

Demo of ISP  
Eclipse GUI  
Command-line Options  
Set-up Audience with LiveDVD

About 30 minutes - by Ganesh

# Understanding Out-of-order Execution

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- For MPI programs, no dynamic tool can maintain in-order (in program order) issue of all instructions and yet cover the full extent of non-determinism!
- Luckily, the MPI semantics are such that out-of-order issue of ‘hijacked’ instructions is OK !
- Theory of ‘happens-before’ (after break) justifies
- Illustration on Crooked-barrier now
  - Real code illustration is through POE-Illustration

# Demonstrate on Auto-send Examples

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- Experiment with Autosend examples
  - MPI\_Autosend1 and Autosend2
    - These execute without deadlocks
    - Stepping in internal issue order shows how ISP performs things
    - The IntraCompletesBefore (recently called IntraHappensBefore) justifies this out-of-order execution
  - MPI\_Autosend3
    - This example deadlocks
    - Pattern is R(from:0,h1); W(h1); B; S(to:0,h2); W(h2)
  - Try variant MPI\_Autosend4
    - Have pattern S(to:0,h2); W(h2); B; R(from:0,h1); W(h1)
    - Run with and without buffering
    - What do you observe?
      - The case with buffering does not deadlock!

## ‘Crooked Barrier’ in POE-Illustration.c

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Process P0

Isend(1, req) ;

**Barrier ;**

Wait(req) ;

Process P1

Irecv(\*, req) ;

**Barrier ;**

Recv(2) ;

Wait(req) ;

Process P2

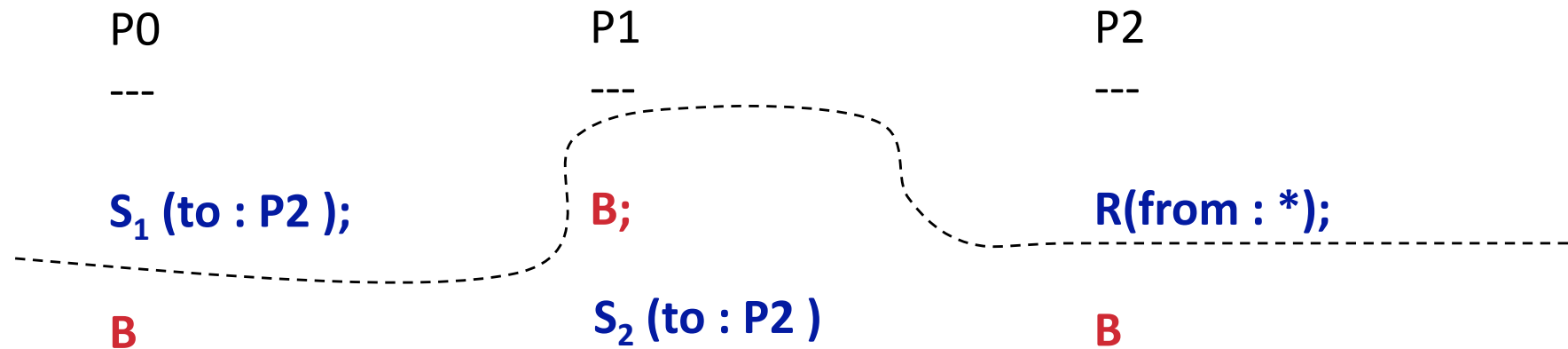
**Barrier ;**

Isend(1, req) ;

Wait(req) ;

# The “Crooked Barrier” example

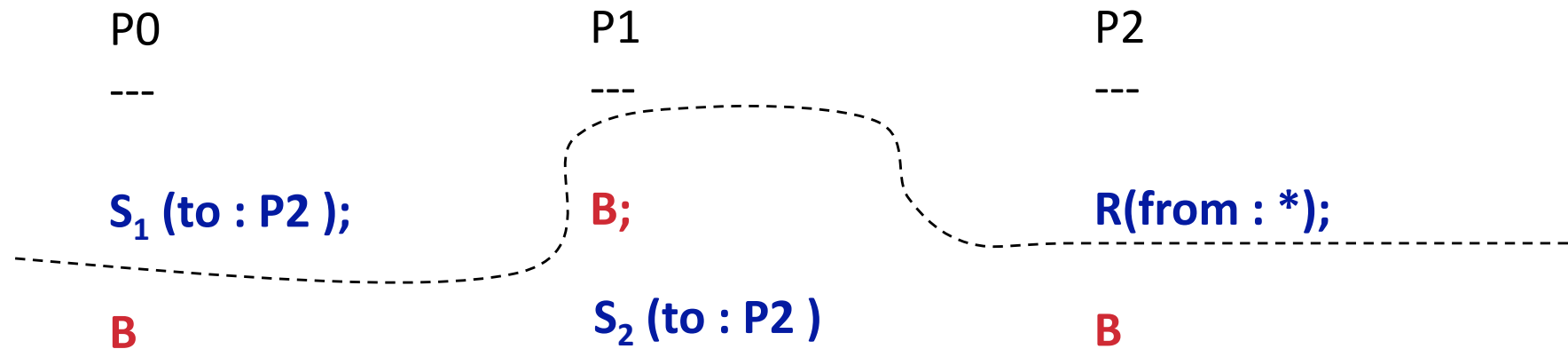
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Can `S2 (to : P2 )` match `R(from : *)` ?

# The “Crooked Barrier” example

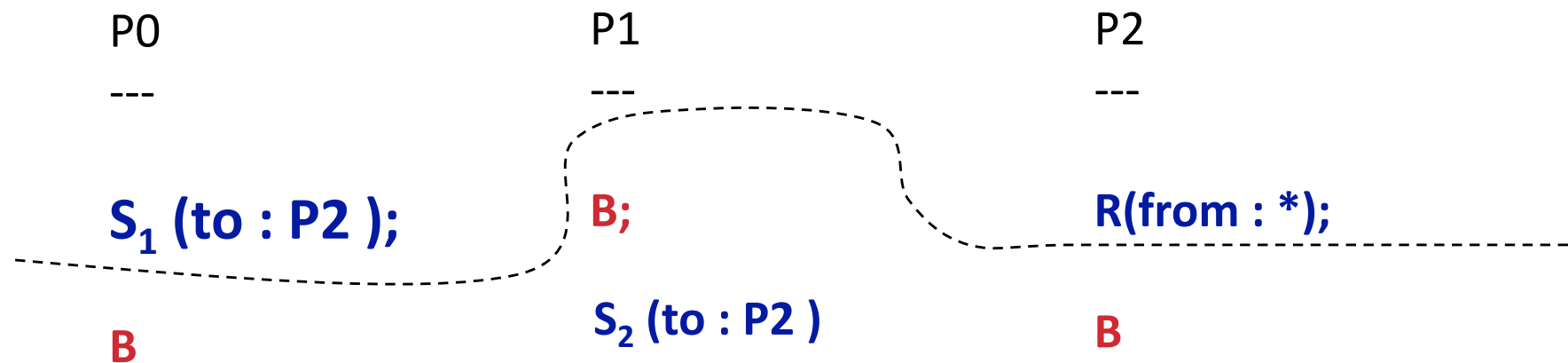
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Can **S<sub>2</sub> (to : P2 )** match **R(from : \*)** ? **YES! Here is how!**

