

Terminology: Denoted and Expressed Values

- A ***denoted value*** is the meaning of a variable
- An ***expressed value*** is the result of an expression

The set of denoted and expressed values can be different

Terminology: Denoted and Expressed Values

- First-order functions
 - denoted values: numbers and functions
 - expressed values: numbers
- Higher-order functions
 - denoted values: numbers and functions
 - expressed values: numbers and functions

Procedure Expressions: Concrete Syntax

$\langle \text{prog} \rangle ::= \langle \text{expr} \rangle$

$\langle \text{expr} \rangle ::= \text{proc } (\langle \text{id} \rangle^{*(,.)}) \langle \text{expr} \rangle$

$\quad ::= (\langle \text{expr} \rangle \langle \text{expr} \rangle^*)$

**let identity = proc(x) x
in (identity 5)**

Procedure Expressions: Abstract Syntax

<prog> ::= (a-program <expr>)
<expr> ::= (proc-exp (list <id>*) <expr>)
::= (app-exp <expr> (list <expr>*))
<val> ::= <num> | <proc>
<proc> ::= (closure (list <id>*) <expr> <env>)

(a-program
 (let-exp (list 'identity)
 (list (proc-exp (list 'x) (var-exp 'x)))
 (app-exp (var-exp 'identity) (list-exp 5))))

Implementing Procedures

(implementation in DrScheme)

New representation of environments:

```
(define-datatype environment environment?  
  (empty-env-record)  
  (extended-env-record  
    (syms (list-of symbol?))  
    (vals (list-of denval?))  
    (env environment?)))
```

Recursion

Suppose we try to write the **fact** function using only **let**

```
let fact = proc(n) if n then *(n, (fact -(n, 1))) else 1  
in (fact 10)
```

The above doesn't work, because **fact** is not bound in the local function

We'll add **letrec**, but first we'll see how to implement **fact** without it...

Recursion with Let

- **Problem:** **fact** can't see itself
- **Note:** anyone calling **fact** can see **fact**
- **Idea:** have the caller supply **fact** to **fact** (along with a number)

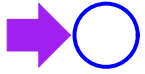
```
let fact = proc(n, f) if n then *(n, (f -(n, 1) f)) else 1
in (fact 10 fact)
```

this works!

What Happened?

- The key insight is delaying some work to the caller
- We can exploit this idea to implement **letrec**, but in a slightly different way
- **letrec** requires a *closure* that refers to itself
- We can delay the actual construction of the closure until it is extracted from the environment

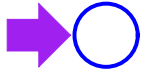
Recursive Environments for Recursive Functions



This isn't going to work

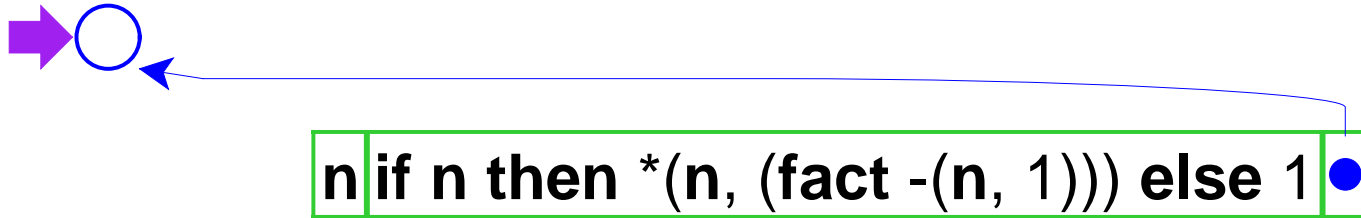
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in (fact 10)
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Recursive Environments for Recursive Functions



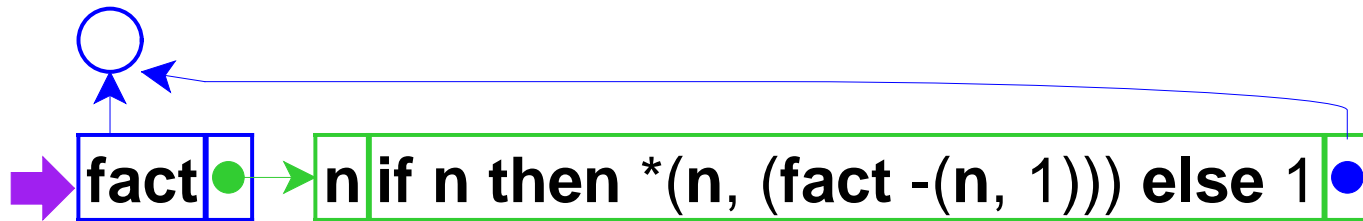
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Recursive Environments for Recursive Functions



