## Assignment in Scheme

So far, we have one form of assignment: vector-set!

$$
\begin{aligned}
& \text { (let ([v (vector } 1: 23)]) \\
& \quad(\text { begin } \\
& \quad(\text { vector-set! } \mathbf{v} 172) \\
& \quad \text { v }) \text { ) } \\
& \rightarrow \rightarrow \\
& \#(1723)
\end{aligned}
$$

## Assignment in Scheme

Scheme actually allows variables to be modified:

```
(let ([x 2])
    (begin
    (set! x 73)
    x))
\longrightarrow
7 3
```

- Don't write Scheme code like that, except for HW6
- But many languages have assignment, and need it


## Assignment in the Book Language

- Add a set expression form:

$$
\text { <expr> }::=\text { set <id> = <expr> }
$$

## Evaluating with Assignment

Can't write this, since we don't have begin in our language
let $x=10$
$y=12$
in (begin set $x=+(x, 1)$
$\mathbf{x}$ )

## Evaluating with Assignment

> Instead, use a binding for a dummy variable d to sequence expressions; initial environment is empty

$$
\begin{aligned}
& \text { let } x=10 \\
& y=12 \\
& \text { in let } \mathbf{d}=\text { set } x=+(x, 1) \\
& \text { in } x
\end{aligned}
$$

## Evaluating with Assignment

## Eval RHS (right-hand side) of the let expression

$$
\begin{aligned}
& \text { let } \mathbf{x}=10 \\
& \mathbf{y}=12 \\
& \text { in let } \mathbf{d}=\text { set } \mathbf{x}=+(\mathbf{x}, 1) \\
& \quad \text { in } \mathbf{x}
\end{aligned}
$$

## Evaluating with Assignment



Extend the current environment with $\mathbf{x}$ and $\mathbf{y}$, and eval body

$$
\begin{aligned}
& \text { let } \mathbf{x}=10 \\
& \quad \mathbf{y}=12 \\
& \text { in let } d=\text { set } x=+(x, 1) \\
& \quad \text { in } x
\end{aligned}
$$

## Evaluating with Assignment



Eval RHS of the let expression
let $x=10$
$y=12$
in let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, 1)$
in $\mathbf{x}$

## Evaluating with Assignment



It modifies the $\mathbf{x}$ in the current lexical scope; we define set to always return 1

$$
\begin{aligned}
& \text { let } \mathbf{x}=10 \\
& \quad \mathbf{y}=12 \\
& \text { in let } \mathbf{d}=\text { set } x=+(x, 1) \\
& \text { in } x
\end{aligned}
$$

## Evaluating with Assignment



Bind $\mathbf{d}$ to the result 1; to eval the body, $\mathbf{x}$, we look it up in the environment as usual, and find 11

$$
\begin{aligned}
& \text { let } x=10 \\
& \quad y=12 \\
& \text { in let } d=\text { set } x=+(x, 1) \\
& \text { in } x
\end{aligned}
$$

## Evaluating with Assignment


> Variables now correspond to boxes in the environment, not fixed values

$$
\begin{aligned}
& \text { let } x=10 \\
& y=12 \\
& \text { in let } \mathbf{d}=\text { set } x=+(x, 1) \\
& \text { in } x
\end{aligned}
$$

## Expressed and Denoted Values

$$
\begin{array}{cl}
\text { <expval> } & ::=\text { <num> } \\
& ::=\text { <proc> } \\
\text { <denval> } & ::=\text { <reference> }
\end{array}
$$

- New datatype:


## (define-datatype reference reference? <br> (a-ref (pos integer?) <br> (vec vector?)))

- New function:
apply-env-ref : env sym -> ref


## Assignment and Closures

An example with proc; again, we start with the empty environment

```
let \(x=10\)
        \(y=12\)
in let \(f=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})\)
    in let \(\mathbf{d}=\) set \(\mathbf{x}=+(\mathbf{x}, 1)\)
        in (f 0)
```