# The “Crooked Barrier” example

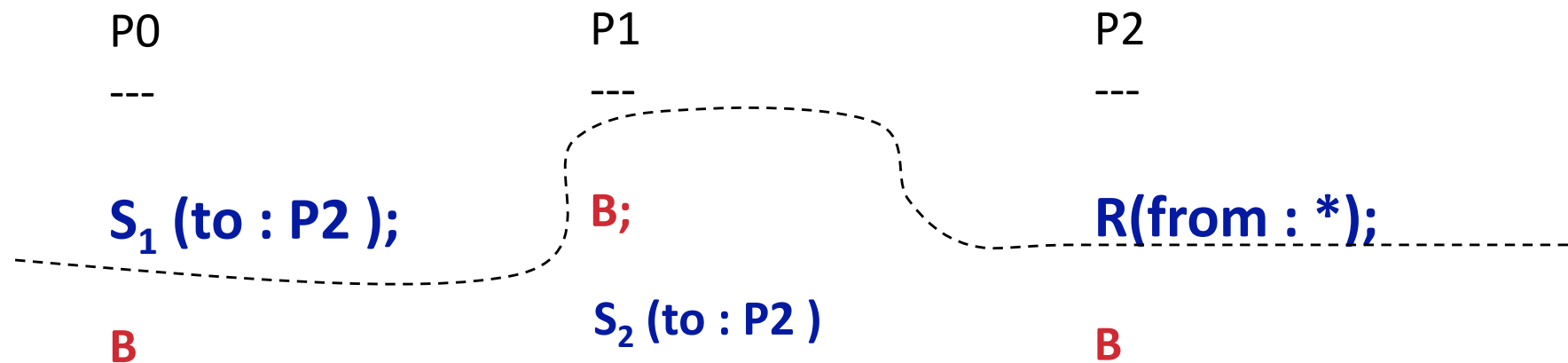
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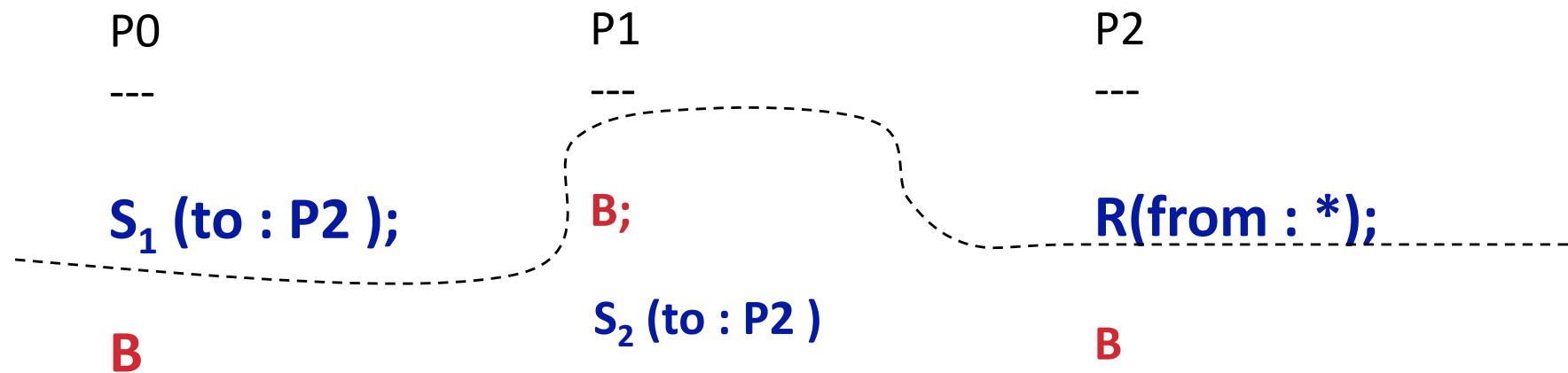
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Can `S2 (to : P2 )` match `R(from : *)` ? YES! Here is how!

# The “Crooked Barrier” example

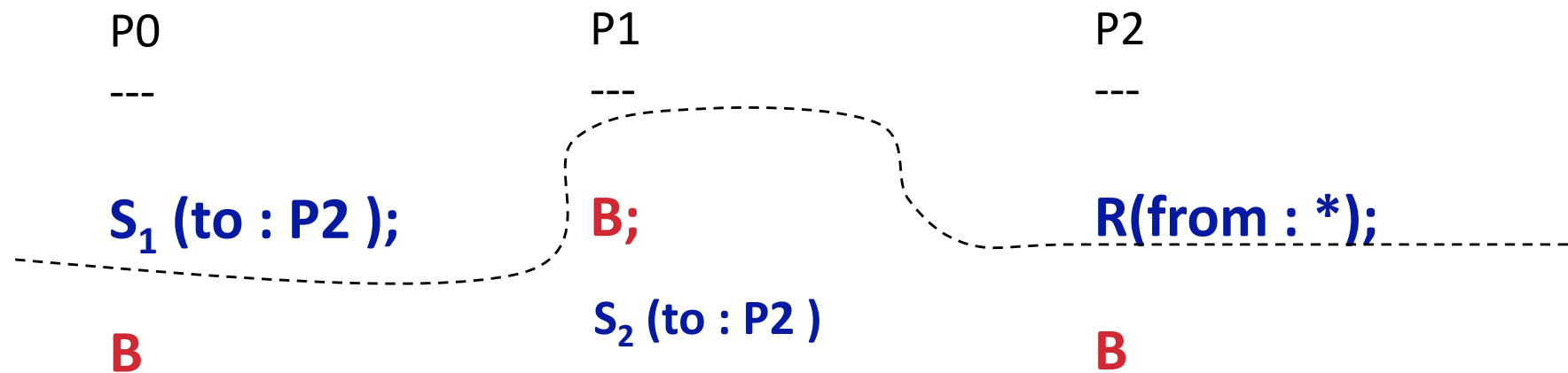
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Can `S2 (to : P2 )` match `R(from : *)` ? **YES! Here is how!**