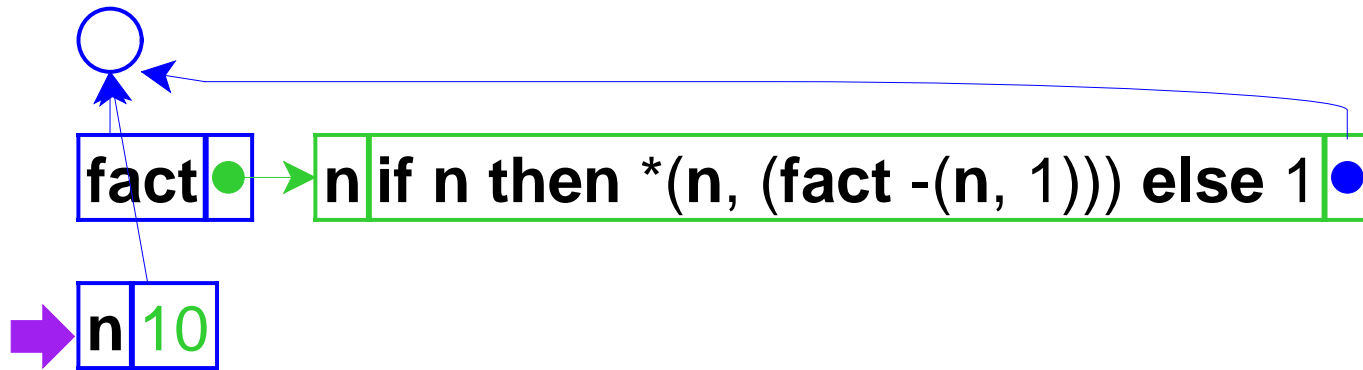
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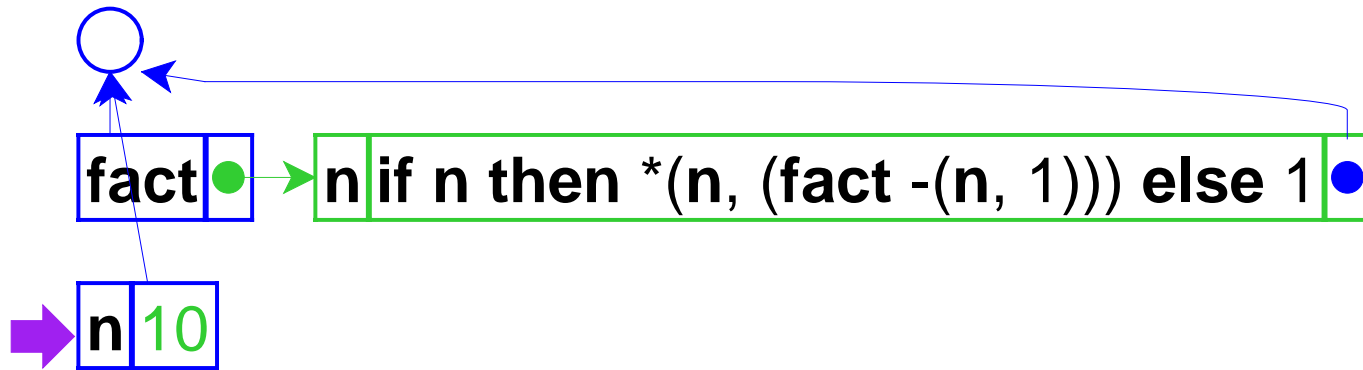
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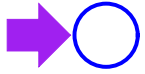
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*No binding for **fact***

