

Today

- ◆ **Course overview**
- ◆ **Program analysis introduction**

Organization

◆ Lectures

- **Tuesdays: I lecture**
- **Thursdays: You present about specific tools**
 - **I'll start this week**
 - **Nathan will go next week**

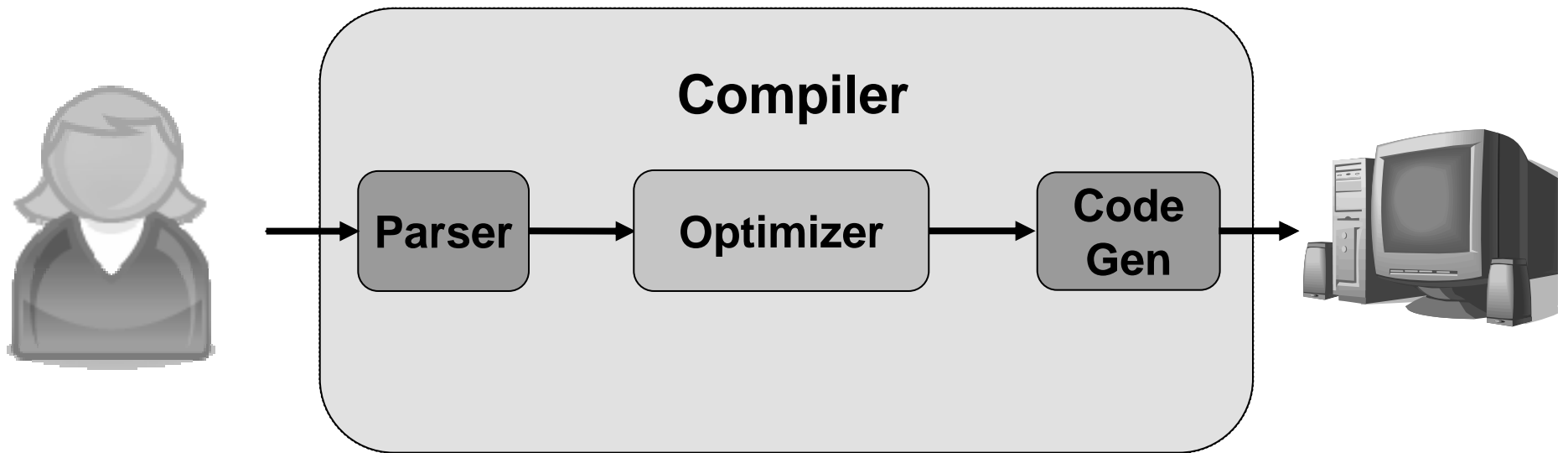
Projects

- ◆ **Each student (or group or two) will complete a project**
- ◆ **Extend an existing analysis tool or compiler**
- ◆ **There are lots of open source tools that make great starting points for projects**
 - **CIL, cXprop, gcc, Saturn, LLVM, Locksmith, Sparse**
- ◆ **Example projects:**
 - **Add a new optimization**
 - **Add a new analysis feature**
 - **Make an existing analysis more precise**
 - **Etc.**

More Projects

- ◆ **1-2 page proposal**
 - **Due mid-semester**
- ◆ **Implementation + in-class presentation**
 - **Due a few weeks before end of semester**

What's in a compiler?



What's in an optimizer?

1. **Compute properties of a program using static analysis**
 2. **Apply semantics-preserving transformations based on these properties**
 - ◆ **Goal: Improve some resource metric: memory usage, execution speed, etc.**
- ◆ **The field of program analysis started with optimizing compilers**
 - **But program analysis is more broadly used today**

“Program analysis offers static compile-time techniques for predicting safe and computable approximations to the set of values or behaviors arising dynamically at run-time when executing a program on a computer.”

**From *Principles of Program Analysis*,
Nielson & Nielson & Hankin**

Program Analysis

- ◆ **Used in:**
 - **Bug finders**
 - **Program verifiers**
 - **Code refactoring tools**
 - **Garbage collectors**
 - **Runtime monitoring system**
 - **And... compilers**

- ◆ **This class is about the fundamental program analysis techniques used in all of these applications**

Course Goals

- ◆ **Understand basic techniques for doing program analyses and transformations**
 - **These techniques are the cornerstone of a variety of program analysis tools**
 - **They may come in handy, no matter what research you end up doing**
- ◆ **Get a feeling for what research is like in the area by reading research papers, and getting your feet wet in a small research project**

