

# HOIST: A System for Automatically Deriving Static Analyzers for Embedded Systems

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- Hoist makes it significantly easier to do static analysis of embedded software
  - E.g. TinyOS
- Automatically derives transfer functions for analyzing object code
  - This is new
  - Hoisted transfer functions are maximally precise
  - Brute-force approach that works well for small architectures

# Before Hoist:

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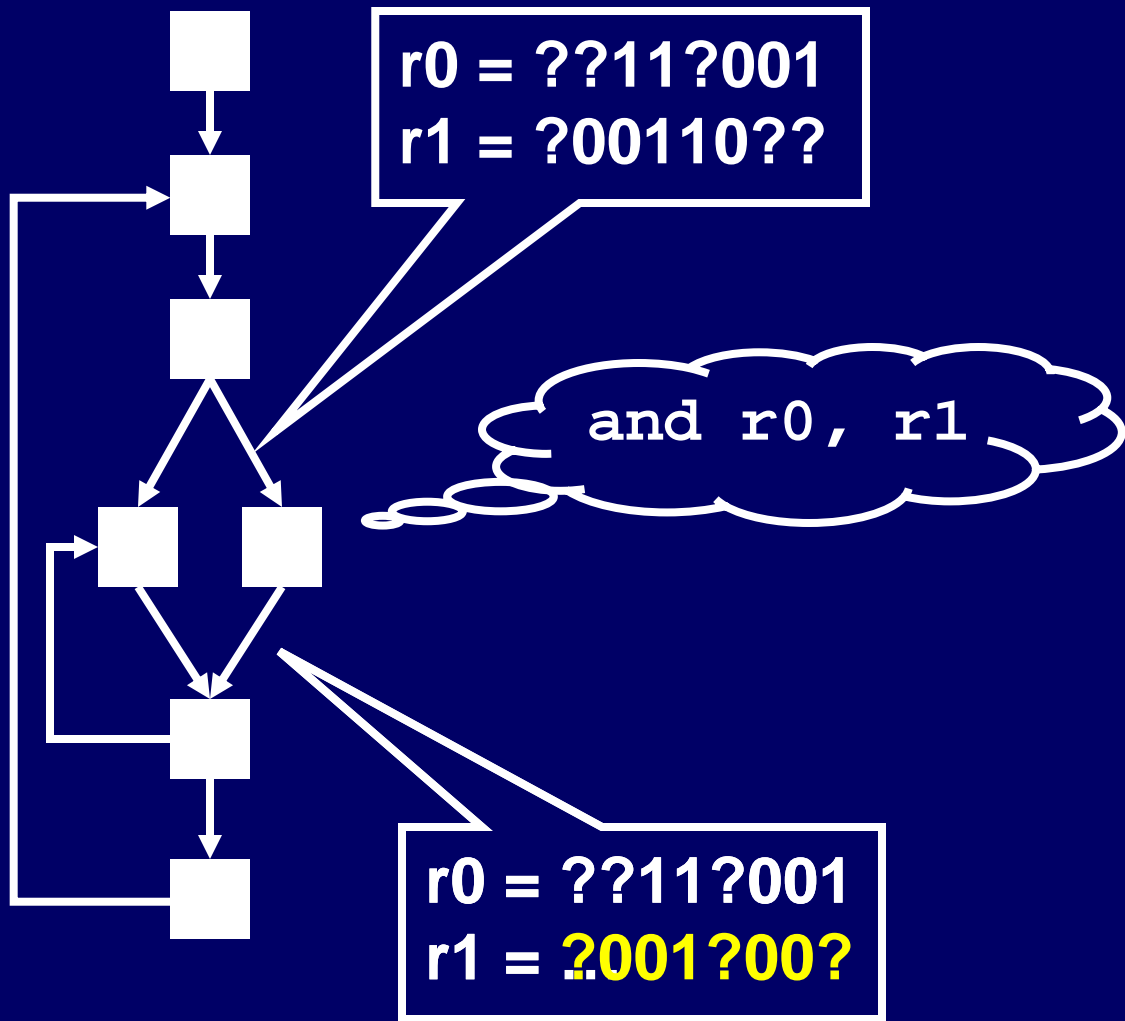
# After:

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# Use Static Analysis to Eliminate...