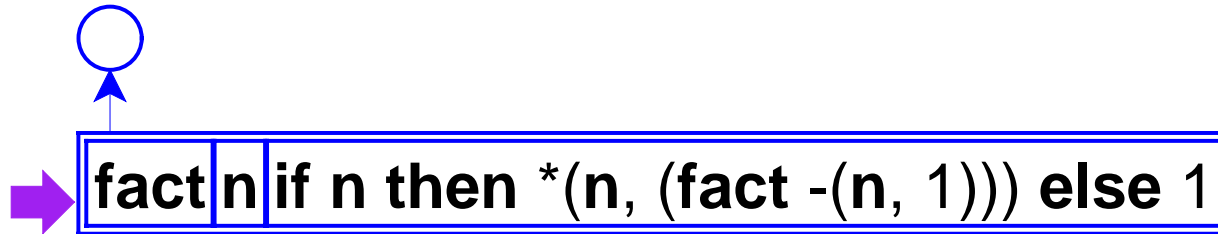
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Recursive Environments for Recursive Functions



```
letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

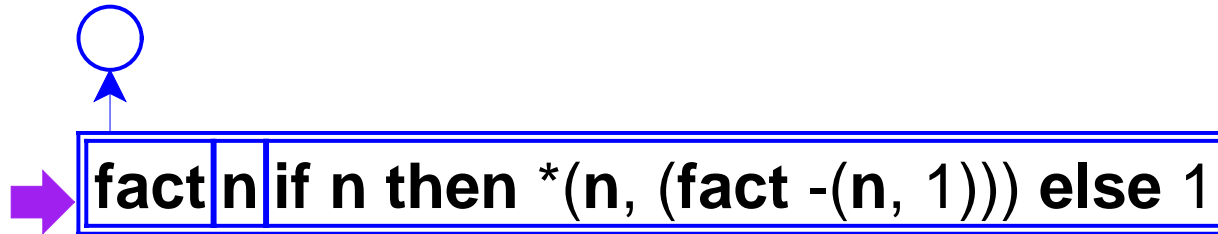
Recursive Environments for Recursive Functions



double box means a recursively extended environment