# The “Crooked Barrier” example

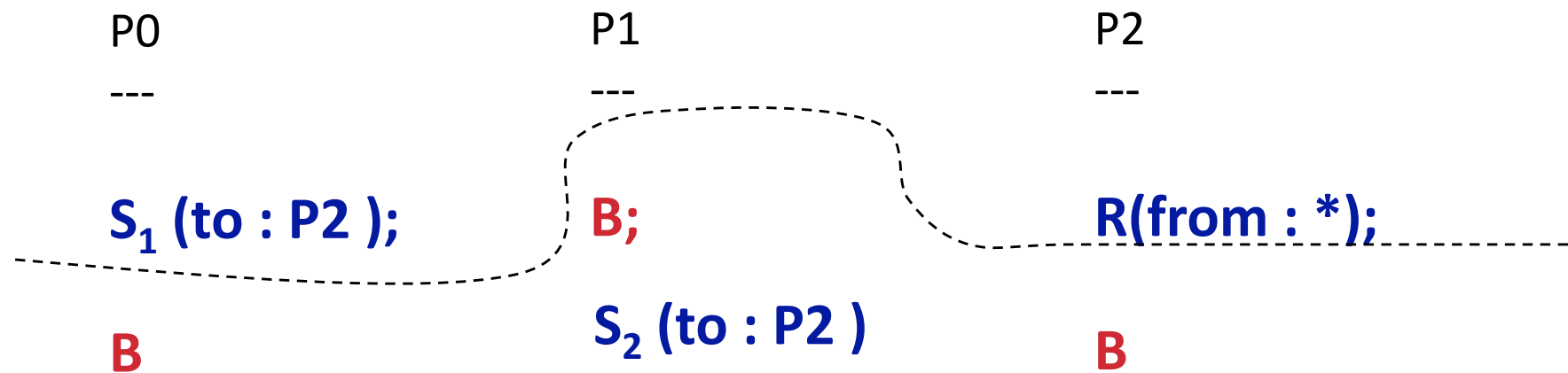
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Can  $S_2$  (to : P2 ) match  $R$ (from : \*) ? YES! Here is how!

# The “Crooked Barrier” example

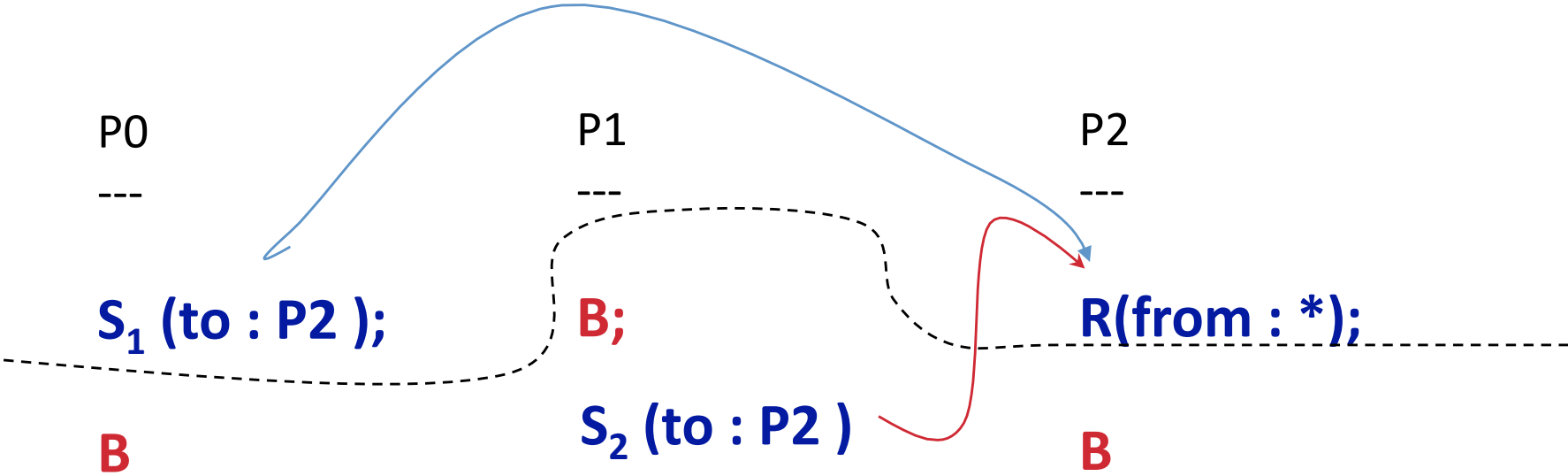
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Can  $S_2(\text{to} : P2)$  match  $R(\text{from} : *)$ ? **YES!** Here is how!

# The “Crooked Barrier” example

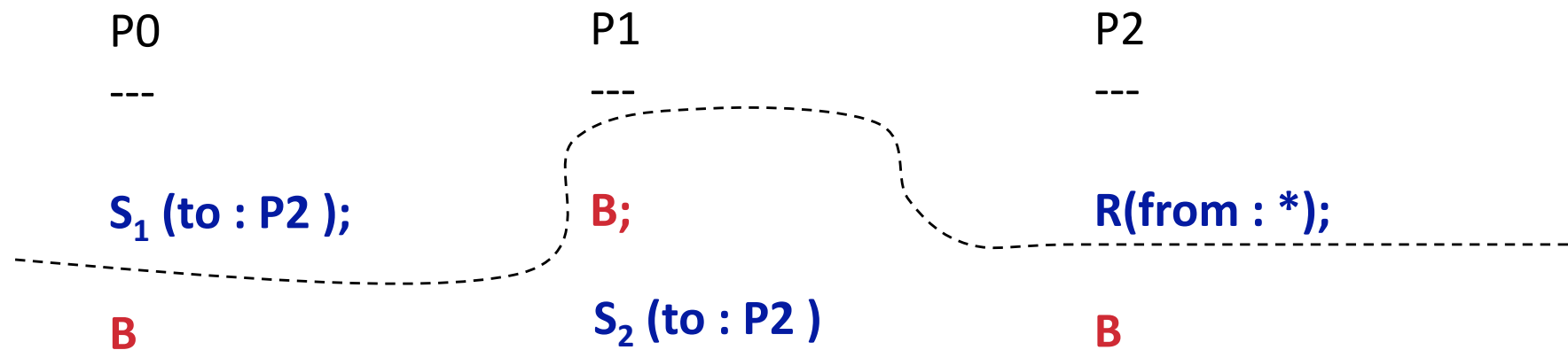
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# The “Crooked Barrier” example

