

An Introduction to x86 ASM

Malware Analysis Seminar

Meeting 1

Cody Cutler, Anton Burtsev

Registers

- General purpose
 - EAX, EBX, ECX, EDX
 - ESI, EDI (index registers, but used as general in 32-bit protected mode)
- Stack
 - EBP, ESP
- Instruction pointer
 - EIP
- Flags
 - EFLAGS
- Segment
 - CS, DS, SS, ES, FS, GS

Syntax

- General form:
 - *mnemonic operand(s)*
`movl %eax, %ebx`
- Operands (0-3) may refer:
 - Registers
 - Memory
 - Immediate

Syntax (contd.)

- Register naming (AT&T is UNIX default)
 - **AT&T:** %eax
 - **Intel:** eax
- Source/Destination Ordering:
 - Load EBX with the value of EAX
 - **AT&T:** movl %eax, %ebx
 - **Intel:** mov ebx, eax

Constant value/immediate value format

- Load EAX with the address of the "C" variable boo
 - **AT&T:** `movl $_boo, %eax`
 - **Intel:** `mov eax, _boo`
 - Note that “_” works for static (global) variables only
- Now let's load ebx with 0xd00d:
 - **AT&T:** `movl $0xd00d, %ebx`
 - **Intel:** `mov ebx, d00dh`

Operator size specification

- You don't want make GAS to guess this wrong
 - **AT&T:** `movw %ax, %bx`
 - **Intel:** `mov bx, ax`

Referencing memory

- 32bit protected mode addressing
 - **AT&T:** `immed32(basepointer,indexpointer,indexscale)`
 - **Intel:** `[basepointer + indexpointer*indexscale + immed32]`
- A global C variable
 - **AT&T:** `_booga`
 - **Intel:** `[_booga]`
- Addressing what a register points to:
 - **AT&T:** `(%eax)`
 - **Intel:** `[eax]`

Referencing memory (contd.)

- Addressing a variable offset by a value in a register:
 - **AT&T:** `_variable(%eax)`
 - **Intel:** `[eax + _variable]`
- Addressing a value in an array of integers (scaling by 4):
 - **AT&T:** `_array(,%eax,4)`
 - **Intel:** `[eax*4 + array]`
- Offsets with immediate value
 - **C code:** `*(p+1)` where `p` is a `char *`
 - **AT&T:** `1(%eax)` where `eax` has the value of `p`
 - **Intel:** `[eax + 1]`

Referencing memory (contd.)

- Addressing a particular char in an array of 8-character records
 - **EAX** holds the number of the record desired.
 - **EBX** has the wanted char's offset within the record.
 - **AT&T:** `_array(%ebx,%eax,8)`
 - **Intel:** `[ebx + eax*8 + _array]`

Arithmetic

- Integers:
 - Two's compliment:
 - Reverse bits, then add one (throw away carry)
 - Original value: 00111000 (+56)
 - Reverse bits: 11000111
 - Add 1: 11001000 (-56)
- Rules of arithmetic are preserved

$$\begin{array}{r} 002C \\ + \underline{FFFF} \\ \hline 002B \end{array} \qquad \begin{array}{r} 44 \\ + \underline{(-1)} \\ \hline 43 \end{array}$$

Carry and overflow

- Overflow
 - Set if the true result of the operation is too big to fit into the destination for signed arithmetic.
- Carry
 - Set if there is a carry in the msb of an addition or a borrow in the msb of a subtraction.
 - Can be used to detect overflow for unsigned arithmetic.

Extended precision arithmetic

- ADC

operand1 = operand1 + carry flag + operand2

- SBB

operand1 = operand1 - carry flag - operand2

- Sum of 64-bit integers in **EDX:EAX** and **EBX:ECX**

`add eax, ecx ; add lower 32-bits`

`adc edx, ebx ; add upper 32-bits and carry`

Control structures

- Control structures decide what to do based on comparisons of data
- **CMP** instruction
 - subtract operands
 - set EFLAGS
- **EFLAGS** register
 - ZF – zero flag
 - CF – carry flag
 - SF – sign flag

Control structures (contd.)

- **Unsigned:** `cmp vleft, vright <=> vleft - vright`
 - `vleft = vright`: ZF (1), CF (0)
 - `vleft > vright`: ZF (0), CF (0) – no borrow
 - `vleft < vright`: ZF (0), CF (1) – borrow
- **Signed:** `cmp vleft, vright <=> vleft - vright`
 - `vleft = vright`: ZF (1), CF (0)
 - `vleft > vright`: ZF (0), SF = CF
 - `vleft < vright`: ZF (0), SF != CF

Branch instructions

- JMP
 - Short:
 - One byte instruction!
 - But jumps only 128 bytes up or down
 - Near:
 - Jump anywhere in a segment
 - 2-byte displacement: jump 32000 bytes
 - 4-byte displacement: jump anywhere in 32-bit mode
 - Far:
 - Jump across segments

Examples

```
if ( EAX == 0 )      cmp eax, 0 ; set flags (ZF set if eax - 0 = 0)
    EBX = 1;         jz thenblock ; if ZF is set branch to thenblock
else                 mov ebx, 2 ; ELSE part of IF
    EBX = 2;         jmp next ; jump over THEN part of IF
thenblock:
    mov ebx, 1 ; THEN part of IF
next:
```