- Concurrency errors
- Deadline misses
- Stack overflow
- Language-level errors
  - Array bound violations
  - Null pointer dereferences
  - Numerical problems
- Everything else Jim Larus talked about!





&	0	1	?
0	0	0	0
1	0	1	?
?	0	?	?

Transfer Function

# Abstract Transfer Functions

$$\begin{array}{l} \text{??11?001} \\ \& \text{?00110??} \\ = \text{?001?00?} \end{array} \quad \begin{array}{l} [ 3.. 6] \\ + [36..60] \\ = [39..66] \end{array}$$

$$\begin{array}{l} \text{??11?001} \\ + \text{?00110??} \\ = \dots \end{array} \quad \begin{array}{l} [ 3.. 6] \\ \& [36..60] \\ = \dots \end{array}$$

Bitwise

Interval

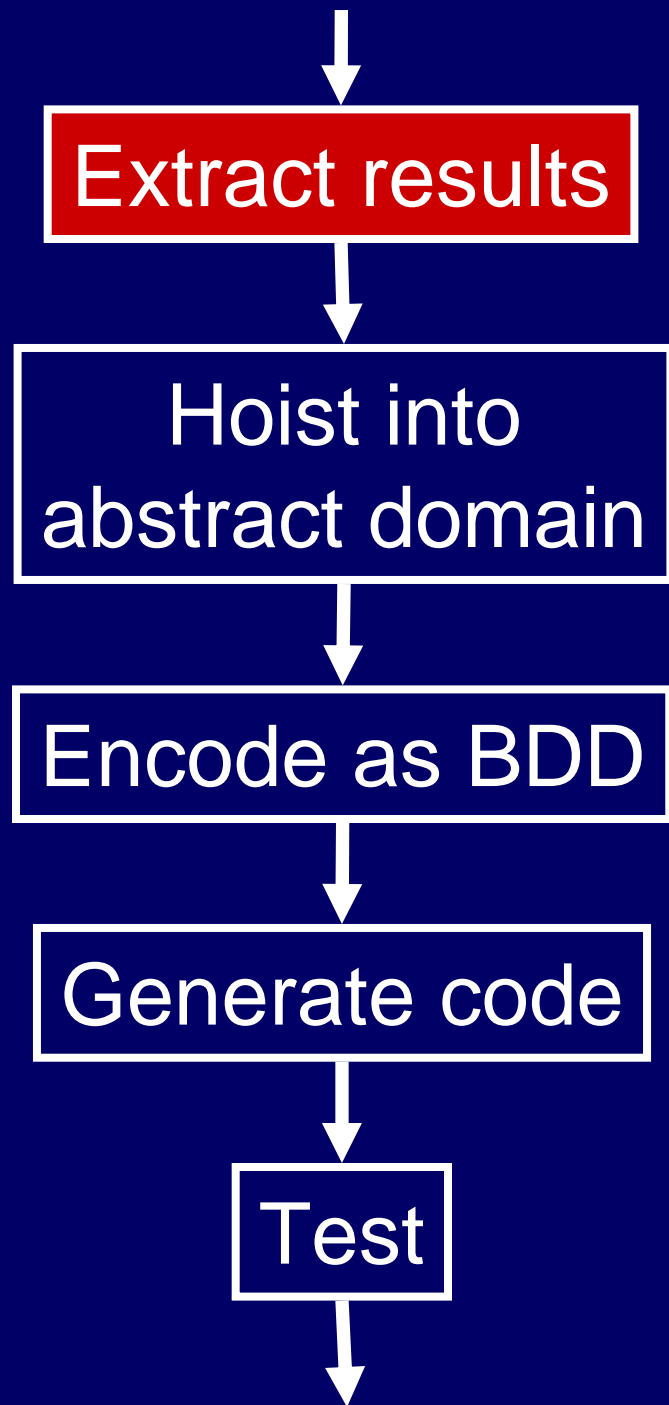
# Transfer Functions can be Hard

- Domain / operation mismatch
- Condition codes – input and output
- Hard to know where precision matters
- Lots of transfer functions:  
# domains \* # instructions \* # architectures
- **Result:** Wasted time, bugs, imprecision