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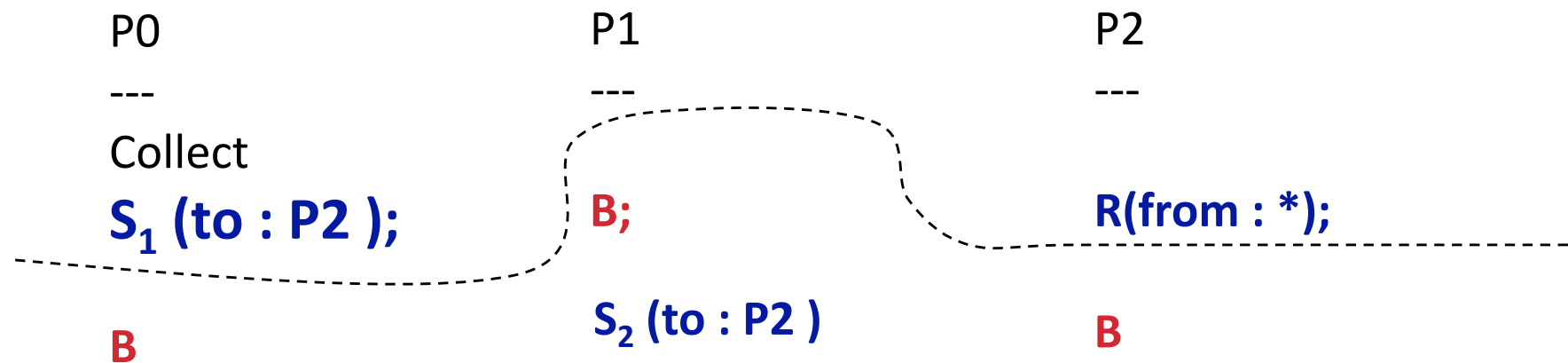


ISP handles this situation through

- Out-of-order execution
- Dynamic Instruction Rewriting

# The “Crooked Barrier” example

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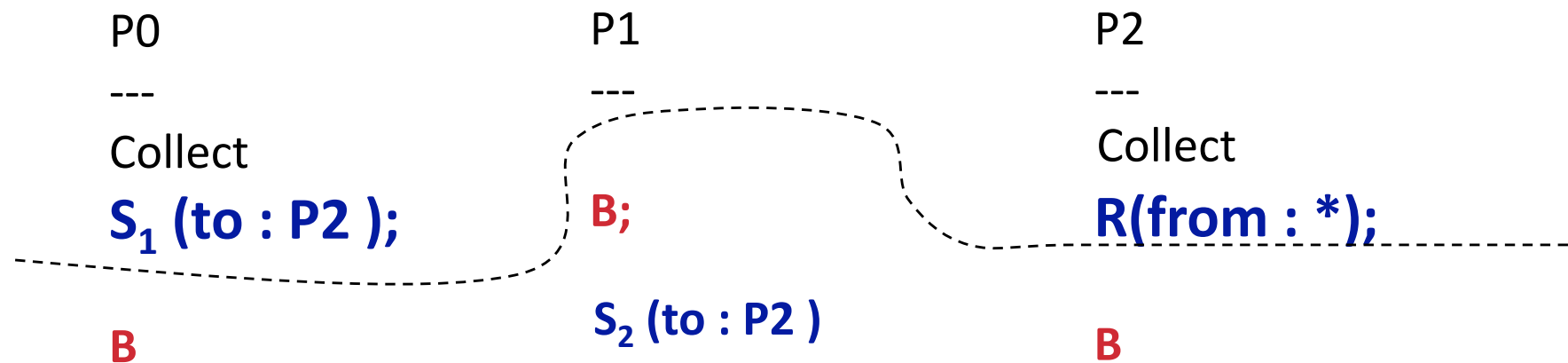


ISP handles this situation through

- Out-of-order execution
- Dynamic Instruction Rewriting

# The “Crooked Barrier” example

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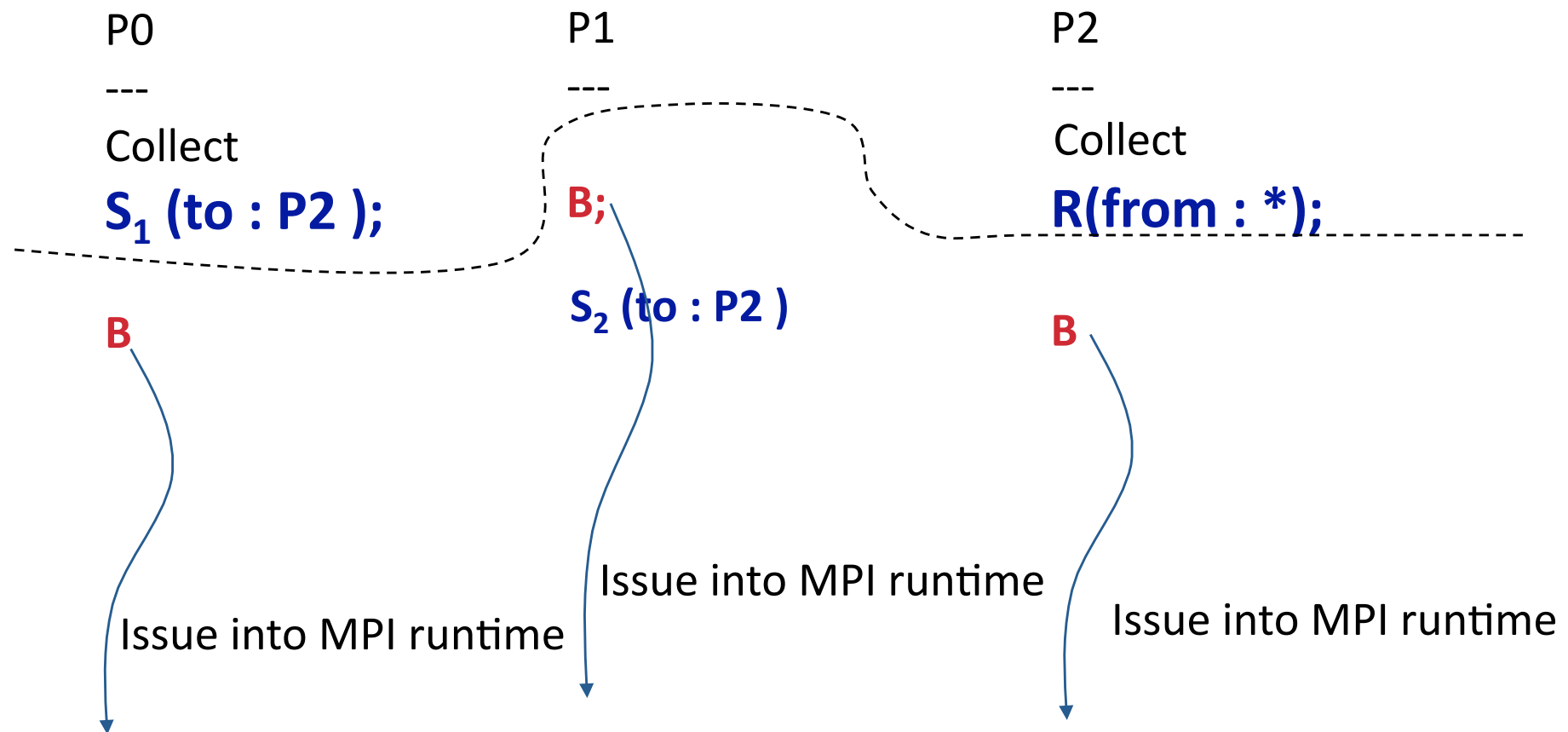