Motivation

- ◆ **Improve program correctness**
 - **Prevent bugs**
 - **Find bugs**
 - **Predict resource usage**
 - **Prove properties**

- ◆ **Improve program performance**
 - **Avoid redundant computations**
 - **Substitute fast computations for slow ones**

- ◆ **Example program property:**
 - “Z holds the value 17 at line 10”
 - Why is this useful?
- ◆ **Challenges: We need to compute this property**
 - **Safely**
 - **Precisely**
 - **Efficiently**

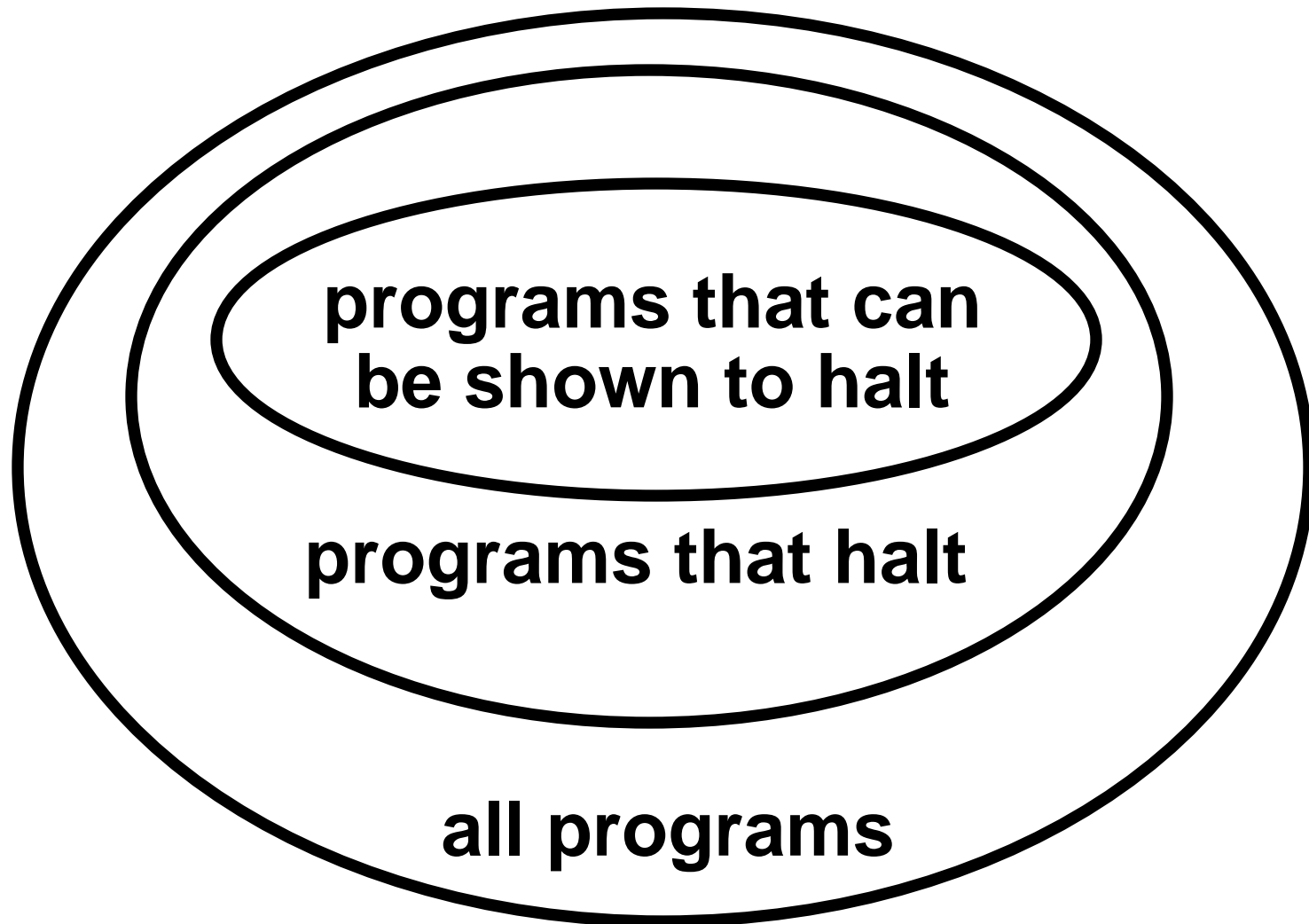
Challenge 1: Safety

- ◆ **“Yes” means “absolutely yes, for all possible executions of the program”**
 - **Ideally, analyzer could emit a proof**
 - **What would this look like?**
- ◆ **“No” means “probably not but I don’t know for sure”**
 - **Proof could not be found**
- ◆ **Note: Sometimes unsafe analyses are useful**