## Assignment and Closures

## Eval RHS of the let expression

```
let \(\mathbf{x}=10\)
        \(\mathbf{y}=12\)
in let \(\mathbf{f}=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})\)
    in let \(\mathbf{d}=\) set \(\mathbf{x}=+(\mathbf{x}, 1)\)
        in (f 0)
```


## Assignment and Closures



Extend the current environment with $\mathbf{x}$ and $\mathbf{y}$, and eval body

```
let \(x=10\)
        \(y=12\)
in let \(\mathbf{f}=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})\)
    in let \(\mathbf{d}=\) set \(\mathbf{x}=+(\mathbf{x}, 1)\)
        in (f 0)
```


## Assignment and Closures



Eval RHS of the let expression...

$$
\begin{aligned}
& \text { let } x=10 \\
& \quad y=12 \\
& \text { in let } f=\operatorname{proc}(z)+(z, x) \\
& \text { in let } d=\operatorname{set} x=+(x, 1) \\
& \quad \text { in }(f 0)
\end{aligned}
$$

## Assignment and Closures


... which creates a closure, pointing to the current environment

```
let \(x=10\)
        \(y=12\)
in let \(\mathbf{f}=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})\)
    in let \(\mathbf{d}=\) set \(\mathbf{x}=+(\mathbf{x}, 1)\)
        in (f 0)
```


## Assignment and Closures



To finish the let, the environment is extended with $\mathbf{f}$ bound to the closure; then evaluate the body

```
let \(x=10\)
    \(y=12\)
in let \(f=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})\)
    in let \(\mathbf{d}=\) set \(\mathbf{x}=+(\mathbf{x}, 1)\)
        in (f 0)
```


## Assignment and Closures



Eval RHS of the let expression...
let $x=10$
$y=12$
in let $\mathbf{f}=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})$
in let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, 1)$ in (f 0)

## Assignment and Closures


... which changes the value of $\mathbf{x}$, then produces 1

$$
\begin{gathered}
\text { let } x=10 \\
y=12
\end{gathered}
$$

$$
\text { in let } f=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})
$$

$$
\text { in let } \mathbf{d}=\text { set } \mathbf{x}=+(\mathbf{x}, 1)
$$ in (f 0)

## Assignment and Closures



To eval the body, (f 0), we look up fin the environment to find a closure, and evaluate 0 to 0

$$
\begin{gathered}
\text { let } x=10 \\
y=12
\end{gathered}
$$

in let $\mathbf{f}=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})$
in let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, 1)$ in (f 0)

## Assignment and Closures



Extend the closure's environment with 0 for $\mathbf{z}$, and evaluate the closure's body in that environment; the result will be 11

```
let \(x=10\)
    \(y=12\)
in let \(\mathbf{f}=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})\)
    in let \(\mathbf{d}=\boldsymbol{s e t} \mathbf{x}=+(\mathbf{x}, 1)\)
        in (f 0)
```


## Assignment and Closures


>By capturing environments, closures capture variables that may change

```
let \(x=10\)
    \(y=12\)
in let \(f=\operatorname{proc}(\mathbf{z})+(\mathbf{z}, \mathbf{x})\)
    in let \(\mathbf{d}=\boldsymbol{s e t} \mathbf{x}=+(\mathbf{x}, 1)\)
        in (f 0)
```


## Assignment and Arguments

Another example with proc, but with the let inside the proc

```
let f = proc(z)
    let x = 10
        in let d = set x = +(x,z)
        in x
    in +((f 1), (f 9))
```


## Assignment and Arguments

Eval RHS of the let expression...
let $\mathbf{f}=\operatorname{proc}(\mathbf{z})$
let $x=10$
in let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$
in $x$
in +((f 1), (f 9))

## Assignment and Arguments

$z$ let $x=10$ in let $d=$ set $x=+(x, z)$ in $x$
... which creates a closure, pointing to the current environment
let $\mathbf{f}=\operatorname{proc}(\mathbf{z})$
let $\mathrm{x}=10$
in let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$
in $x$
in +((f 1), (f 9))