ISP handles this situation through

- Out-of-order execution
- Dynamic Instruction Rewriting

# The “Crooked Barrier” example

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# The “Crooked Barrier” example

---

P0

---

Collect

**S<sub>1</sub> (to : P2 );**

B

P1

---

B;

**S<sub>2</sub> (to : P2 )**

P2

---

Collect

**R(from : \*);**

B

# The “Crooked Barrier” example

---

P0

---

Collect

**S<sub>1</sub> (to : P2 );**

B

P1

---

B;

Collect

**S<sub>2</sub> (to : P2 )**

P2

---

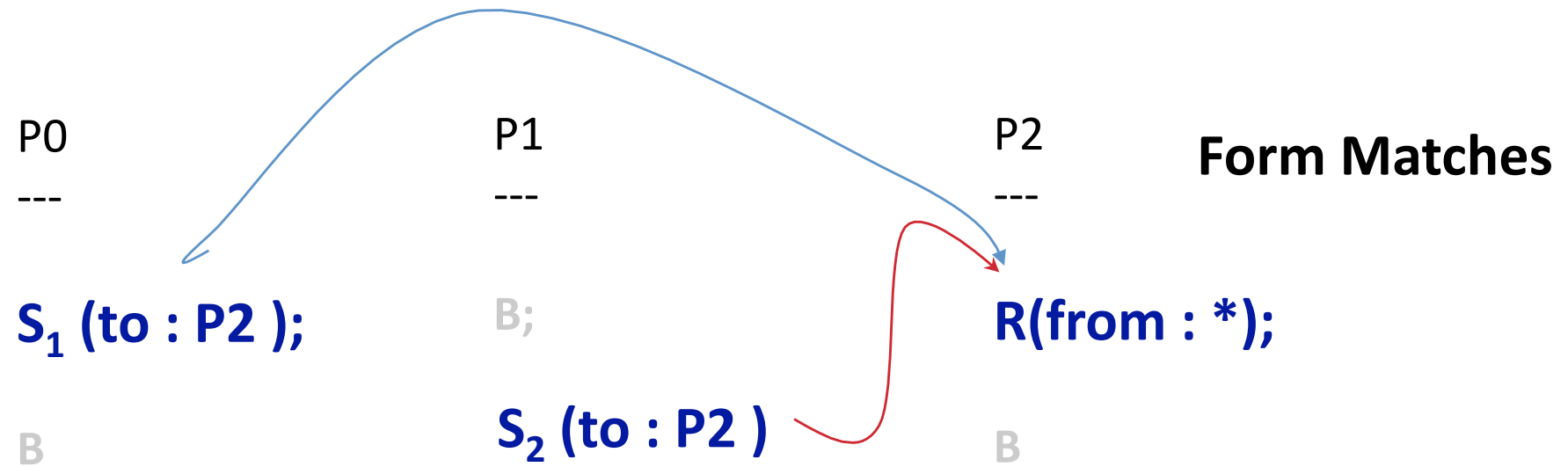
Collect

**R(from : \*);**

B

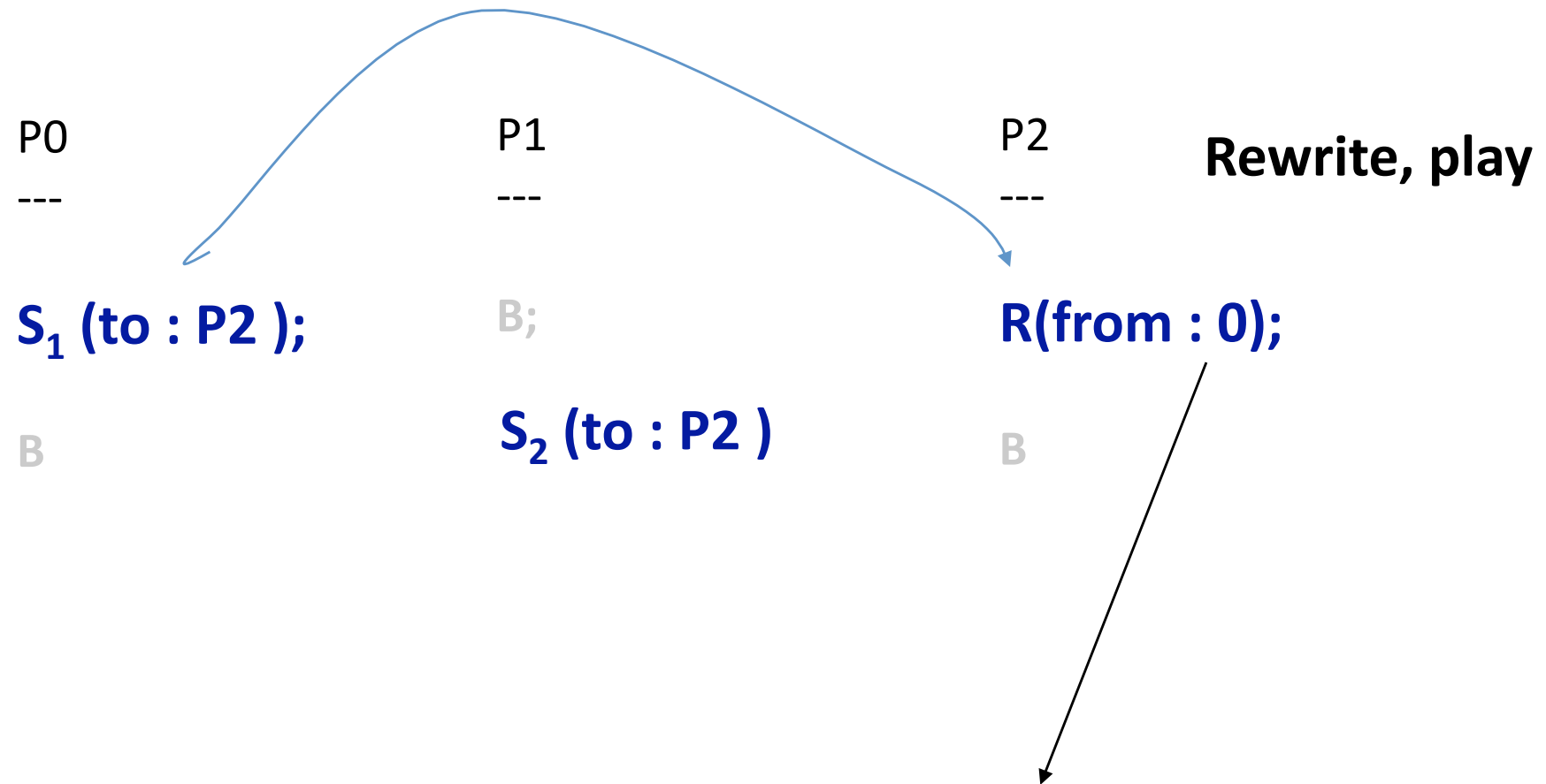
# The “Crooked Barrier” example

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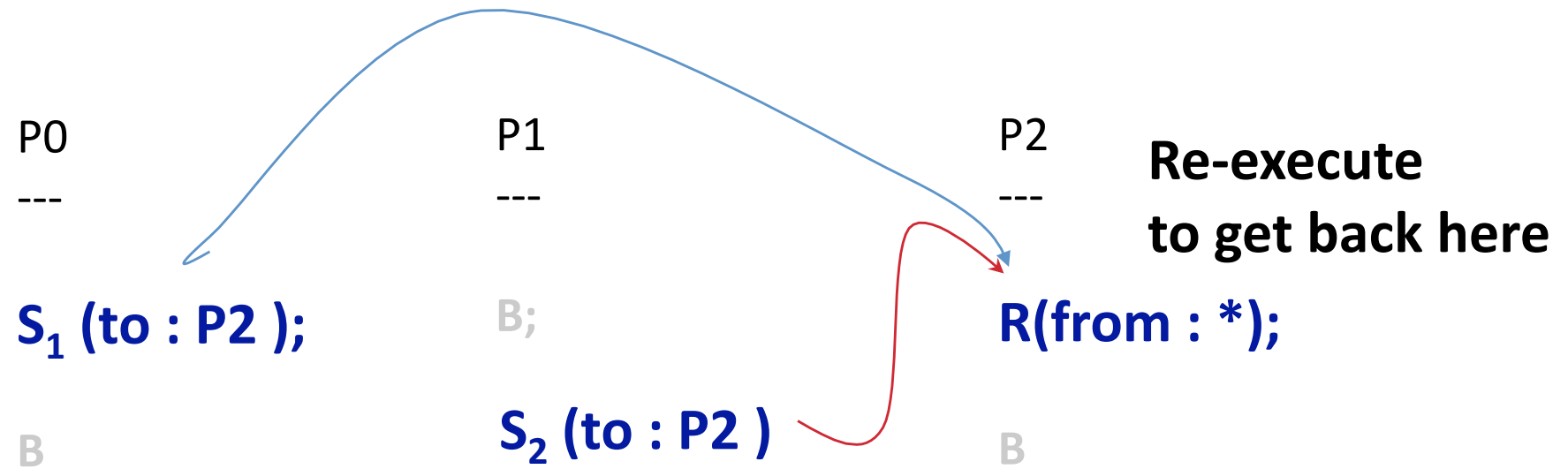
# The “Crooked Barrier” example

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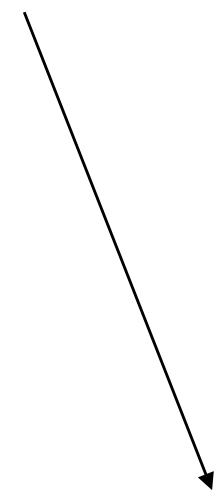
---

P0  
---  
Collect  
**S<sub>1</sub> (to : P2 );**  
B

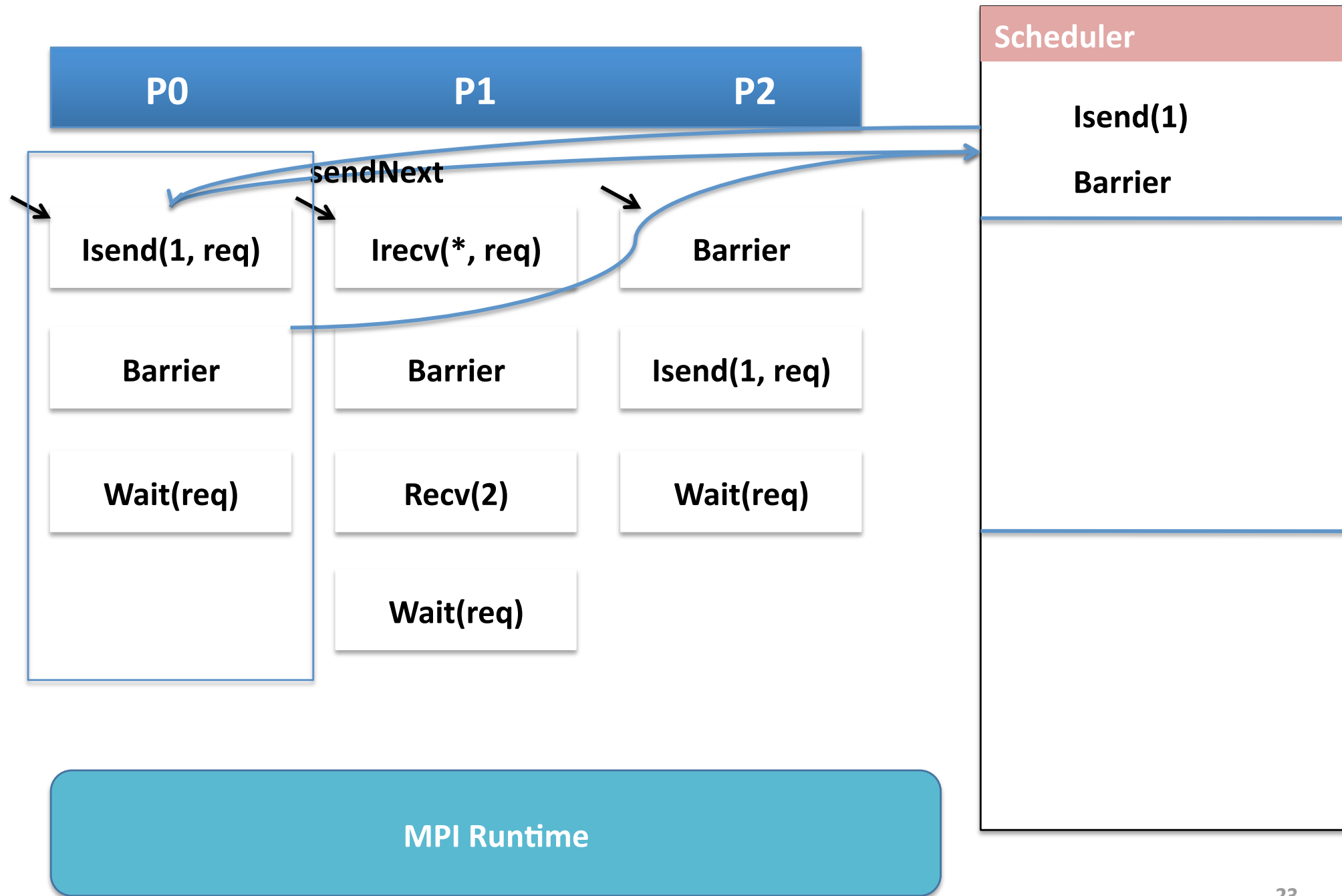
P1  
---  
B;  
Collect  
**S<sub>2</sub> (to : P2 )**