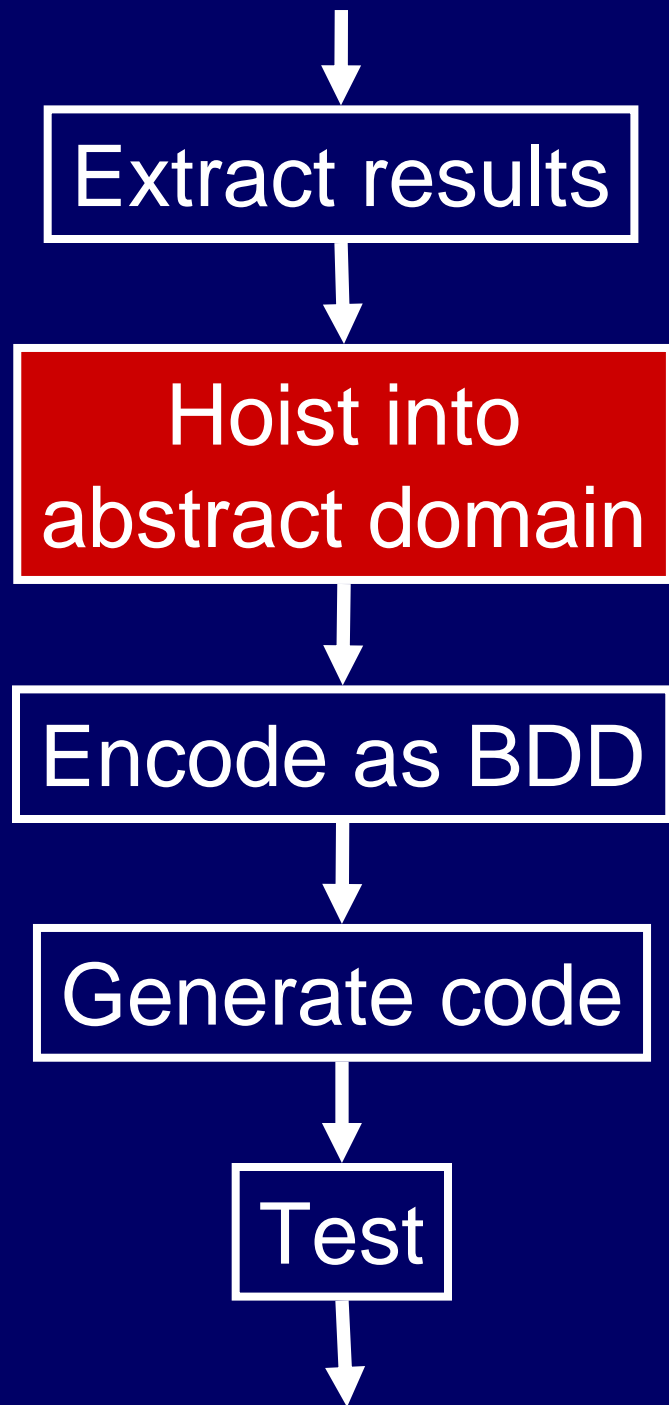
# Hoist Contributions

- Derive transfer functions with
  - Near-zero developer effort
  - Maximal precision
  - Sufficient performance
  - High confidence in correctness