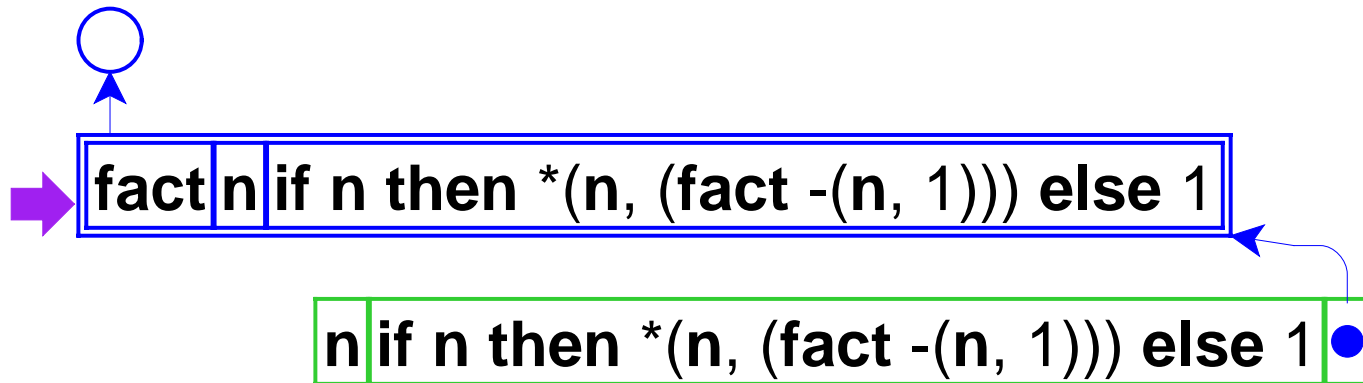
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Recursive Environments for Recursive Functions



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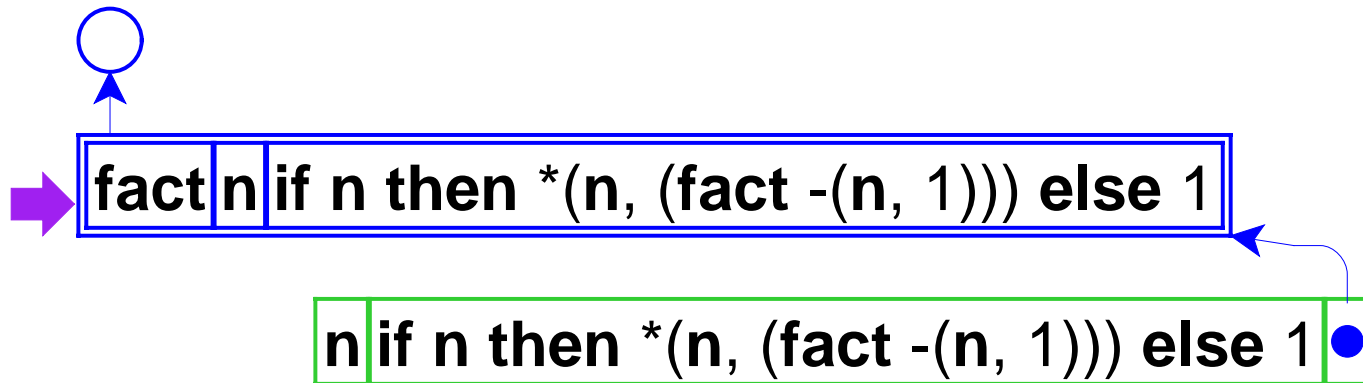
Recursive Environments for Recursive Functions



*every lookup of **fact**
generates a closure*

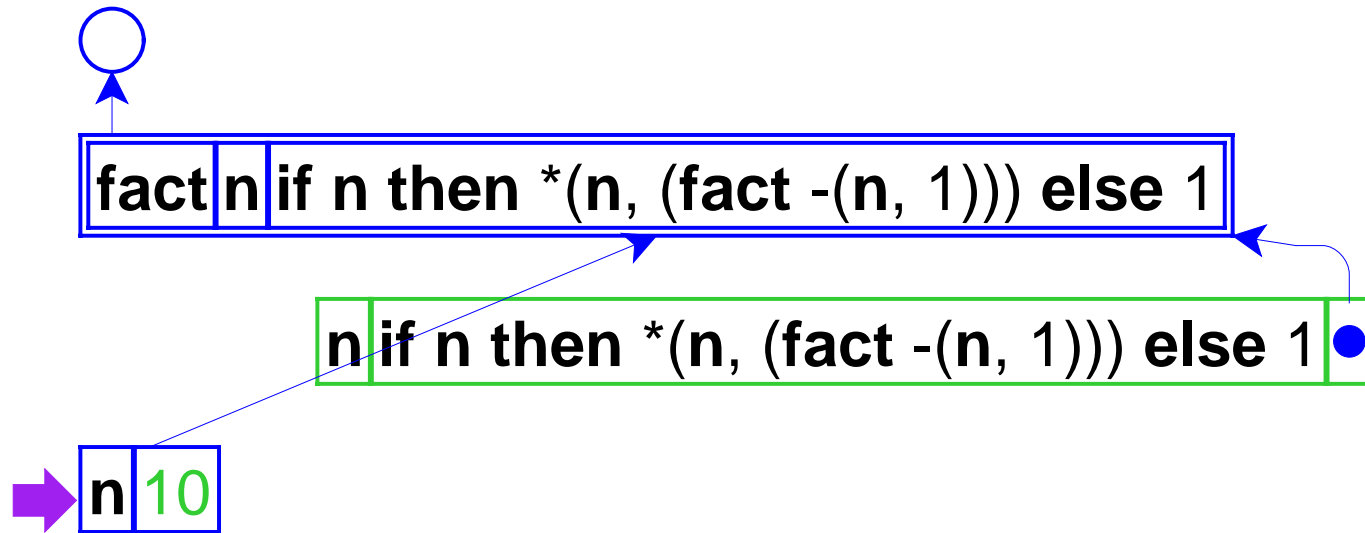