P2  
---  
Collect  
**R(from : 1);**  
B

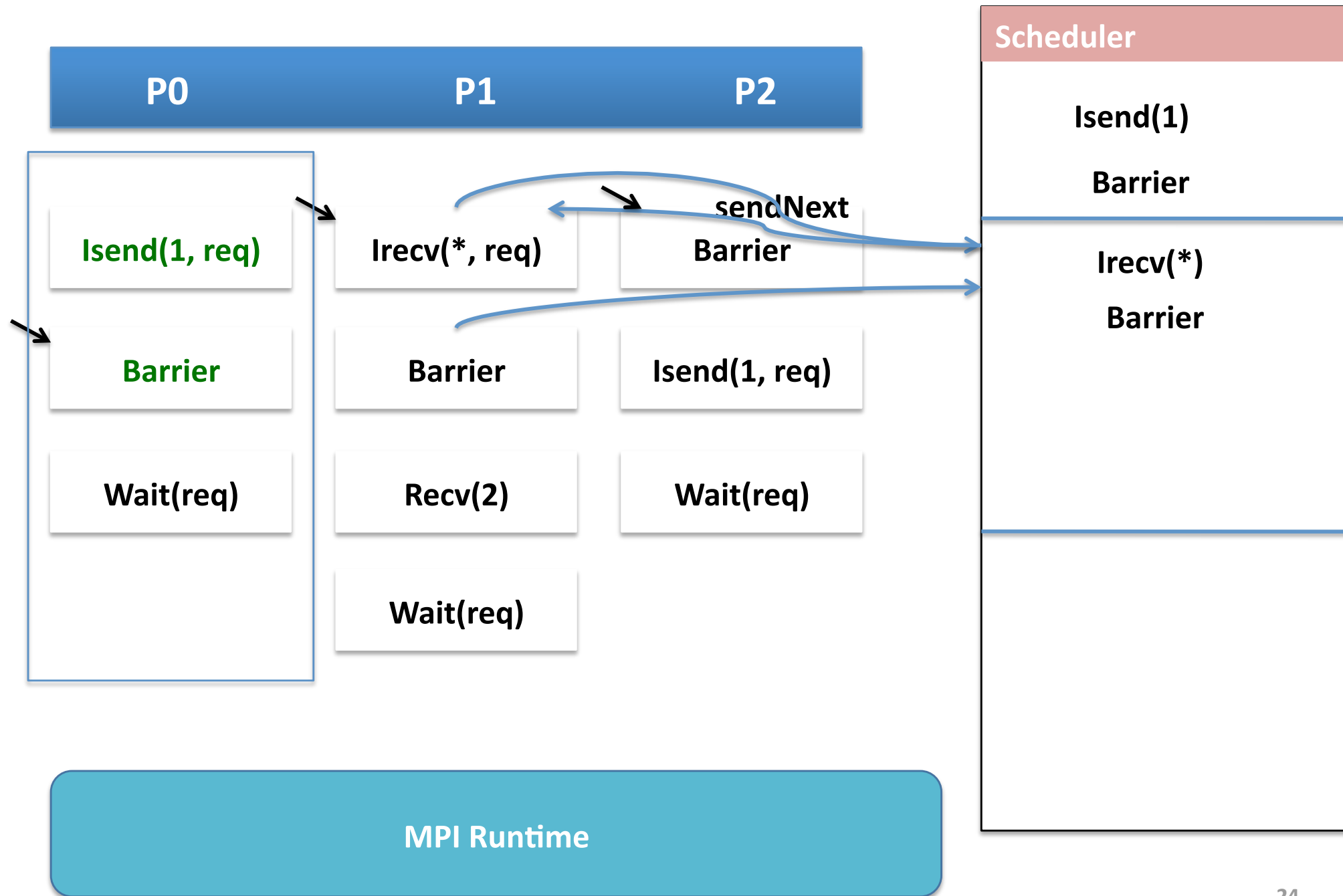
**Rewrite, play**



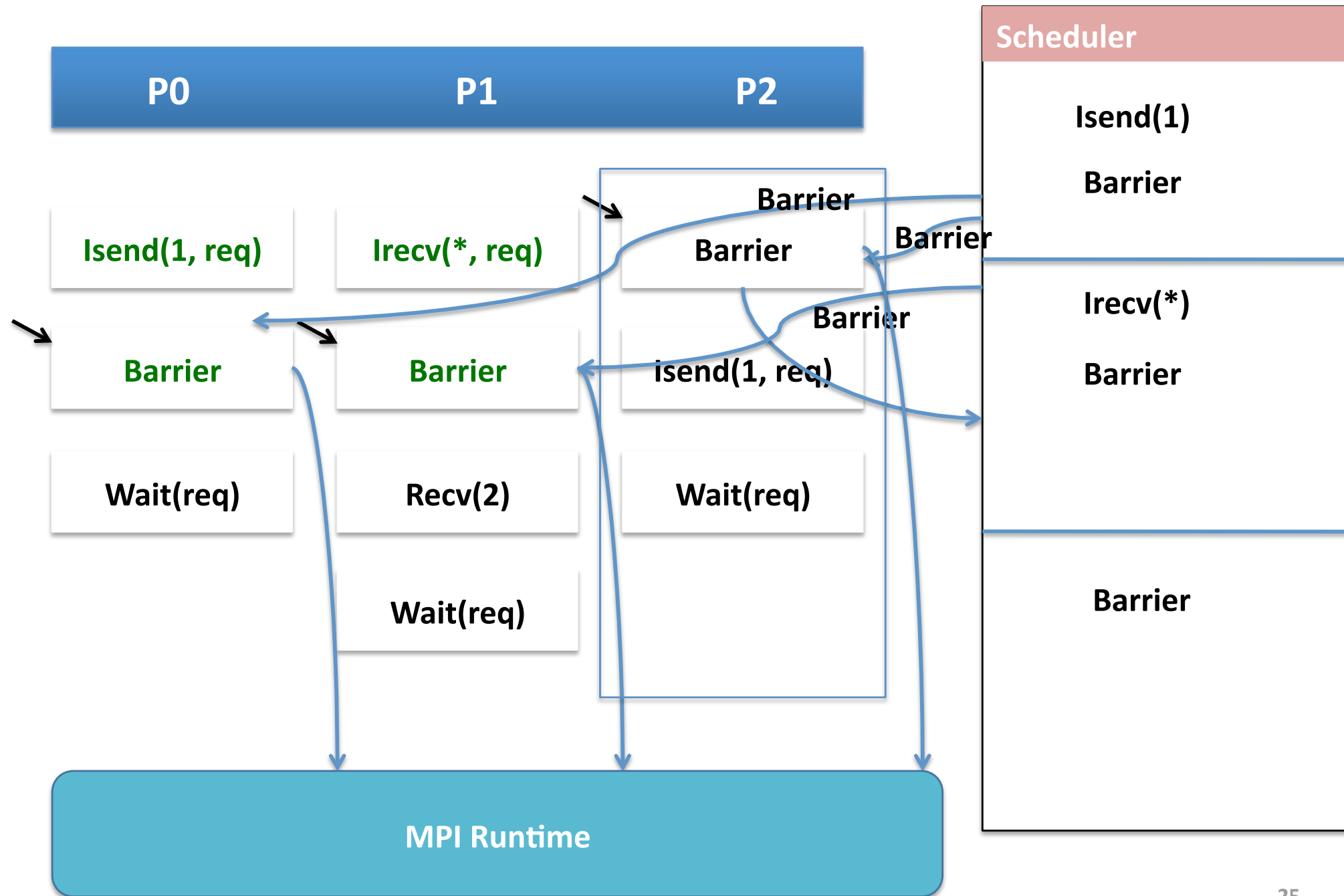
# Hijack Calls, Generate Relevant Interleavings



