

L20 -- Markov Chain  
[Jeff Phillips - Utah - Data Mining]

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Graph  $G = (E, V)$

$V =$  vertices  $\{a, b, c, d, e, f, g, h\}$

$E =$  edges  $\{(a, b), (a, c), (a, d), (b, d), (c, d), (c, e), (e, f), (e, g), (f, g), (f, h)\}$   
unordered pairs

Draw graph:

	a	b	c	d	e	f	g	h
a	0	1	1	1	0	0	0	0
b	1	0	0	1	0	0	0	0
c	1	0	0	1	1	0	0	0
d	1	1	1	0	0	0	0	0
e	0	0	1	0	0	1	1	0
f	0	0	0	0	1	0	1	1
g	0	0	0	0	1	1	0	0
h	0	0	0	0	0	1	0	0

**\*\*adjacency matrix\*\***

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Each  $v$  in  $V$  is a state.

If at  $b$ , represent state as

$q = [0 \ 1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0]^T$

Can "think of" fractional state

$q = [1/2 \ 0 \ 0 \ 1/2 \ 0 \ 0 \ 0 \ 0]^T$

1/2 at  $a$  and 1/2 at  $d$

probability of being in each state:

each  $q[i] \geq 0$  and  $\sum_i q[i] = 1$

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Transition matrix  $P =$  normalized adjacency matrix

	a	b	c	d	e	f	g	h
a	0	1/2	1/3	1/3	0	0	0	0
b	1/3	0	0	1/3	0	0	0	0
c	1/3	0	0	1/3	1/3	0	0	0
d	1/3	1/2	1/3	0	0	0	0	0
e	0	0	1/3	0	0	1/3	1/2	0
f	0	0	0	0	1/3	0	1/2	1
g	0	0	0	0	1/3	1/3	0	0
h	0	0	0	0	0	1/3	0	0

then given a state  $q$ , we can "transition" to the next state by

$$q_1 = P * q$$

This one "step" of a "Markov Chain".

"Markov" means that each state only depends on previous state.

next step

$$q_2 = P * q_1 \quad \text{or}$$

$$= P * P * q \quad \text{or}$$

$$= P^2 * q$$

$$q_n = P^n * q$$

where  $P^n = P * P * P * \dots n \text{ times} \dots * P$

Can think of as a randomized random walk.

+ start state  $q = q_0$ .

+ each step, takes one path at random

+  $q_n$  is probability distribution of state after  $n$  steps

+ thus each column of  $P^n$  positive, sums to 1 for all  $n$

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Markov Chain is **ergodic** if

exists some  $t$  such that for all  $n \geq t$  then

each entry in  $P^n$  is positive.

--> for any  $q$ , then

$$q_n = P^n q$$

is positive in all elements

--> after  $t$  steps, always have *some* probability of being anywhere.

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When is a chain not ergodic?

+ cyclic

$$P = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 \end{bmatrix}$$

always alternates states in even/odd states

--> can be larger and more irregular, uncommon in practice

+ has absorbing + transient states

$P$  based on *directed* graph

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 \\ 1/2 & 0 & 1/2 & 1 \\ 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1/2 & 0 & 1/2 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 1/2 & 1/2 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 0 & 0 & 0 \end{bmatrix}$$

state  $d$  always goes to  $b$ , but can never return to  $d$ .

also...

$$P = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 \\ 1/2 & 0 & 1/2 & 1/2 \\ 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 0 & 1/2 \end{bmatrix}$$

may stay at d (w.p. 1/2) but state "seeps" from d to b (and then a,c)

(a,b,c) = absorbing, d = transient

+ not connected

$$P = \begin{bmatrix} 1/2 & 1/2 & 0 & 0 \\ 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 2/3 & 1/2 \\ 0 & 0 & 1/3 & 1/2 \end{bmatrix}$$

(a,b) cannot reach (c,d) and vice-versa

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Consider an ergodic Markov Chain (P,q)

**\*\*AMAZING\*\*** property

let  $P^{*} = P^n$  as  $n \rightarrow \infty$

then  $q_{*} = P^{*} q$

is **\*\*NOT\*\*** dependent on q

--> That is, for all starting states q, the final state is  $q_{*}$

--> as we do a random walk, we will eventually be in the same expected state.

Note that  $q_{*} = P^{*} q = P^{*+1} q$

so  $q_{*} = P q_{*}$

--> If state distribution is initially  $q_{*}$ , then already in final distribution.

$q_{*}$  second eigenvector of P

second eigenvalue determines rate of convergence

--> smaller <-> faster convergence

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Metropolis Algorithm (MCMC)

Metropolis, Rosenbluth, Rosenbluth, Teller, Teller 1953

(Boltzman dist, Manhattan project)

Hastings 1970

(more general)

each state v in V has weight associated with it

$$w(v) \quad \sum_{\{v \text{ in } V\}} w(v) = W$$

Want to land in state  $v$  w.p.  $w(v)/W$

-->  $V$  might be very large, and  $W$  unknown.

-->  $V$  can be "continuous"

"probe-only" can only measure  $w(v)$  at any one state

Strategy: design special Markov Chain so  $q_*[v] = w(v)/W$

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 Start  $v_0$  in  $V$  ( $q = [0 \ 0 \ 0 \ \dots \ 1 \ \dots \ 0 \ 0]^T$ )

choose neighbor  $u$  (proportional to  $K(v,u)$ )

if  $(w(u) \geq w(v_i))$  -->  $v_{\{i+1\}} = u$

else w.p.  $w(u)/w(v)$  -->  $v_{\{i+1\}} = u$

else -->  $v_{\{i+1\}} = v_i$

if ergodic:

there exists some  $t$  s.t. for  $i \geq t$

$$\Pr[v_i = v] = w(v)/W$$

NOTE: not in limit, but for some finite  $t$  (even for continuous)  $V$   
 through AMAZING "coupling from past"

But  $t$  is hard to find.

Often goal is to create many samples:

formal: run for  $t_+$  steps, take sample, ...

run for another  $t_+$  steps, take sample, ... repeat

in practice: run for 1000 steps (burn in),  
 take next 5000 steps as random samples

has "auto-correlation" but eventually more time efficient than  $tN$  steps for  $N$   
 samples

and  $t$  unknown.

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"inherently sequential" makes very hard to parallelize

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Applies even if  $V$  is continuous