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Recursive Environments for Recursive Functions



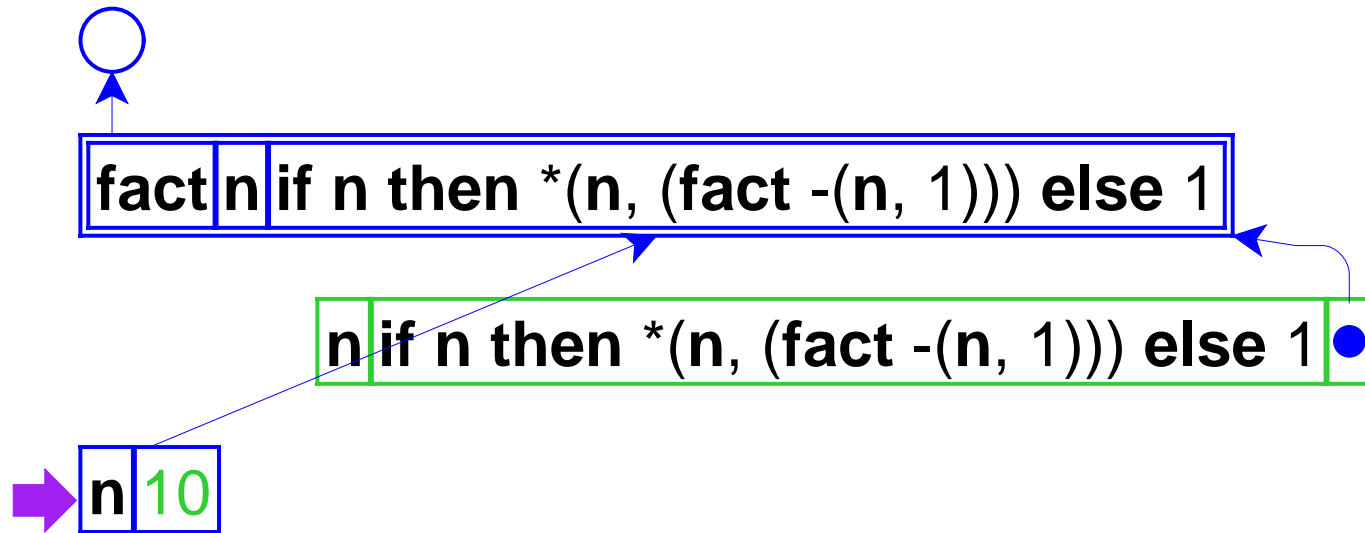
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Recursive Environments for Recursive Functions



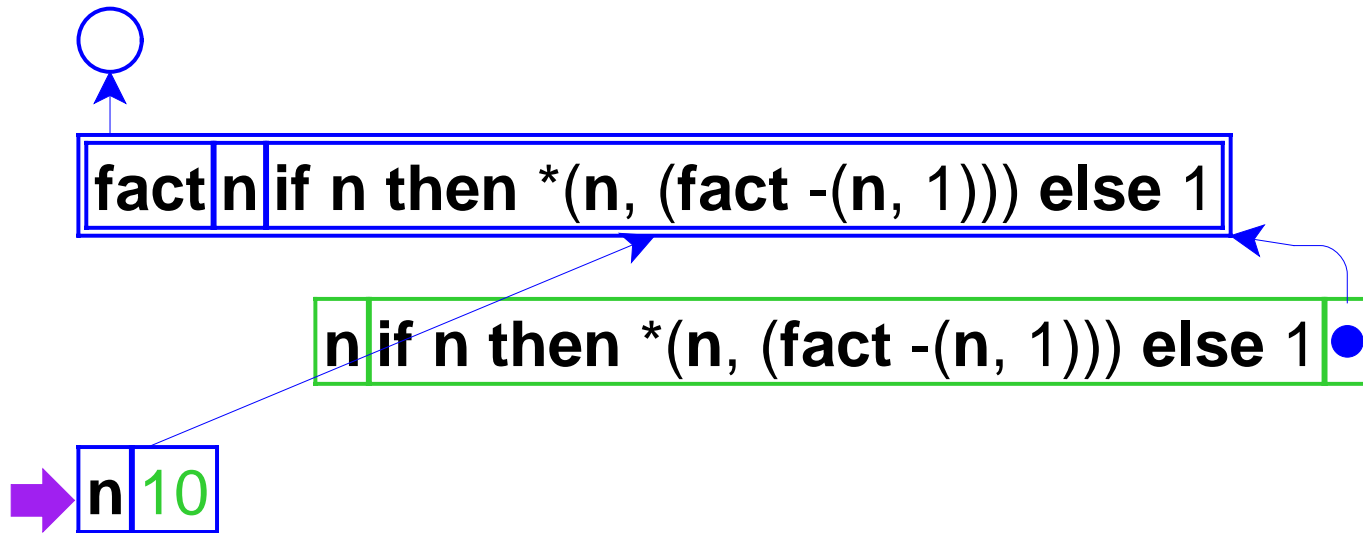
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Recursive Environments for Recursive Functions