## Assignment and Arguments



Bind the closure to $\mathbf{f}$ and eval the body

```
let f = proc(z)
    let x = 10
    in let d = set x = +(x,z)
        in x
in +((f 1), (f 9))
```


## Assignment and Arguments



Evaluate the first operand, (f 1)
let $f=\operatorname{proc}(z)$
let $\mathrm{x}=10$
in let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$
in $x$
in +((f 1), (f 9))

## Assignment and Arguments



Take the closure for $\mathbf{f}$, extend its environment with a binding for $\mathbf{z}$, and eval the closure's body

```
let f = proc(z)
    let x = 10
    in let d = set x = +(x,z)
        in x
in +((f 1),(f 9))
```


## Assignment and Arguments



Eval the RHS

```
let f = proc(z)
    let x = 10
        in let d = set x = +(x,z)
        in x
    in +((f 1), (f 9))
```


## Assignment and Arguments



Add the binding for $\mathbf{x}$ and eval the inner body

```
let f= proc(z)
    let x = 10
    in let d = set x = +(x,z)
        in x
in +((f 1),(f 9))
```


## Assignment and Arguments



Eval RHS...

```
let f = proc(z)
    let x = 10
        in let d = set x = +(x,z)
        in x
in +((f 1),(f 9))
```


## Assignment and Arguments


... which modifies the value of $\mathbf{x}$

```
let f = proc(z)
    let x = 10
    in let d = set x = +(x,z)
        in x
in +((f 1), (f 9))
```


## Assignment and Arguments



Bind d to 1 and evaluate $\mathbf{x}$, which produces 11

```
let f = proc(z)
    let x = 10
        in let d = set x = +(x,z)
        in x
in +((f 1), (f 9))
```


## Assignment and Arguments



First operand is 11 ; now evaluate the second operand, (f 9)

```
let f = proc(z)
    let x = 10
    in let d = set x = +(x,z)
        in x
in +((f 1), (f 9))
```


## Assignment and Arguments



Again, take the closure for $\mathbf{f}$, extend the closure's environment with a binding for $\mathbf{z}$, and eval the closure's body

```
let \(\mathbf{f}=\operatorname{proc}(\mathbf{z})\)
    let \(x=10\)
    in let \(\mathbf{d}=\) set \(\mathbf{x}=+(\mathbf{x}, \mathbf{z})\)
        in \(x\)
    in +((f 1), (f 9))
```


## Assignment and Arguments



29


Add a binding for $\mathbf{x}$, then eval the inner body

```
let f = proc(z)
    let x = 10
        in let d = set x = +(x,z)
        in x
in +((f 1),(f 9))
```


## Assignment and Arguments



Again the d RHS modifies the value of $\mathbf{x}$, but using the new $\mathbf{z}$ and $\mathbf{x}$

```
let f = proc(z)
    let x = 10
    in let d = set x = +(x,z)
                in x
in +((f 1),(f 9))
```


## Assignment and Arguments



| 29 |
| :--- | :--- |



Bind $\mathbf{d}$ to 1 and evaluate $\mathbf{x}$, which produces 19
let $f=\operatorname{proc}(\mathbf{z})$ let $x=10$
in let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$
in $x$
in +((f 1), (f 9))

## Assignment and Arguments



29


So the operands are 11 and 19; The final result is 30
let $\mathbf{f}=\operatorname{proc}(\mathbf{z})$ let $x=10$
in let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$
in $\mathbf{x}$
in +((f 1), (f 9))

## Assignment and Arguments



| 29 |
| :--- | :--- |


$>$ Every evaluation of a binding expression creates a new variable (box)
let $f=\operatorname{proc}(\mathbf{z})$ let $x=10$
in let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$ in $x$
in +((f 1), (f 9))

## Assignment and Locals within Procedures

An example with a procedure in a procedure
let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$ let $d=$ set $x=+(x, z)$ in $x$
in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in...

## Assignment and Locals within Procedures

## Eval RHS of the let expression...

let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$
let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$ in $\mathbf{x}$
in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in...