Challenge 2: Precision

- ◆ **Imprecise analysis is easy**
 - Just say “I don’t know” every time
- ◆ **How much precision is necessary?**
 - For example, do we need to prove that $Z==17$, or is it ok just to prove that $Z > 0$?
- ◆ **Imprecision is inevitable**
 - All interesting program analysis questions are undecidable when phrased precisely

Precision Cont'd



**programs that can
be shown to halt**

programs that halt

all programs

Challenge 3: Efficiency

- ◆ **Inherent conflict between precision and efficiency**
 - How long to analyze all of Windows XP?
 - Many state-of-the-art analysis techniques do not scale to really big programs
- ◆ **Claim: Only modular analyses truly scale**
 - Even $O(n)$ may be unacceptable if the analysis is not modular
 - Example of a modular analysis: Compute per-function summaries, then analyze the summaries together

Awful Example

```
int foo (void) {  
    int x = 0;  
    return x++ + x++;  
}
```

- ◆ **What does this function return?**
- ◆ **C compiler can pick an order in which to evaluate exprs**
- ◆ **An analyzer outside of the compiler must evaluate all possibilities**

Language Semantics

- ◆ **Semantics tells us what a program in the language means**
 - “Mathematical model that predicts the behavior of a program.”
- ◆ **Scope is limited – does not tell us:**
 - If the program is correct
 - What system calls the program makes
 - What instructions the program uses
 - How fast the program is, or how much memory it uses
- ◆ **The field of static analysis was developed to support optimization**
 - But w/o language semantics
 - This led to confusing papers and broken optimizers

Loose Semantics

- ◆ A “conforming” C program produces the same results on all compiler / platform combinations
- ◆ The C standard tells us what conforming programs look like
- ◆ Not all legal programs conform

```
int x[10];  
y = x[15];
```
- ◆ C semantics does not tell us what value is in `y` after the assignment
 - Of even if the program keeps running

Loose Semantics

- ◆ **Claim: The compiler for a well-designed language can statically check for conformance**
- ◆ **Does C have this property?**
 - **No: It is undecidable whether a C program conforms**
 - **Oops**
 - **Optimizing C compilers break non-conforming programs without informing the user**

More Semantics

- ◆ **Analyses only apply to conforming programs**
 - **Most large programs will contain non-conforming behavior**
- ◆ **Unfortunate implication: Can't actually predict the behavior of a C program**
 - **That is, the mathematical model doesn't apply to most real programs**
 - **Thus the buffer overflow vulnerabilities that make life difficult on the Internet**
- ◆ **In this course we'll be taking a pretty relaxed view of semantics**

Course Topics

- ◆ **Representations**
 - **Abstract Syntax Tree**
 - **Control Flow Graph**
 - **Dataflow Graph**
 - **Static Single Assignment**
 - **Control Dependence Graph**
 - **Program Dependence Graph**

Course Topics

- ◆ **Analysis/Transformation Algorithms**
 - **Dataflow Analysis**
 - **Interprocedural analysis**
 - **Pointer analysis**
 - **Abstract interpretation**
 - **Constraint-based analysis**
 - **Interaction between transformations and analyses**

Course Topics

- ◆ **Applications**
 - **Scalar optimizations**
 - **Loop optimizations**
 - **Program verification**
 - **Bug finding**
 - **Garbage collection**

Analysis Issues

- ◆ **How much of the program do we see?**
 - All?
 - One file at a time?
 - One library at a time?
- ◆ **Any additional inputs?**
 - Human help?
 - Profile data?

Analysis Issues

- ◆ **Analysis/compilation model**
 - **Separate compilation/analysis**
 - **quick, but no opportunities for interprocedural analysis**
 - **Link-time**
 - **allows interprocedural and whole program analysis**
 - **but what about shared precompiled libraries?**
 - **and what about compile-time?**
 - **Run-time**
 - **best optimization/analysis potential (can even use run-time state as additional information)**
 - **can handle run-time extensions to the program**
 - **but severe pressure to limit compilation time**

Summary

- ◆ **For next Tuesday: Read “Lecture Notes on Static Analysis,” pages 1-34**