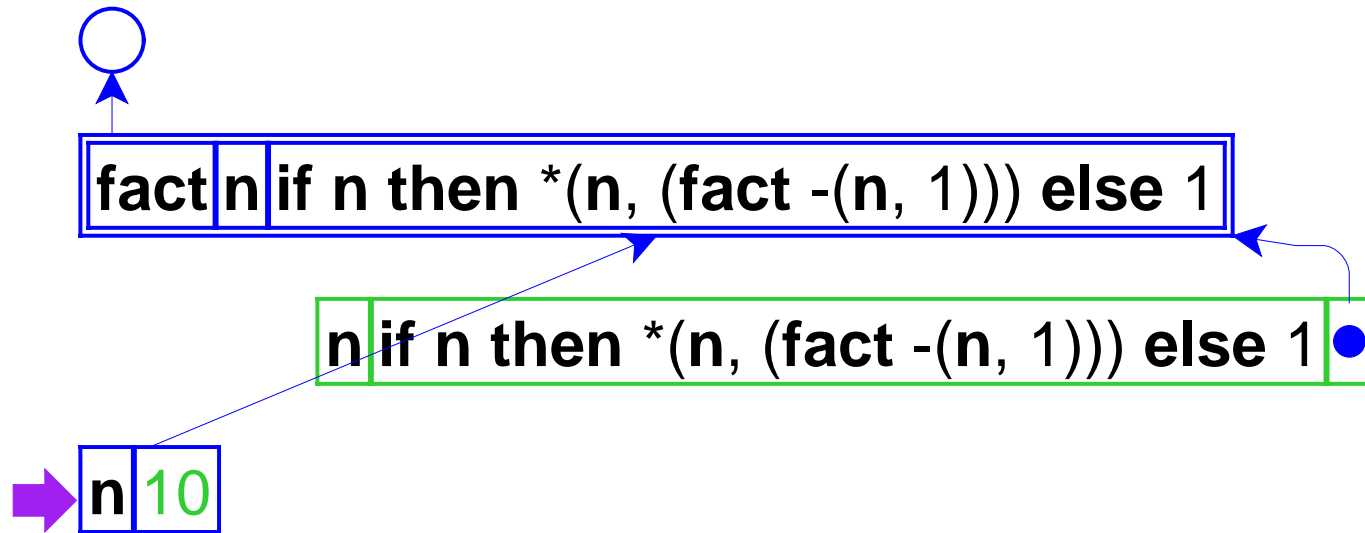
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Recursive Environments for Recursive Functions



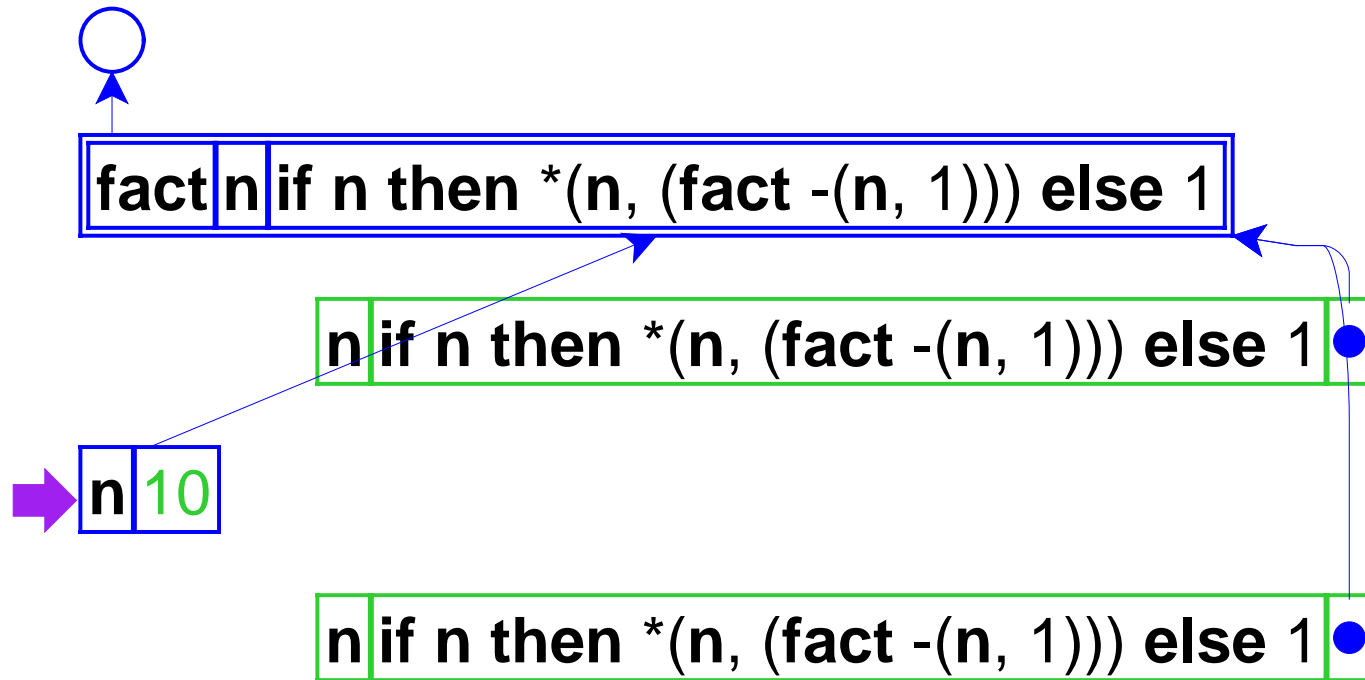
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Recursive Environments for Recursive Functions



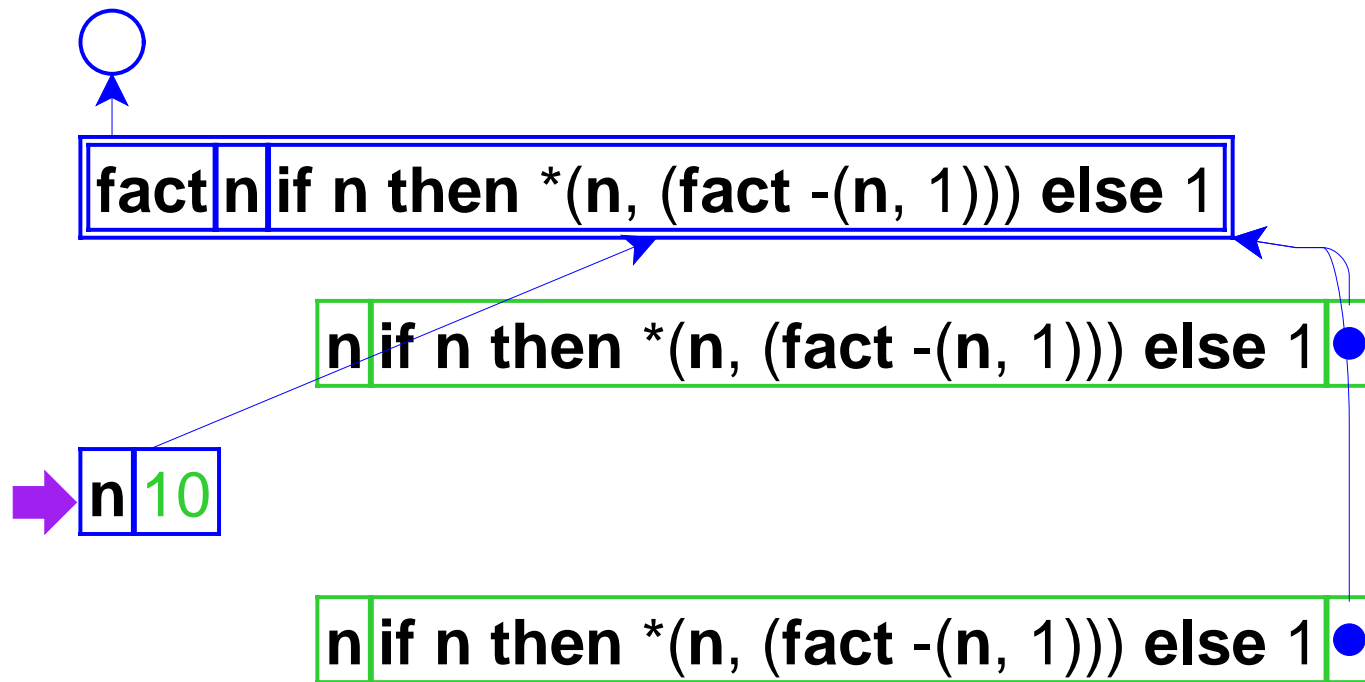