Comparison instructions

- JE branches if $v_{\text{left}} = v_{\text{right}}$
- JNE branches if $v_{\text{left}} \neq v_{\text{right}}$
- JL, JNGE branches if $v_{\text{left}} < v_{\text{right}}$
- JLE, JNG branches if $v_{\text{left}} \leq v_{\text{right}}$
- JG, JNLE branches if $v_{\text{left}} > v_{\text{right}}$
- JGE, JNL branches if $v_{\text{left}} \geq v_{\text{right}}$

Loops

- LOOP
 - Decrements ECX, if ECX != 0 branches to label
- LOOPE, LOOPZ
 - Decrements ECX (FLAGS register is not modified), if ECX != 0 and ZF = 1, branches
- LOOPNE, LOOPNZ
 - Decrements ECX (FLAGS unchanged), if ECX != 0 and ZF = 0, branches

Loop example

```
sum = 0;  
for ( i=10; i >0; i-- )  
    sum += i;
```

```
mov eax, 0 ; eax is sum  
mov ecx, 10 ; ecx is i
```

```
loop_start:
```

```
    add eax, ecx  
    loop loop_start
```

Stack

- SS
 - Specifies stack segment (usually same as data)
- ESP
 - Contains the address of the data that would be removed from the stack
- PUSH/POP
 - Insert/remove data on the stack
 - Subtract/add 4 to ESP

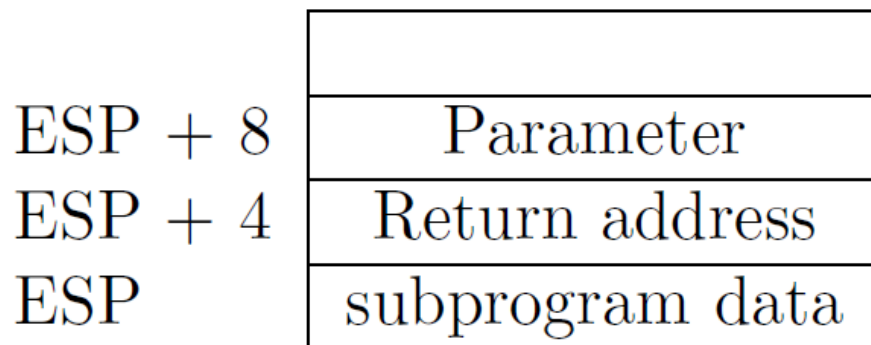
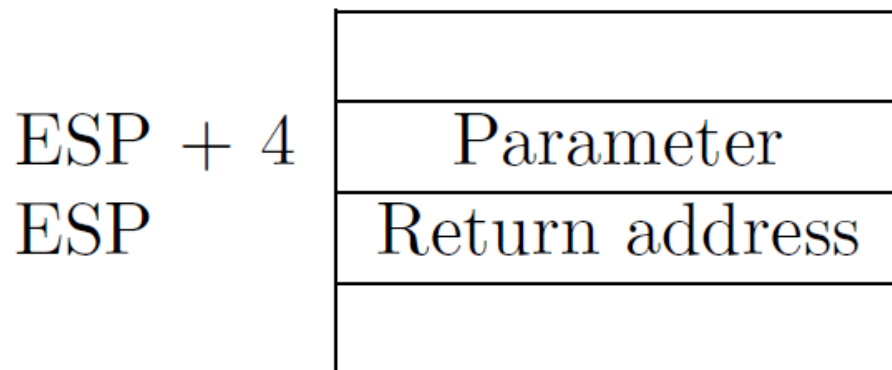
Call/return

- CALL
 - Makes an unconditional jump to a subprogram and pushes the address of the next instruction on the stack
- RET
 - Pops off an address and jumps to that address

Calling conventions

- Goal: reentrant programs
 - Conventions differ from compiler, optimizations, etc.
- Call/return are used for function invocations
- Parameters passed on the stack
 - Pushed onto the stack before the CALL instruction

Stack bottom pointer



Initially parameter is

- [ESP + 4]

Later as the function pushes things on the stack it changes, e.g.

- [ESP + 8]
- Use dedicated register **EBP**

Prologue/epilogue

```
subprogram_label:  
    push ebp        ; save original EBP value on stack  
    mov ebp, esp    ; new EBP = ESP  
; subprogram code  
    pop ebp         ; restore original EBP value  
    ret
```

- Example invocation

```
push dword 1      ; pass 1 as parameter  
call fun  
add esp, 4        ; remove parameter from stack
```


Local variables

- Stored right after the saved EBP value in the stack
- Allocated by subtracting the number of bytes required from ESP

```
subprogram_label:
    push ebp                ; save original EBP value on stack
    mov ebp, esp           ; new EBP = ESP
    sub esp, LOCAL_BYTES ; = # bytes needed by locals
; subprogram code
    mov esp, ebp           ; deallocate locals
    pop ebp                ; restore original EBP value
    ret
```

Enter/leave

- ENTER
 - prologue code
- LEAVE
 - Epilogue

```
subprogram_label:  
    enter LOCAL_BYTES, 0    ; = # bytes needed by locals  
; subprogram code  
    leave  
    ret
```

Examples