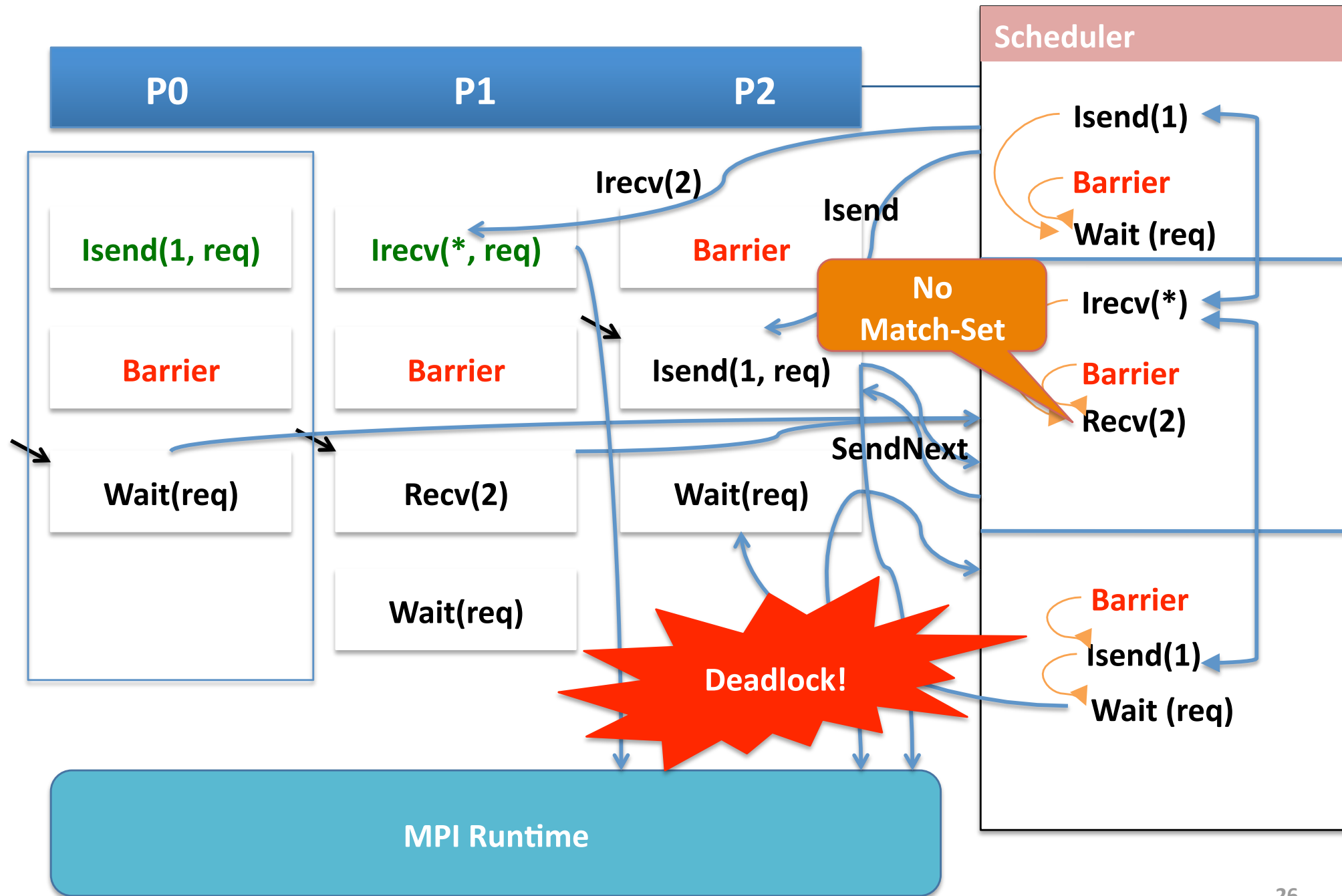
# Hijack Calls, Generate Relevant Interleavings



# Hijack Calls, Generate Relevant Interleavings



# Hijack Calls, Generate Relevant Interleavings



## Example 1: Any-src-can-deadlock9.c (Umpire) lucky.c and unlucky.c are just variations

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- **Distribute LiveDVDs**
  - Help audience boot into it
- **Source code of any-src-can-deadlock.c**
  - Demonstrate Verification Options
  - Observe completes-before
- **Source code of POE-illustration.c**
  - Observe completes-before
  - Observe internal issue order and program order

# LiveDVD Instructions

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- For Windows, boot using LiveDVD
  - Select 'try Linux without installing'
- For MAC, select boot device with LiveDVD present inside drive
  - Select to be DVD
  - Power down
  - Power-up with DVD in drive
- For MAC, you can restore boot device by pressing and holding down ALT / Option when booting again

# Observations of examples lucky / unlucky

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- ISP finds deadlock
- In Java GUI, do these
  - Watch the execution deadlock
  - Learn how to step
    - Step into different interleavings
    - Step through the traces of an interleaving
  - Do the above with and without rank locking
  - Do the above by locking ranks and watching
    - According to Program Order
    - According to Internal Issue Order
- Watch the IntraCB and InterCB

End of C