin{aligned}
& \text { (define ( } \mathrm{f} \text { ) } \\
& \text { (cond }
\end{aligned}
$$

[false x]))

What's the result of clicking Execute ?
Nothing

## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
(define (f x)
    (cond
        [false x]))
(f 10)
```

What's the result of clicking Execute ?

## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
(define (f x)
    (cond
    [false x]))
(f 10)
```

What's the result of clicking Execute?
cond: all questions were false

## Errors in DrScheme

## DrScheme complains about a function body

sometimes before the function is used
sometimes only when the function is called
Why?

## Errors in DrScheme

DrScheme complains about a function body
sometimes before the function is used

- sometimes only when the function is called

Why?
Because some errors are syntax errors and some errors are run-time errors

## Syntax Errors

A syntax error is like a question that isn't a well-formed sentence

- $\mathbf{f}(\mathbf{x})=\mathbf{x}+0$
- DrScheme doesn't understand this notation, just like...
- "Parlez vous Francais?"
- English-only speaker doesn't understand this notation


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- English-only speaker doesn't understand this notation
- (define (f x) (/ x (0)))
- Parens around a zero make no sense to DrScheme, just like...
- "Does rain dog cat?"
- Not enough verbs for this to make sense in English


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- (define (f x) (/ x (0)))
- Parens around a zero make no sense to DrScheme, just like...
- "Does rain dog cat?"
- Not enough verbs for this to make sense in English

When DrScheme sees a syntax error, it refuses to evaluate

## Run-Time Errors

A run-time error is like a well-formed question with no answer

- (/ 12 0)
- A clear request to DrScheme, but no answer, just like...
- "Why are you wearing a green hat?"

There's no answer if l'm wearing a blue hat

## Run-Time Errors

A run-time error is like a well-formed question with no answer

- (/ 12 0)
- A clear request to DrScheme, but no answer, just like...
- "Why are you wearing a green hat?"
- There's no answer if I'm wearing a blue hat
- (cond [false 10])
- There's no reasonable choice for DrScheme, just like...
- "If you can't understand me, what’s your name?"
- No one who understands the question should answer


## Run-Time Errors

A run-time error is like a well-formed question with no answer

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- There's no reasonable choice for DrScheme, just like...
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DrScheme evaluates around run-time errors until forced to answer

## The Difference between DrScheme and English

In a (good) programming language, all errors are well-defined, and the rules are relatively simple

- DrScheme has a simple, well-defined grammar, and deviations from the grammar are syntax errors
- The reduction rules for each construct and primitive operation are well-defined, producing either a value or an error


## Beginner Scheme Grammar

A <var> is a name, a <con> is a constant, and a <prm> is an operator name

A <defn> is one of
(define (<var> <var> ... <var>) <exp>)
(define <var> <exp>)
(define-struct <var> (<var> ... <var>))
A <exp> is one of

```
<var>
<con>
(<prm> <exp> ... <exp>)
(<var> <exp> . . . <exp>)
(cond [<exp> <exp>] . . [<exp> <exp>])
(cond [<exp> <exp>] ... [else <exp>])
(and <exp> . . . <exp>)
(or <exp> ... <exp>)
```


## Evaluation Rules: and/or

```
(and true) }->\mathrm{ true
(and true question ... question)
    ->(and question ... question)
(and false question ... question) }->\mathrm{ false
(or false) }->\mathrm{ false
(or false question ... question)
    (or question ... question)
(or true question ... question) }->\mathrm{ true
```


## Evaluation Rules: and/or

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(or false) }->\mathrm{ false
(or false question ... question)
    (or question ... question)
(or true question ... question) }->\mathrm{ true
```

Note that

$$
\text { (and } 7 \text { false) }
$$

fits the grammar, but has no matching evaluation rule, so it produces a run-time error

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## Execution in DrScheme

Suppose that DrScheme's definition window contains

```
; f : num -> num
(define (f x)
    (+ x 2))
(f 'apple)
```

What's the result of clicking Execute?

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What's the result of clicking Execute ?
+: expects a <number>, given 'apple

## Execution in DrScheme

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; f : num -> num
(define (f x)
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```

What's the result of clicking Execute ?
+: expects a <number>, given 'apple

But this is really a contract violation at the call to $\mathbf{f}$
The implementor of $\mathbf{f}$ might want to clarify that this error is someone else's fault, not a bug in $\mathbf{f}$

## Defensive Programming

```
; f : num -> num
(define (real-f x)
    (+ x 2))
(define (f x)
    (cond
        [(number? x) (real-f x)]
        [else (error 'f "not a number")]))
(f 'apple)
```


## Defensive Programming

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; f : num -> num
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                f: not a number
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## Defensive Programming

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The error function triggers a run-time error

## Defensive Programming

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; f : num -> num
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(f 'apple)
f: not a number
```

The error function triggers a run-time error
You don't have to program defensively in this course, but it sometimes helps to defend against your own mistakes!