## Assignment and Locals within Procedures

$$
x \operatorname{proc}(z) \text { let } d=\text { set } x=+(x, z) \text { in } x \bullet
$$

... which creates a closure, pointing to the current environment
let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$
let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$ in $\mathbf{x}$
in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in...

## Assignment and Locals within Procedures

$\mathrm{mk} / \mathrm{x} \operatorname{proc}(\mathrm{z})$ let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$ in $\mathbf{x} \bullet$

To finish the let, the environment is extended with $\mathbf{m k}$ bound to the closure, then evaluate the body

```
let \(\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})\)
    let \(\mathbf{d}=\) set \(\mathbf{x}=+(\mathbf{x}, \mathbf{z})\) in \(\mathbf{x}\)
in let \(\mathbf{f}=(\mathrm{mk} 10)\)
    in let \(\mathbf{g}=(\mathbf{m k} 12) \quad\) in...
```


## Assignment and Locals within Procedures



Eval RHS, a function call; look up mk...
let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$
let $d=$ set $x=+(x, z)$ in $x$
in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in...

## Assignment and Locals within Procedures

$m \mathrm{~m} / \mathrm{x} \operatorname{proc}(\mathbf{z})$ let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$ in $\mathbf{x} \cdot \boldsymbol{\bullet}$


It's a closure, so extend the closure's environment with 10, and eval the closure's body
let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$
let $d=$ set $x=+(x, z)$ in $x$
in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in...

## Assignment and Locals within Procedures

Qr

$$
\mathrm{mk} \quad \mathrm{x} \operatorname{proc}(\mathbf{z}) \text { let } \mathbf{d}=\text { set } \mathbf{x}=+(\mathbf{x}, \mathbf{z}) \text { in } \mathbf{x} \cdot \boldsymbol{\bullet}
$$

| $\mathbf{x} 10$ |
| ---: | :--- |

$$
\mathbf{z} \text { let d = set } \mathbf{x}=+(\mathbf{x}, \mathbf{z}) \text { in } \mathbf{x} \cdot
$$

Note that the variable $\mathbf{x}$ is in the closure's environment
let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$

$$
\text { let } \mathbf{d}=\text { set } \mathbf{x}=+(\mathbf{x}, \mathbf{z}) \text { in } \mathbf{x}
$$

in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in...

## Assignment and Locals within Procedures



Bind $\mathbf{f}$ to the closure, and evaluate the body
let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$ let $d=$ set $x=+(x, z)$ in $x$
in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in ...

## Assignment and Locals within Procedures



Eval RHS of the let expression, another call to $\mathbf{m k}$; same as before...
let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$ let $d=$ set $x=+(x, z)$ in $x$
in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in...

## Assignment and Locals within Procedures



Extend mk's env with a new $\mathbf{x}$ and get a closure, this time bound to $\mathbf{g}$

## Assignment and Locals within Procedures

$\mathrm{mk} \quad \mathrm{xproc}(\mathbf{z})$ let $\mathbf{d}=$ set $\mathbf{x}=+(\mathbf{x}, \mathbf{z})$ in $\mathbf{x} \cdot$
$\times 10$

$\mathrm{x} \mid 12$
g| $>$ let d = set $x=+(x, z)$ in $x \cdot$

At this point, $\mathbf{f}$ and $g$ have private versions of $\mathbf{x}$
let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$
let $d=$ set $x=+(x, z)$ in $x$
in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in ...

## Assignment and Locals within Procedures


> Closures can capture generated variables, effectively getting private state
let $\mathbf{m k}=\operatorname{proc}(\mathbf{x}) \operatorname{proc}(\mathbf{z})$ let $d=$ set $x=+(x, z)$ in $x$
in let $\mathbf{f}=(\mathbf{m k} 10)$
in let $\mathbf{g}=(\mathbf{m k} 12) \quad$ in...

## Assignment Summary

- Variables now denote references (a.k.a. locations), not values
- Lexical scope still works