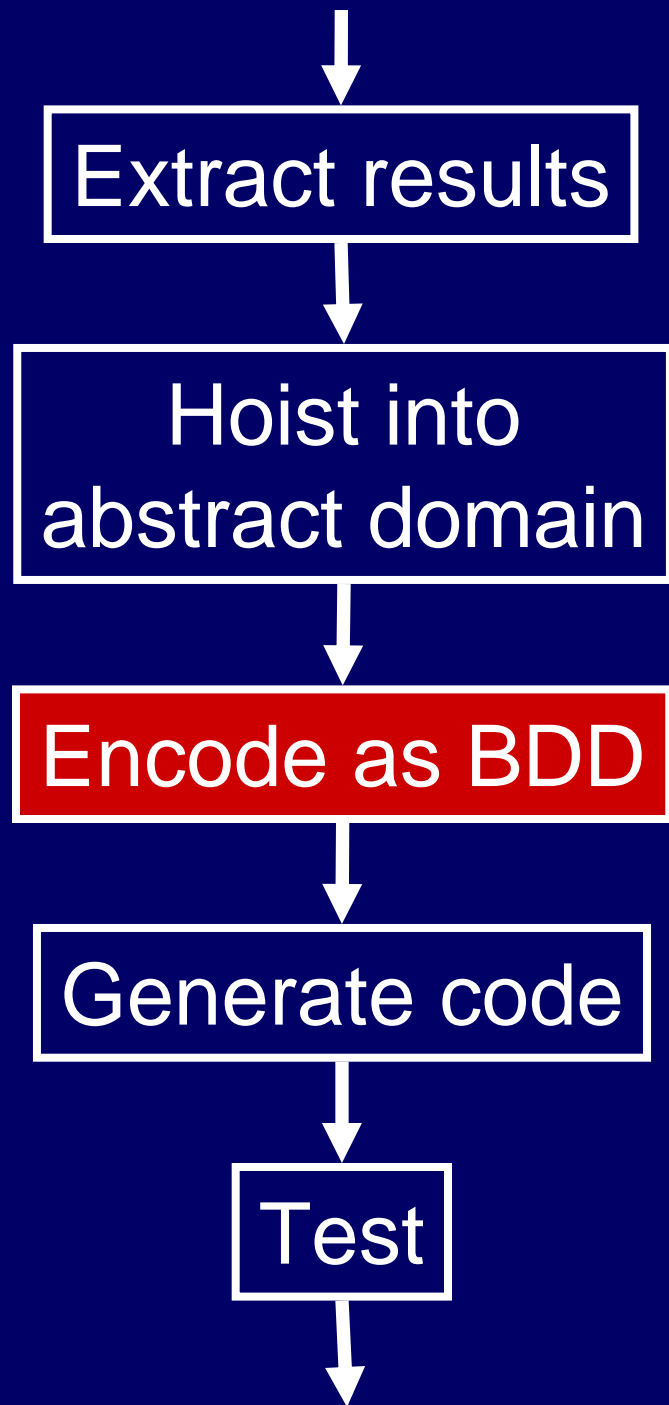




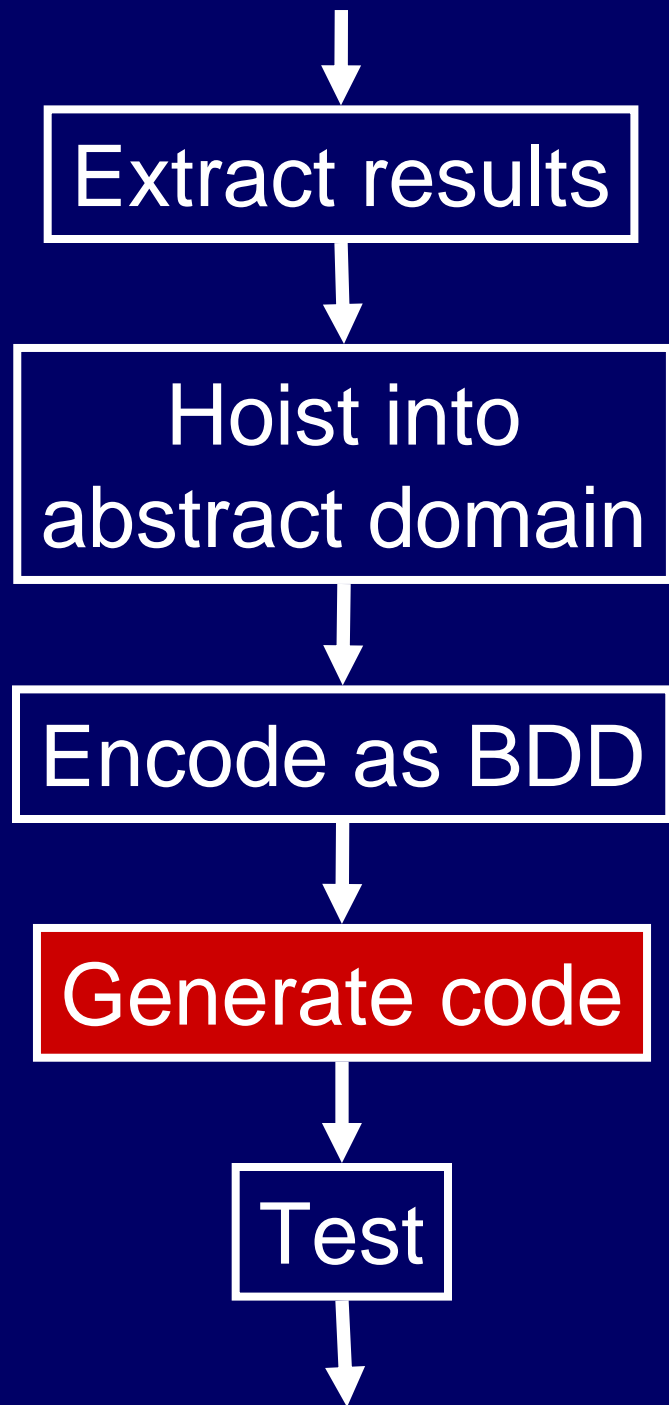
- Extract complete result table for instruction
  - Dest register + cond codes
- Ideas:
  - No high-level model of instruction