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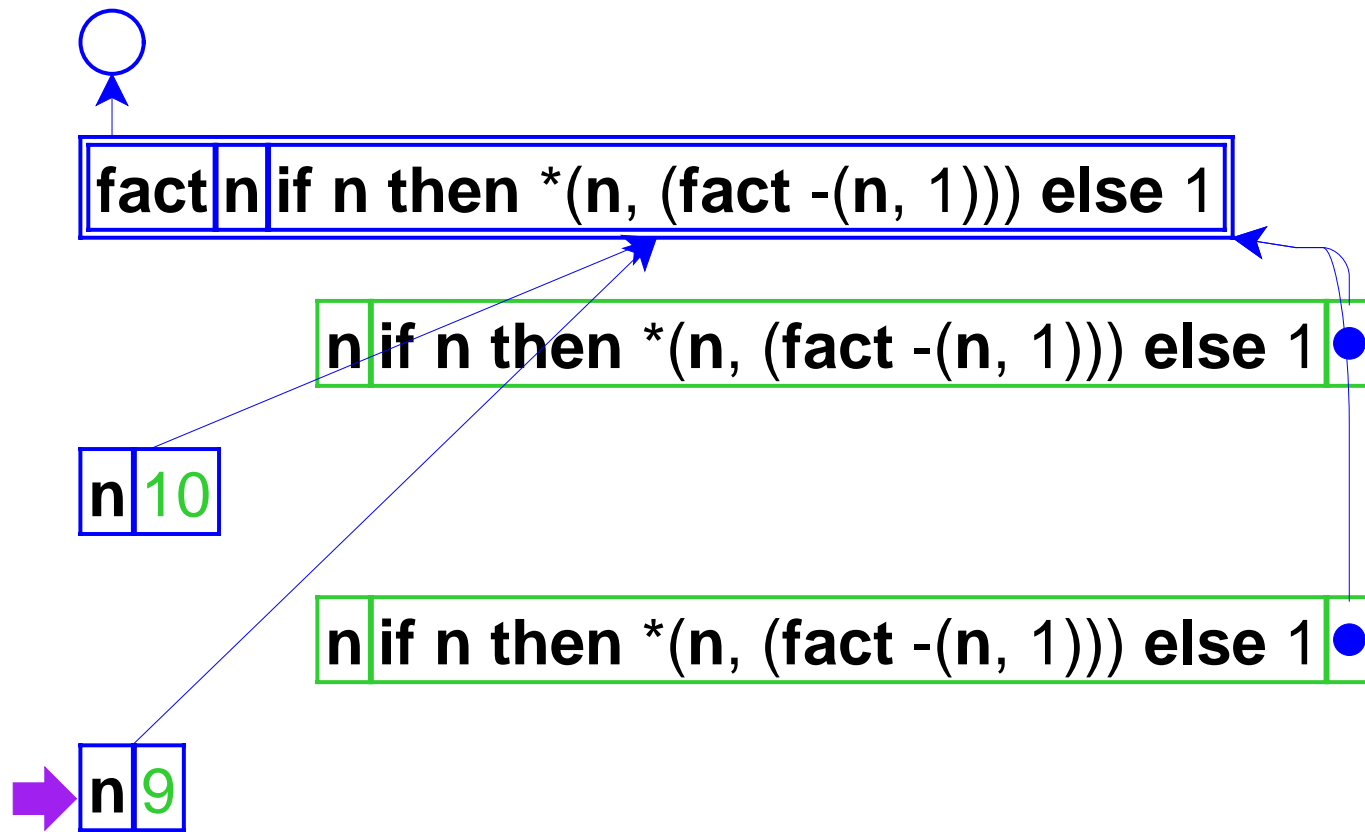
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Recursive Environments for Recursive Functions



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in (fact 10)
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Recursive Environments for Recursive Functions



```
letrec fact = proc(n) if n then *(n, (fact -(n, 1))) else 1
in (fact 10)
```

Implementing Recursively Extended Envs

```
(define-datatype environment environment?
  (empty-env-record)
  (extended-env-record
    (syms (list-of symbol?))
    (vals (list-of denval?))
    (env environment?))
  (recursively-extended-env-record
    (proc-names (list-of symbol?))
    (idss (list-of (list-of symbol?)))
    (bodies (list-of expression?))
    (env environment?)))
```

Implementing letrec

(implement in DrScheme)

Back to Recursion with Let...

- Allowing functions to be values is a powerful idea
- As it turns out, we don't even need **let**!

let $\langle id \rangle_1 = \langle expr \rangle_1 \dots \langle id \rangle_n = \langle expr \rangle_n$ **in** $\langle expr \rangle$

is the same as

(proc($\langle id \rangle_1, \dots \langle id \rangle_n$) $\langle expr \rangle$ $\langle expr \rangle_1 \dots \langle expr \rangle_n$)

Back to Recursion with Let...

- Allowing functions to be values is a powerful idea
- As it turns out, we don't even need **let**!

(**let** ([<id>₁ <expr>₁] ... [<id>_n = <expr>_n]) <expr>)

is the same as

((**lambda** (<id>₁ ... <id>_n) <expr>) <expr>₁ ... <expr>_n)

The Lambda Calculus

- We don't even need functions of multiple arguments...

$$((\mathbf{lambda} \langle id \rangle_1 \dots \langle id \rangle_n \langle expr \rangle) \\ \langle expr \rangle_1 \dots \langle expr \rangle_n)$$

is the same as

$$(((\mathbf{lambda} \langle id \rangle_1) \dots (\mathbf{lambda} \langle id \rangle_n \langle expr \rangle)) \\ \langle expr \rangle_1) \dots \\ \langle expr \rangle_n)$$

Passing multiple arguments one-at-a-time is called ***currying***

The ***lambda calculus*** has only single-argument **lambda** and single-argument function calls, and it's computationally complete