  - Brute force



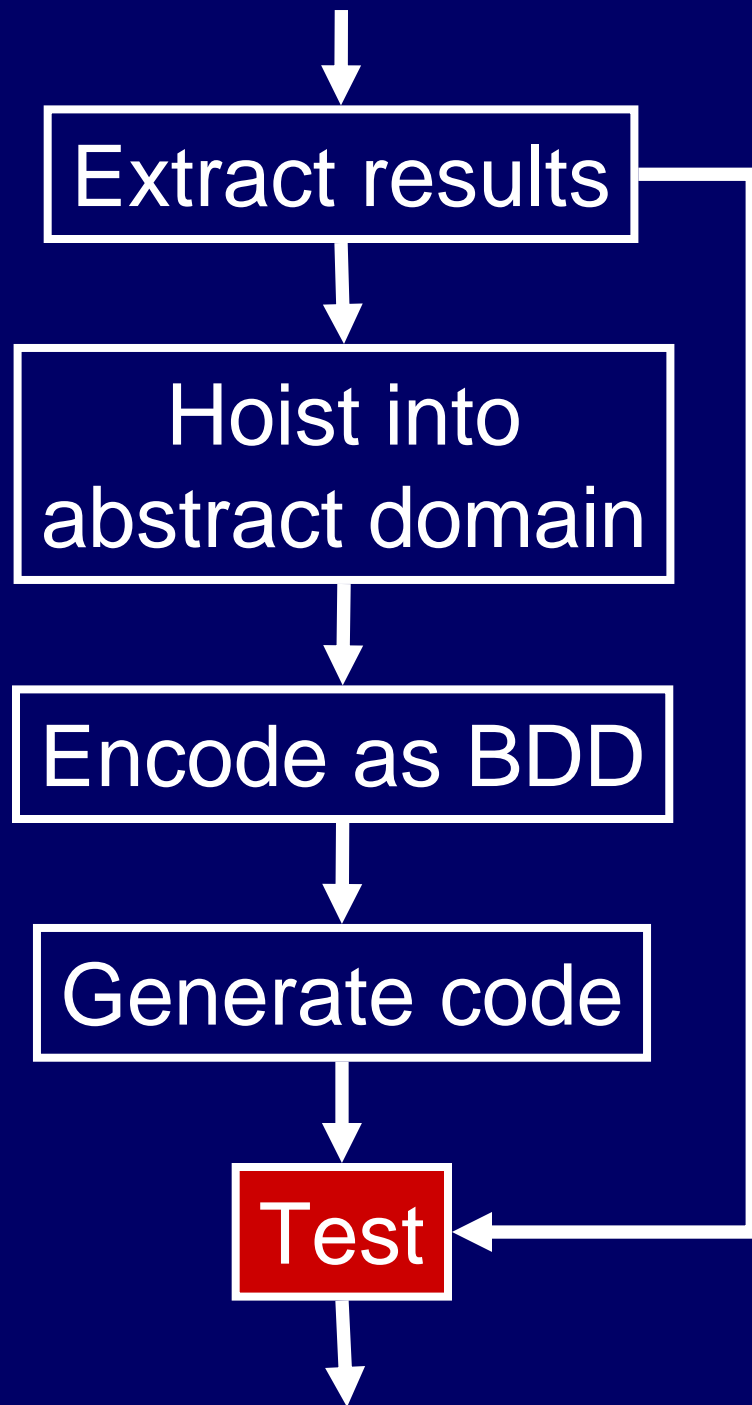
- Generate complete abstract transfer function
- Ideas:
  - Recursive decomposition of abstract domain
  - Speedup through dynamic programming



- Binary decision diagrams can compactly represent many functions
- Encode transfer function as vector of BDDs
- Ideas:
  - Variable ordering matters
  - Operation ordering matters



- Turn BDD into code implementing the transfer function



- Probabilistically or exhaustively verify
  - Correctness
  - Maximal precision
- Original result table is ground-truth

# Hoisting Atmel AVR Architecture

- Up to 45 minutes to Hoist a bitwise operation
- Up to 34 hours to Hoist an interval operation
- Dominated by BDD library
- Parallelizes trivially across operations

# Performance at Analysis Time

- Analyze programs that ship with TinyOS for worst-case stack depth
  - Analysis time increases from 8.3s to 8.9s for the program that takes longest to analyze

# Precision in Bitwise Domain

- Fed random bitwise values to Hoisted and hand-written operations
  - 59% more known bits in result register
  - 130% more known bits in condition codes
- Analyzed 26 TinyOS programs
  - 8% more known bits in result register
  - 40% more known bits in condition codes
- Hand-written operations had been tuned for months



# Twist #1: Pseudo-Unary Ops

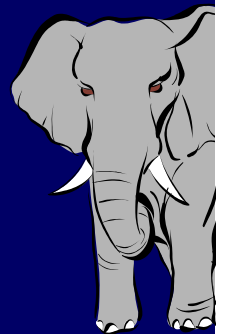
- Problem:
  - `xor 0?10??11, 0?10??11 == 0?00??00`However:
  - `xor r3, r3 == 00000000`
  - Oops! Maximal precision doesn't help here
- Solution: Create a pseudo-unary version of each binary operation
  - E.g. `xor1`, `sub1`, `and1`, `or1`
  - Without these, analysis fails miserably
  - Not fun to implement these by hand

# Twist #2: Interacting Domains

- If a register contains  
[160..210] and ???11011
- We can show that it actually contains  
[187..187] and 10111011
- In general: Use Hoist to create a reduced product of the interval and bitwise domains
  - [Cousot & Cousot 79] says this is impossible
  - For finite domains we can brute-force it
  - Maximally precise

# Elephant in the Closet

- Hoist does not scale to machines bigger than 8 bits
  - 8 bit is important: Many architectures, huge sales volume, used in critical systems
- Current work
  - Replace BDDs with high-level symbolic representation
  - Gain scalability but lose many other advantages of Hoist



# Conclusions

- Reduce barriers to entry for analyzing embedded software
- Hoist generates transfer functions for interval and bitwise domains
  - Near-zero specification effort, maximal precision
- We use Hoisted operations in day-to-day development / use of our static analyzer
  - Biggest benefit is never wondering if the transfer functions are the problem